

Approximate Voronoi cells for lattices, revisited

Thijs Laarhoven

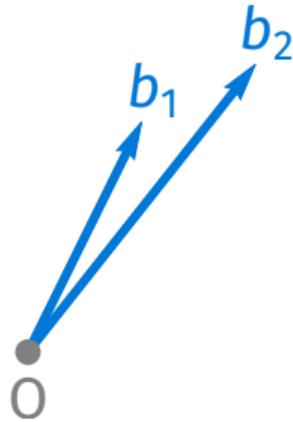
`mail@thijs.com`

`http://www.thijs.com/`

MathCrypt 2019, Santa Barbara, USA
(August 18, 2019)

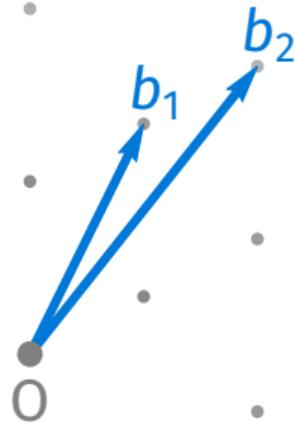
Lattices

Basics



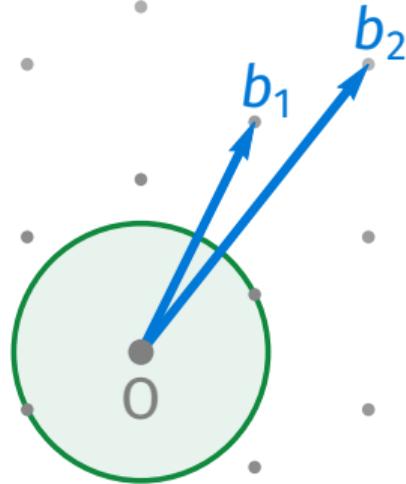
Lattices

Basics



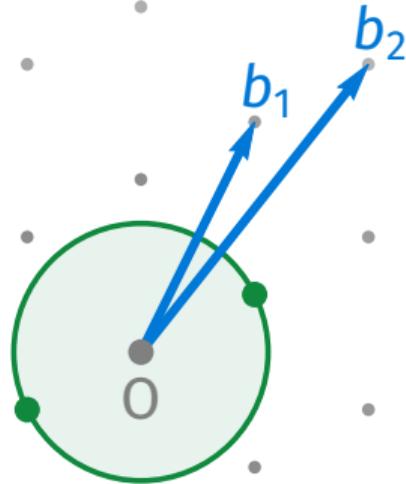
Lattice problems

Shortest Vector Problem (SVP)



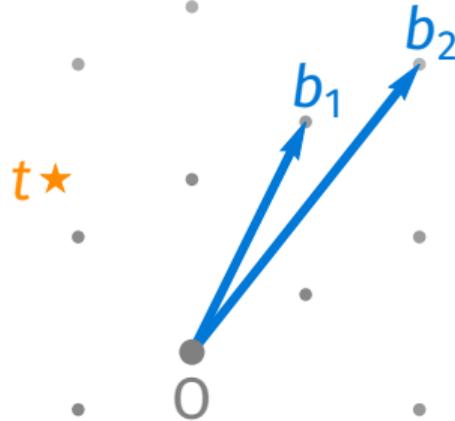
Lattice problems

Shortest Vector Problem (SVP)



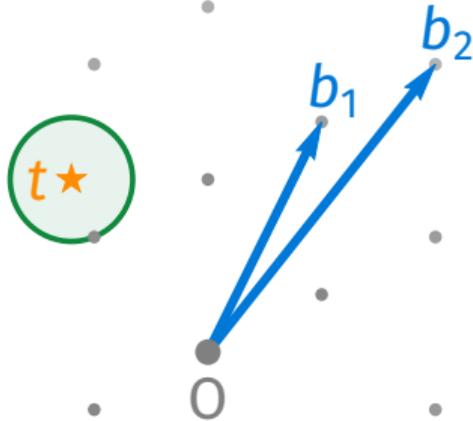
Lattice problems

Closest Vector Problem (CVP)



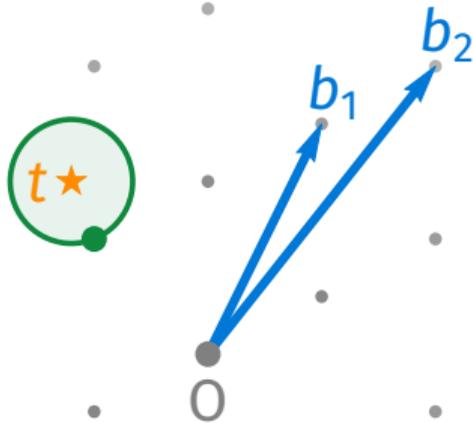
Lattice problems

Closest Vector Problem (CVP)



Lattice problems

Closest Vector Problem (CVP)



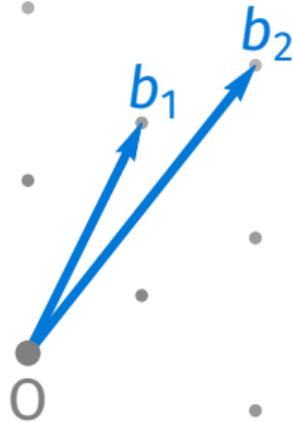
Lattice problems

Asymptotics for SVP and CVP

	Algorithm	$\log_2(\text{Time})$	$\log_2(\text{Space})$	Experiments
Worst-case SVP	Enumeration [Poh81, Kan83, ..., MW15, AN17]	$O(n \log n)$	$O(\log n)$	152
	AKS-sieve [AKS01, NV08, MV10, HPS11]	$3.398n$	$1.985n$	–
	Birthday sieves [PS09, HPS11]	$2.465n$	$1.233n$	–
	Enumeration/DGS hybrid [CCL17]	$2.048n$	$0.500n$	–
	Voronoi cell algorithm [AEVZ02, MV10b, BD15]	$2.000n$	$1.000n$	40
	Quantum sieve [LMP13, LMP15]	$1.799n$	$1.286n$	–
	Quantum enum/DGS [CCL17]	$1.256n$	0.500n	–
	Discrete Gaussian sampling [ADRS15, ADS15, AS18]	1.000n	$1.000n$	–
Average-case SVP	The Nguyen–Vidick sieve [NV08]	$0.415n$	$0.208n$	50
	GaussSieve [MV10, ..., IKMT14, BNvdP16, YKYC17]	$0.415n$	$0.208n$	130*
	Triple sieve [BLS16, HK17]	$0.396n$	$0.189n$	80
	Kleinjung sieve [Kle14]	$0.379n$	$0.189n$	116
	Leveled sieving [WLTB11, ZPH13]	$0.378n$	$0.283n$	–
	Overlattice sieve [BGJ14]	$0.377n$	$0.293n$	90
	Triple sieve with NNS [HK17, HKL18]	$0.359n$	0.189n	76
	Single filters [DL17, ADH+19]	$0.349n$	$0.246n$	155
	Hyperplane LSH [Cha02, FBB+14, Laa15, ..., LM18]	$0.337n$	$0.337n$	107
	Hypercube LSH [TT07, Laa17]	$0.322n$	$0.322n$	–
	May–Ozerov NNS [MO15, BGJ15]	$0.311n$	$0.311n$	–
	Quantum sieve [LMP13]	$0.311n$	$0.208n$	–
	Spherical LSH [AINR14, LdW15]	$0.297n$	$0.297n$	–
	Cross-polytope LSH [TT07, AILRS15, BL16, KW17]	$0.297n$	$0.297n$	80
	Spherical LSF [BDGL16, MLB17, ALRW17, DSvW19]	0.292n	$0.292n$	157
Quantum NNS sieve [LMP15, Laa16]	0.265n	$0.265n$	–	

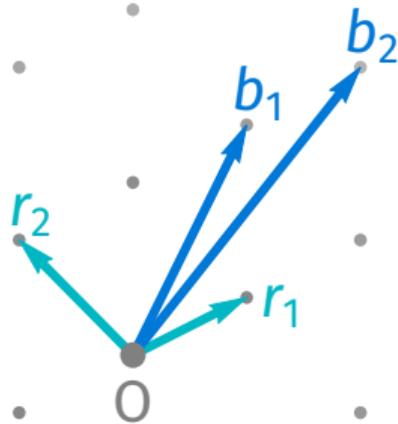
Lattice problems

Closest Vector Problem with Preprocessing (CVPP)



Lattice problems

Closest Vector Problem with Preprocessing (CVPP)



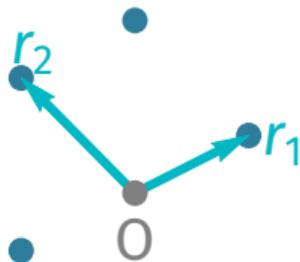
Lattice problems

Closest Vector Problem with Preprocessing (CVPP)



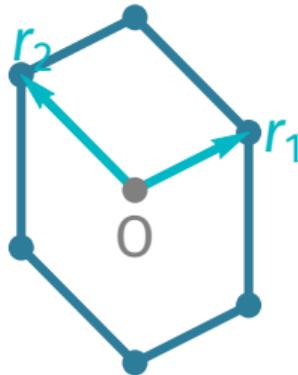
Lattice problems

Closest Vector Problem with Preprocessing (CVPP)



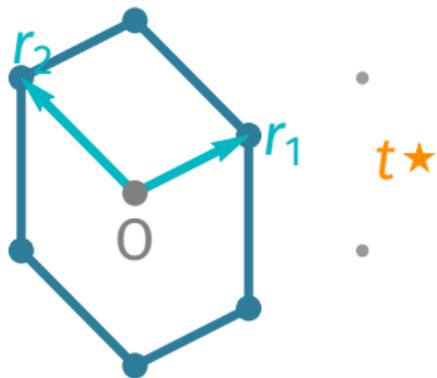
Lattice problems

Closest Vector Problem with Preprocessing (CVPP)



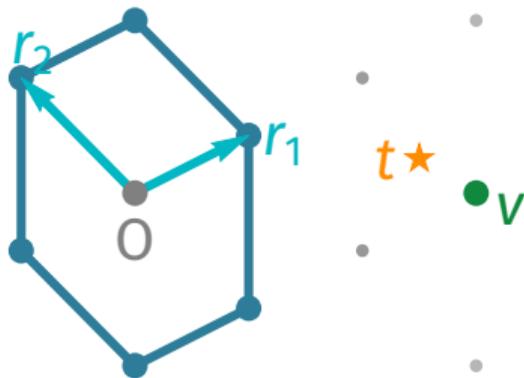
Lattice problems

Closest Vector Problem with Preprocessing (CVPP)



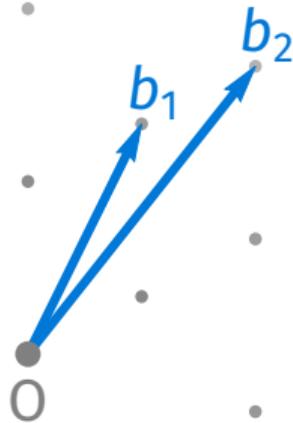
Lattice problems

Closest Vector Problem with Preprocessing (CVPP)



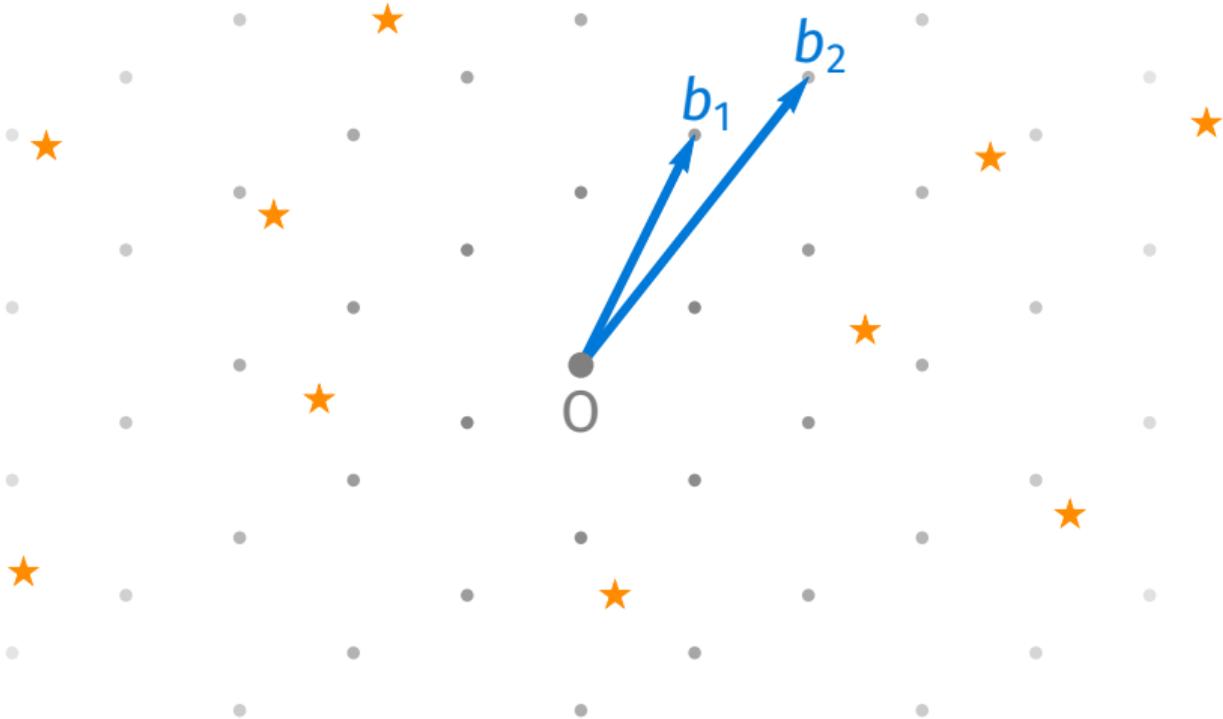
Lattice problems

Batch Closest Vector Problem



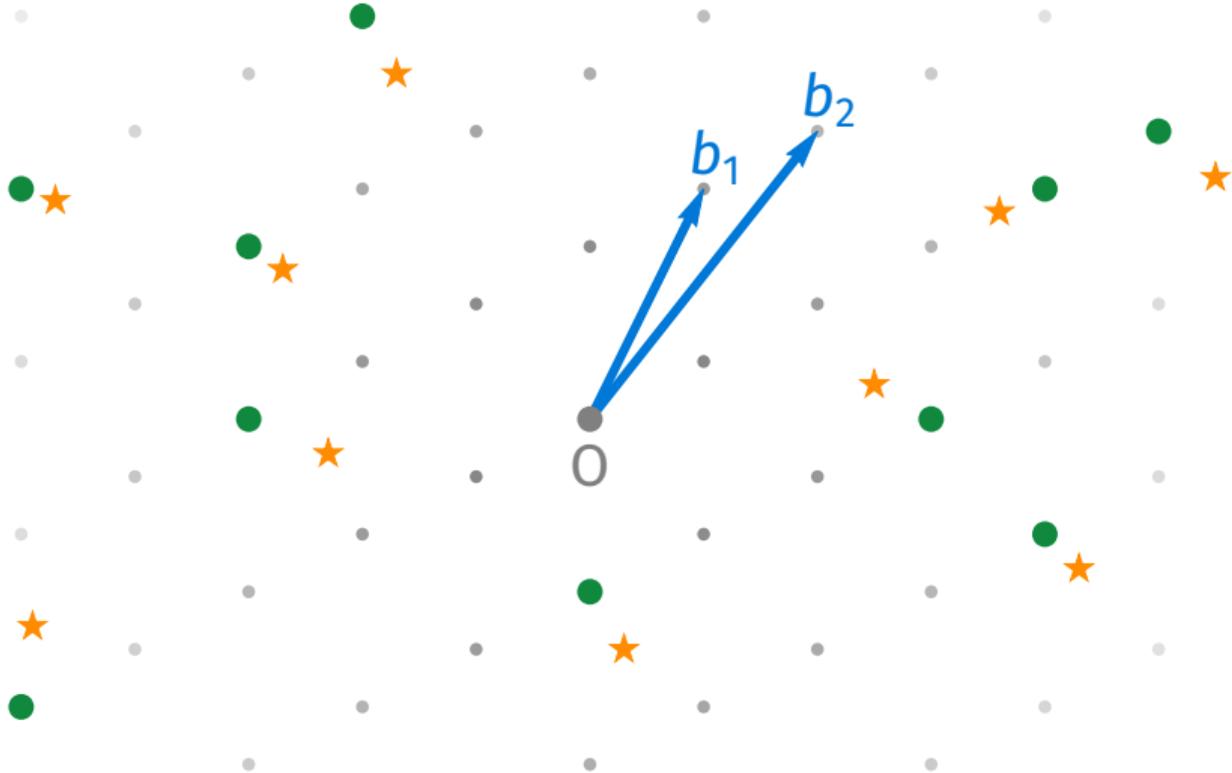
Lattice problems

Batch Closest Vector Problem



Lattice problems

Batch Closest Vector Problem



Lattice problems

Why study CVPP?

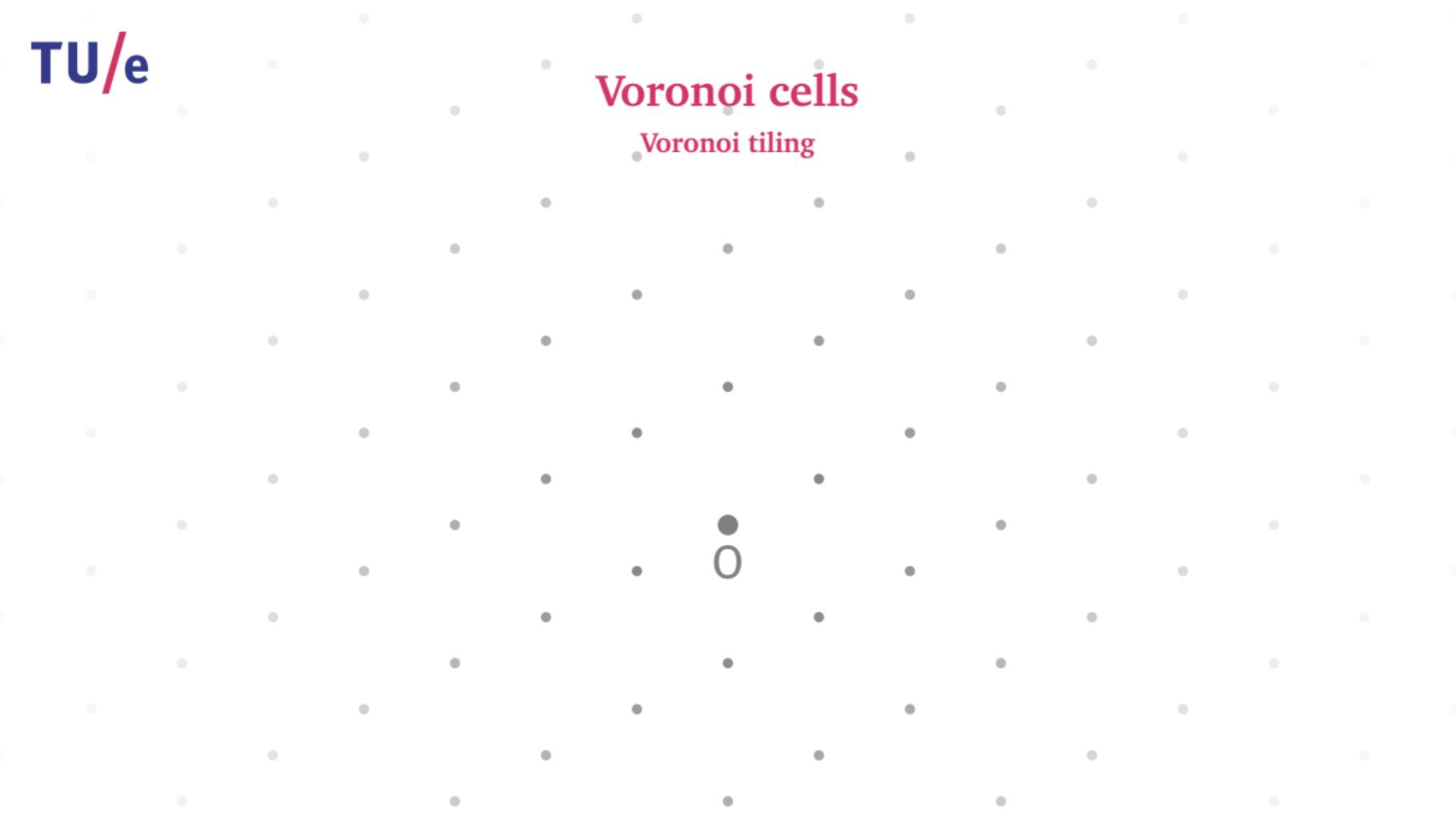
Concrete applications

- Speeding up lattice enumeration for SVP or CVP [GNR10]
- Solving approximate SVP on ideal lattices [PHS19]
- Computing class group actions in a relation lattice [BKV19]

Commonly a lattice basis (public key) is known long before the target vectors (encryptions, signatures)

Voronoi cells

Voronoi tiling



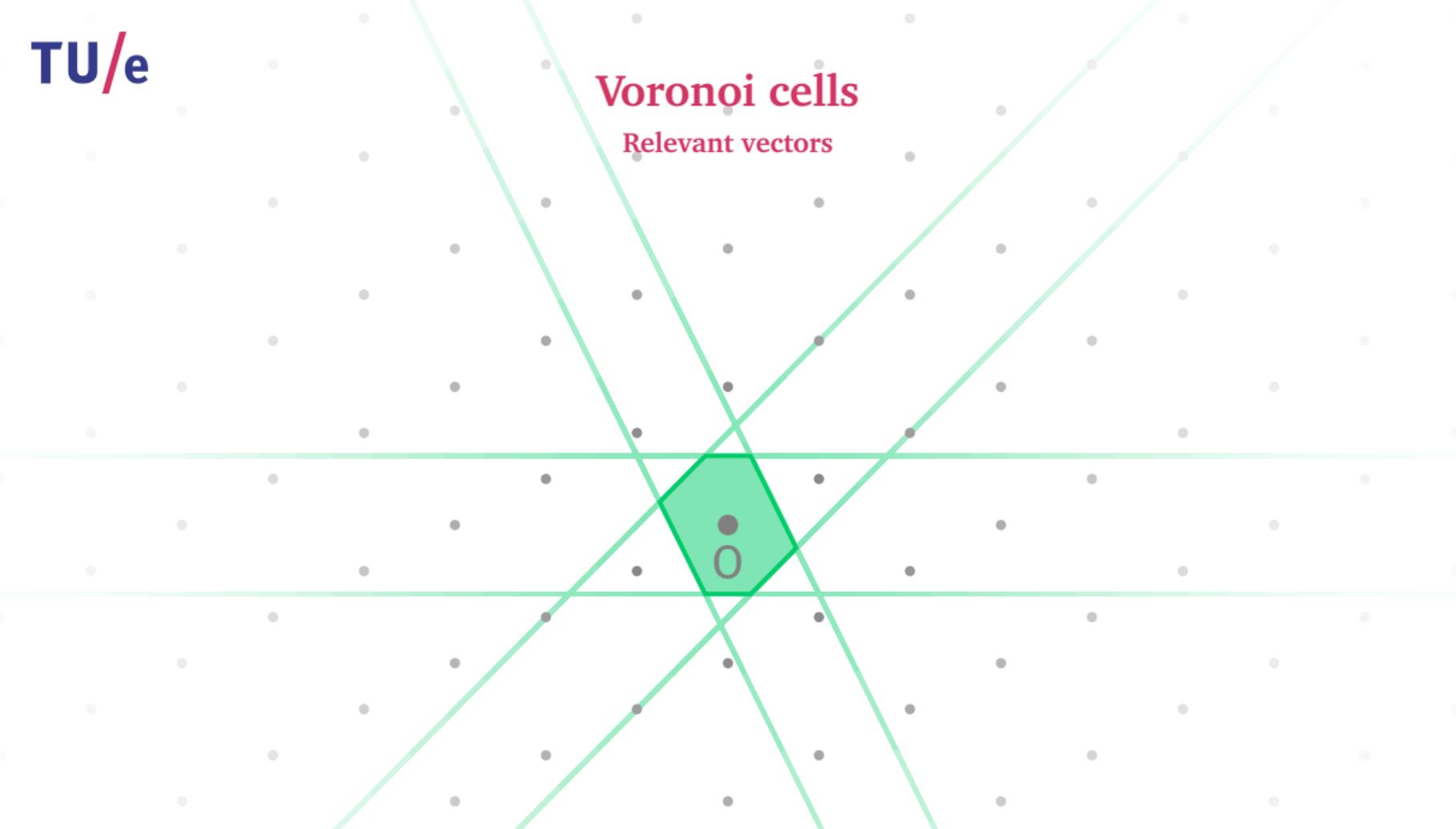
0

Voronoi cells

Voronoi tiling

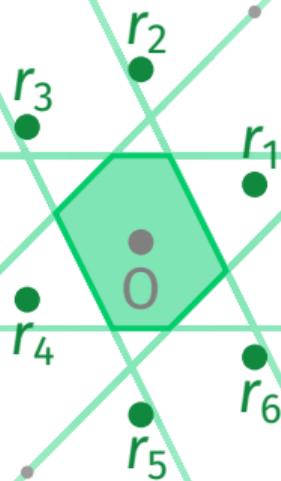


Voronoi cells
Relevant vectors



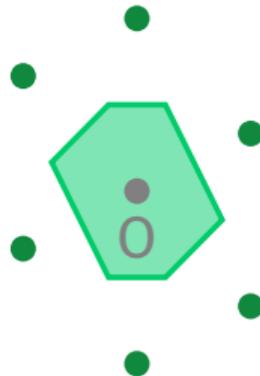
Voronoi cells

Relevant vectors



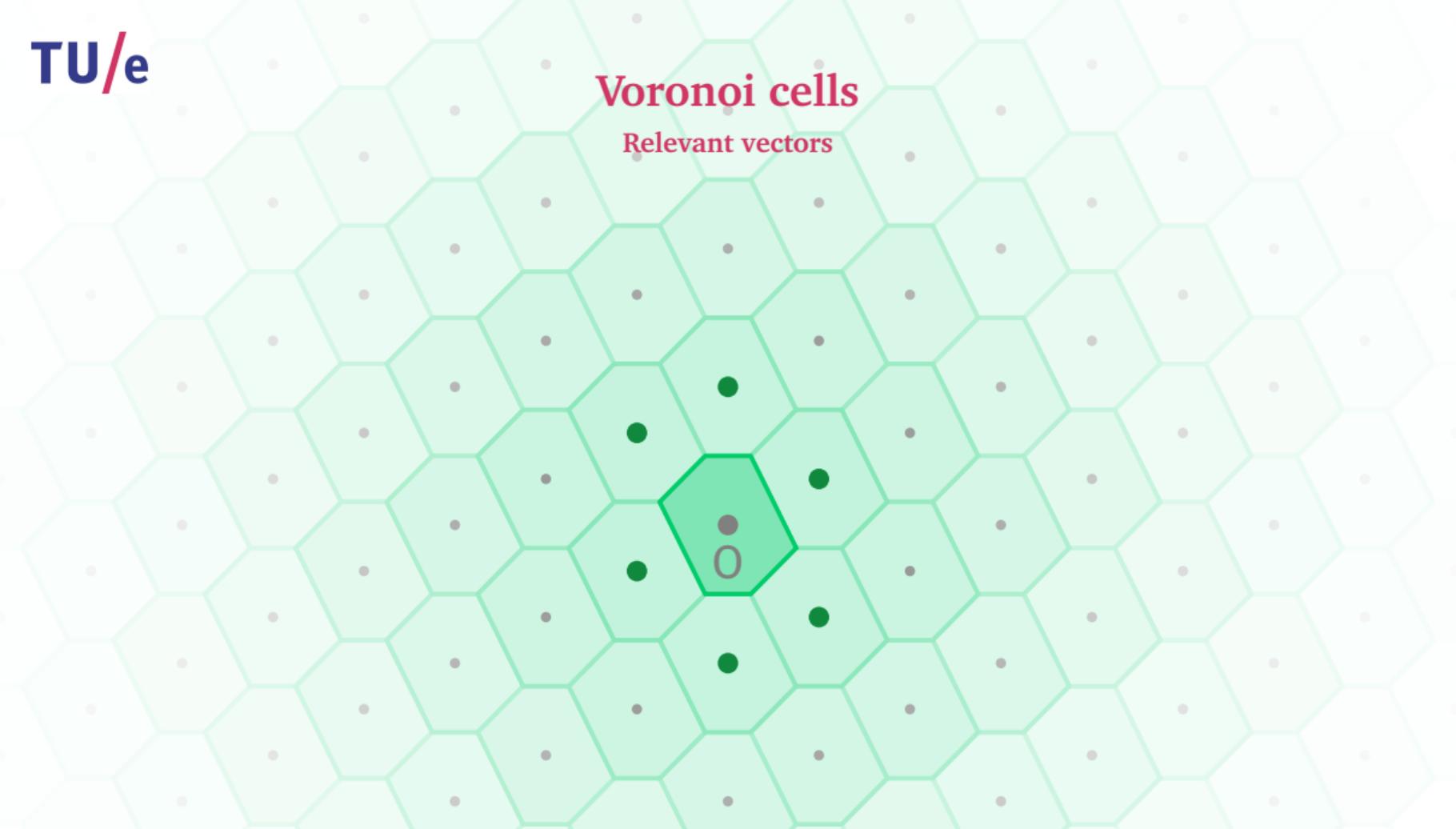
Voronoi cells

Relevant vectors



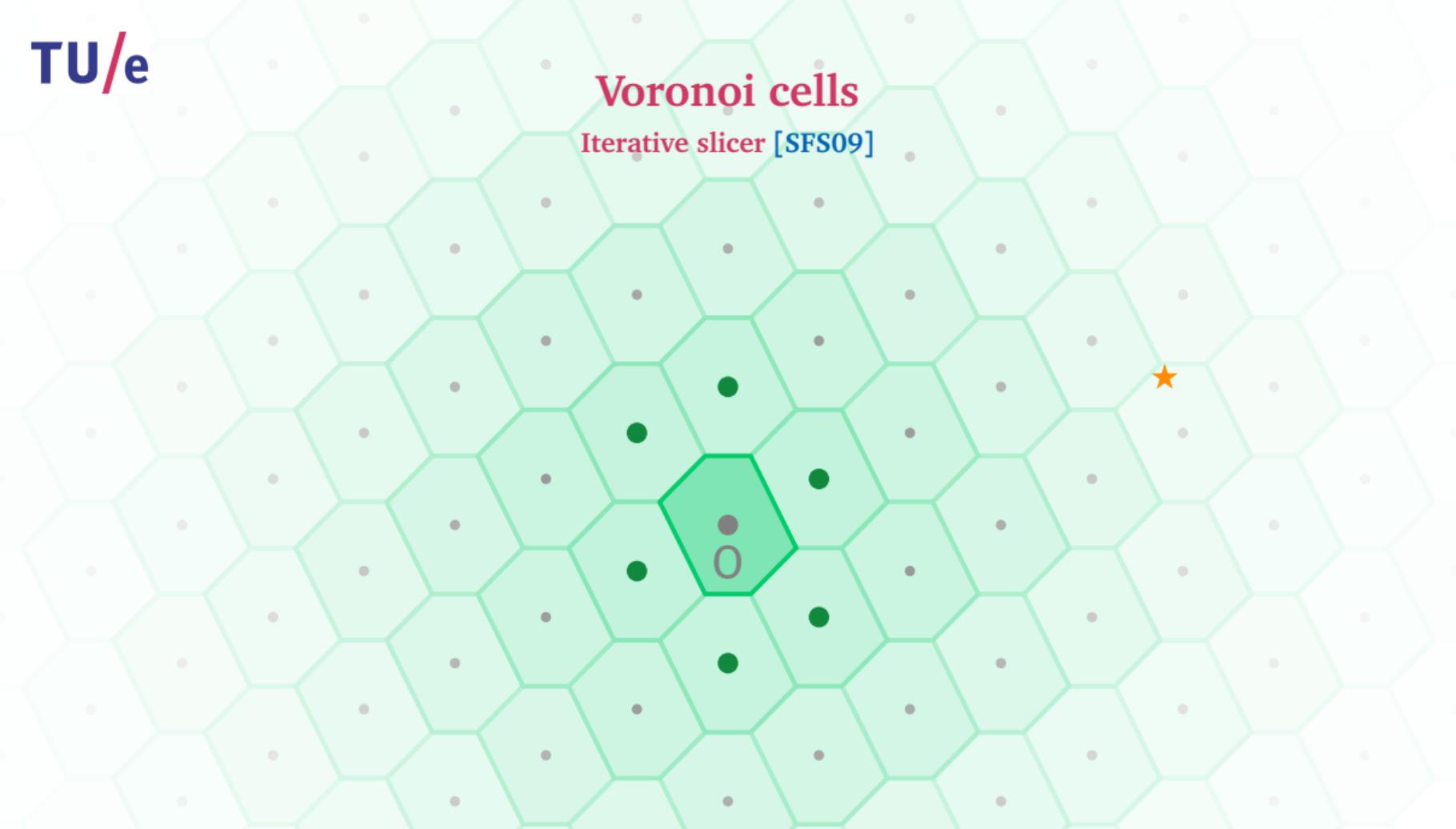
Voronoi cells

Relevant vectors



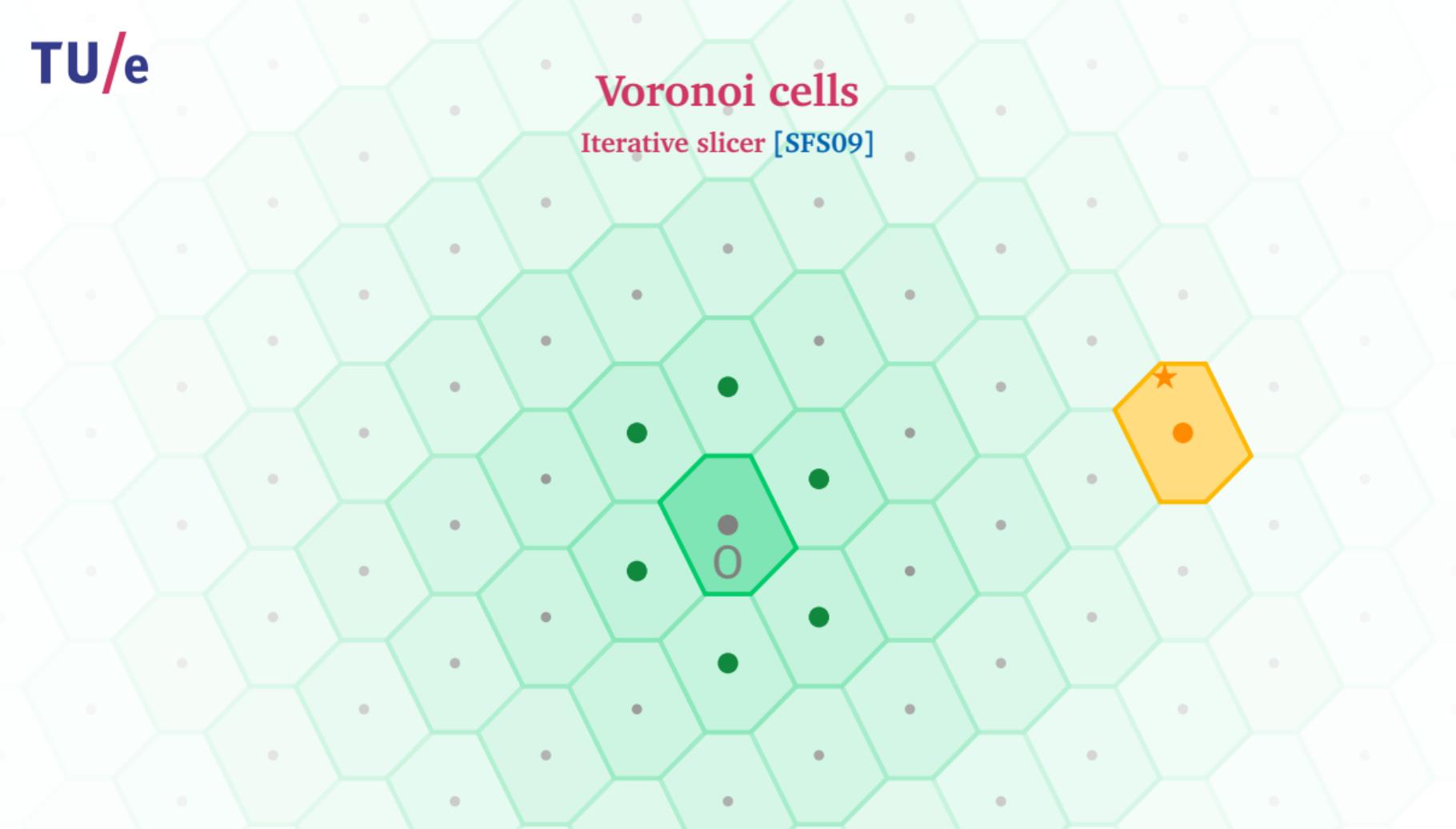
Voronoi cells

Iterative slicer [SFS09]



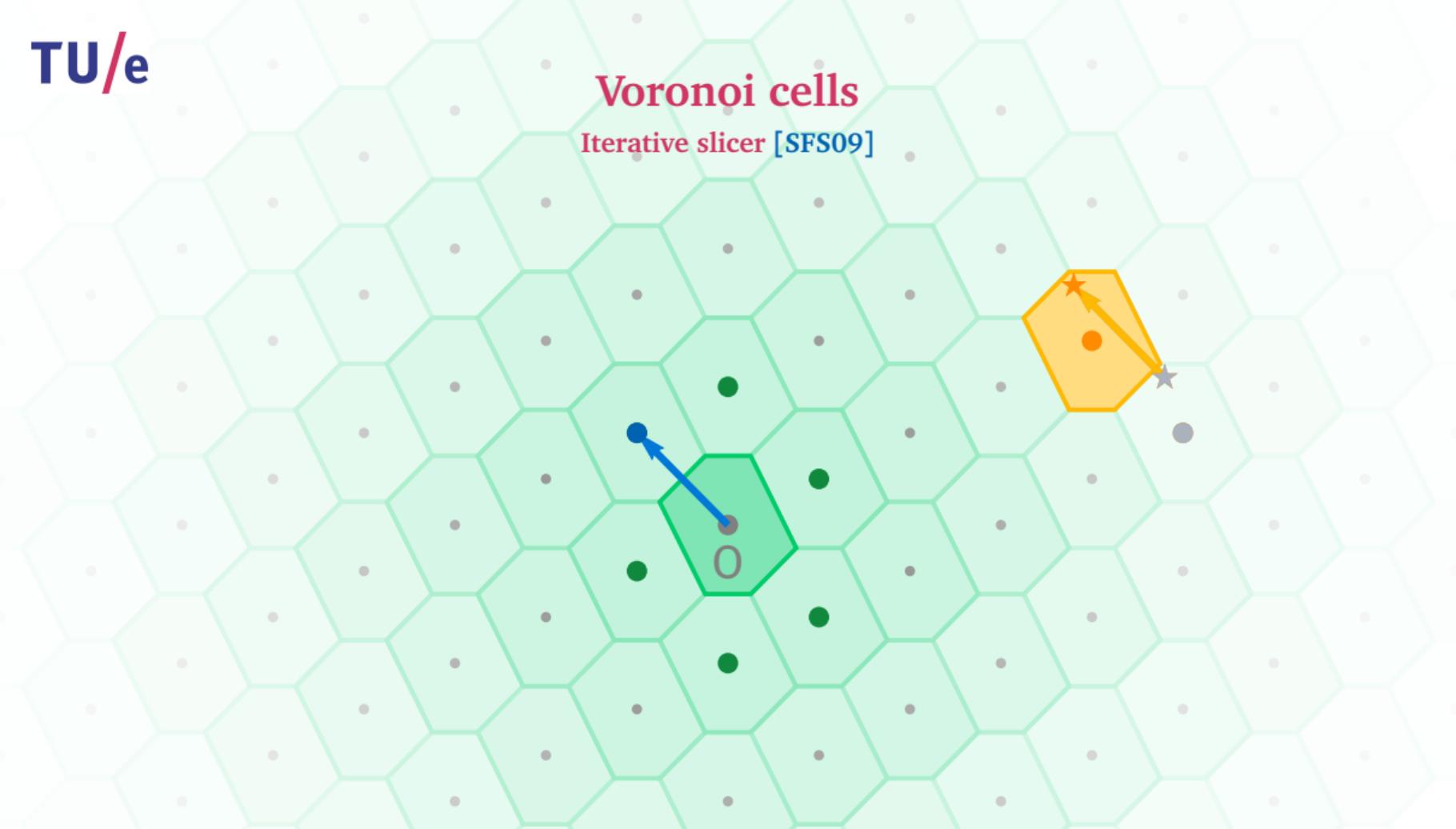
Voronoi cells

Iterative slicer [SFS09]



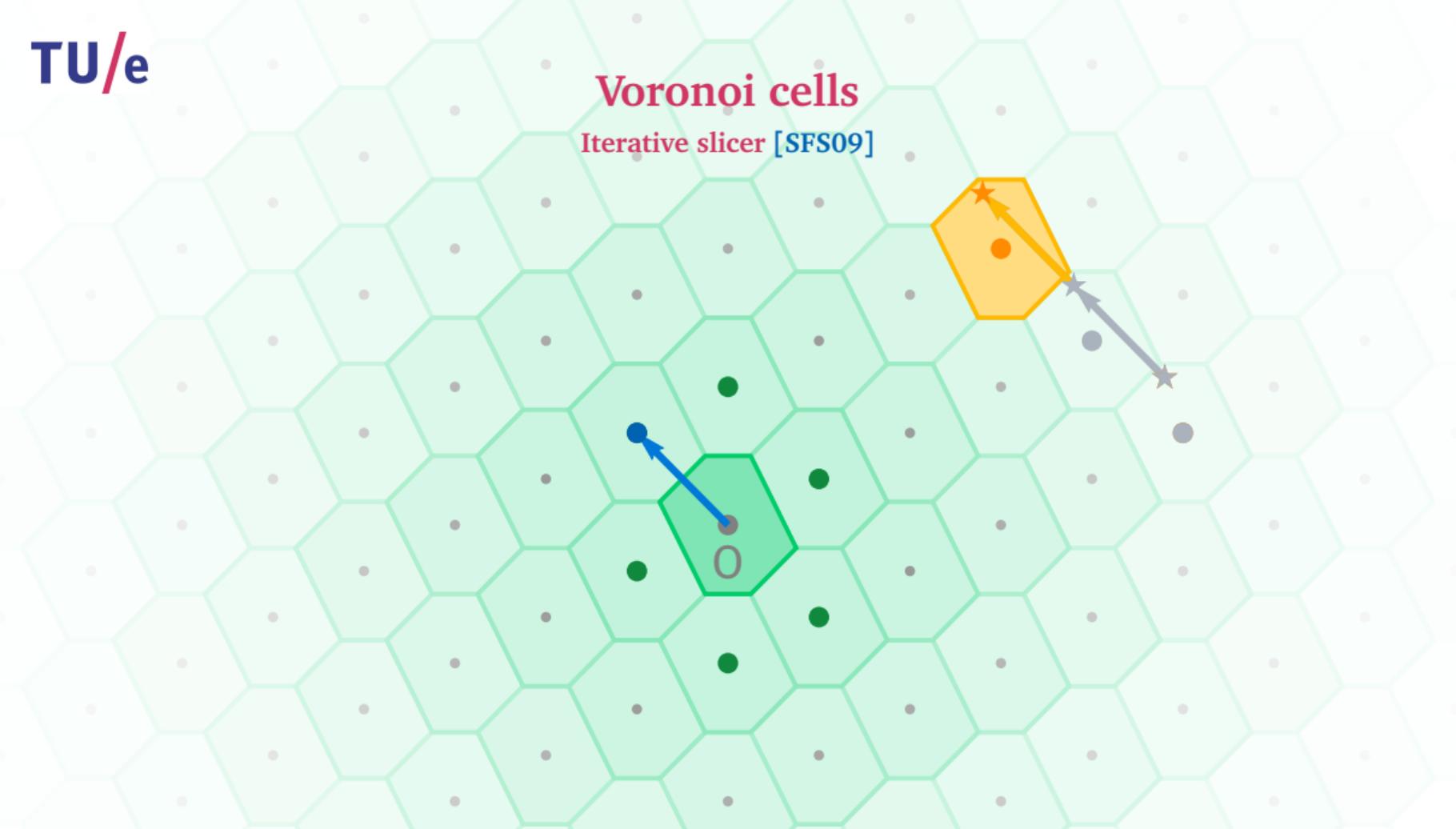
Voronoi cells

Iterative slicer [SFS09]



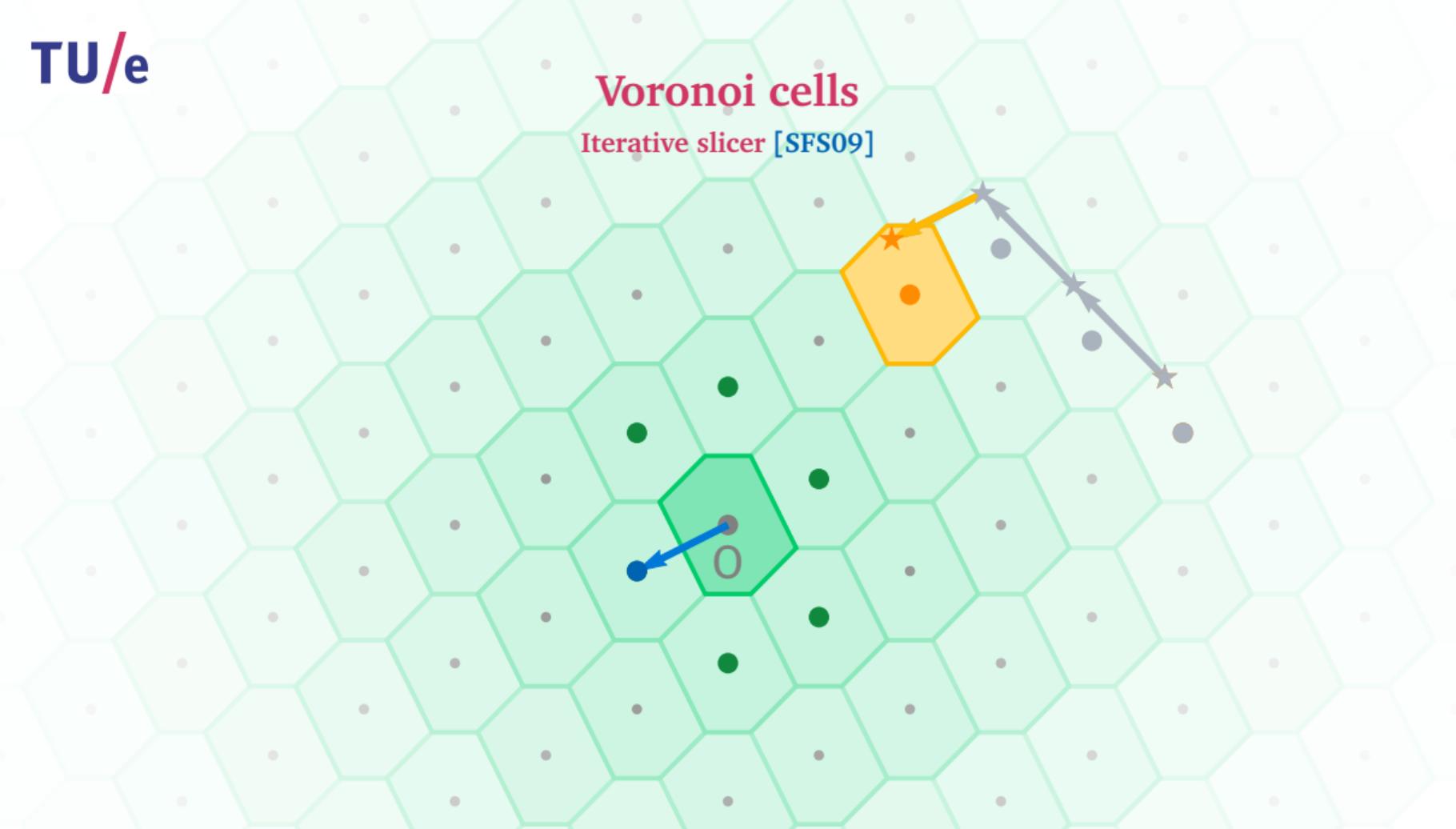
Voronoi cells

Iterative slicer [SFS09]



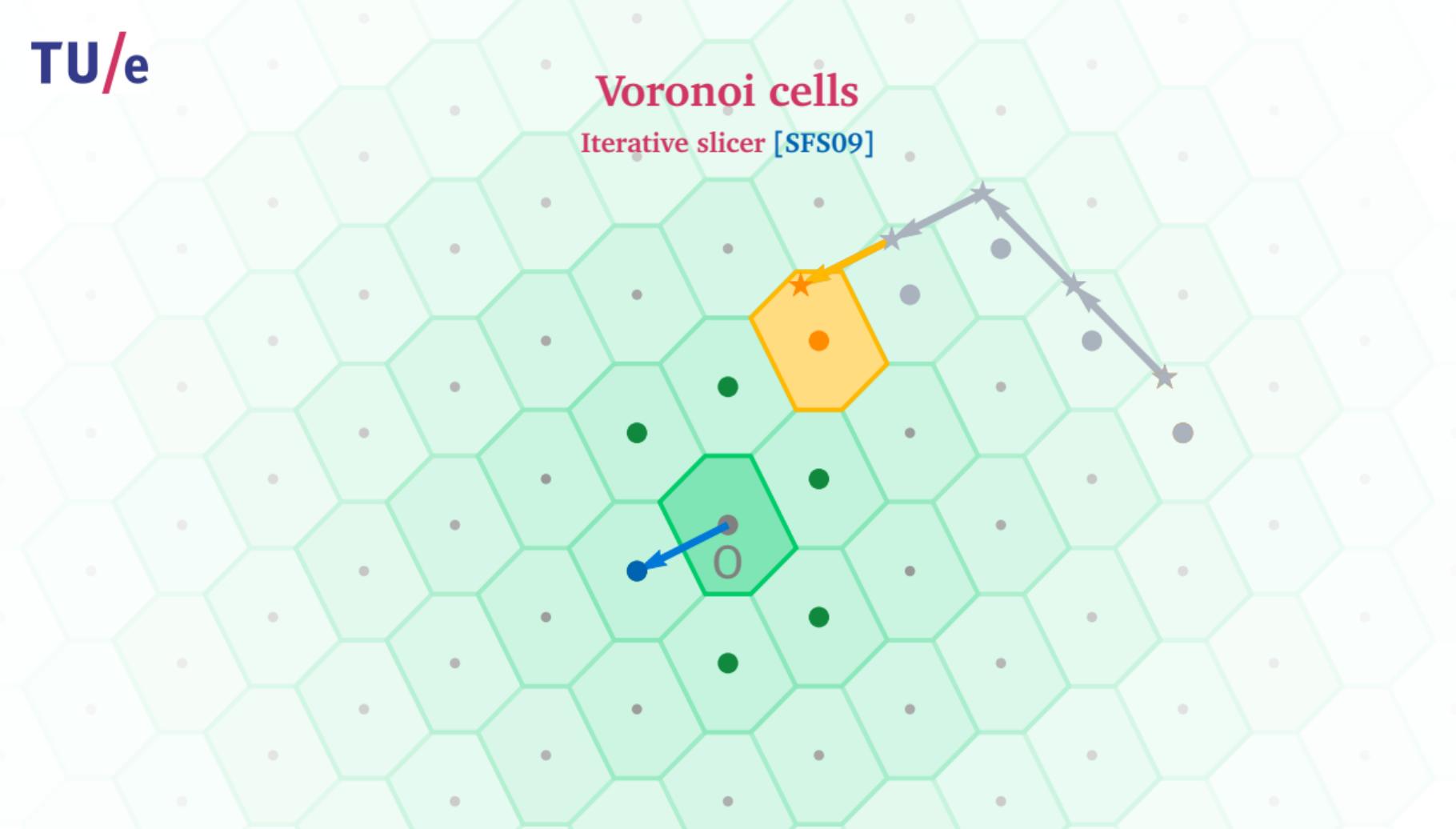
Voronoi cells

Iterative slicer [SFS09]



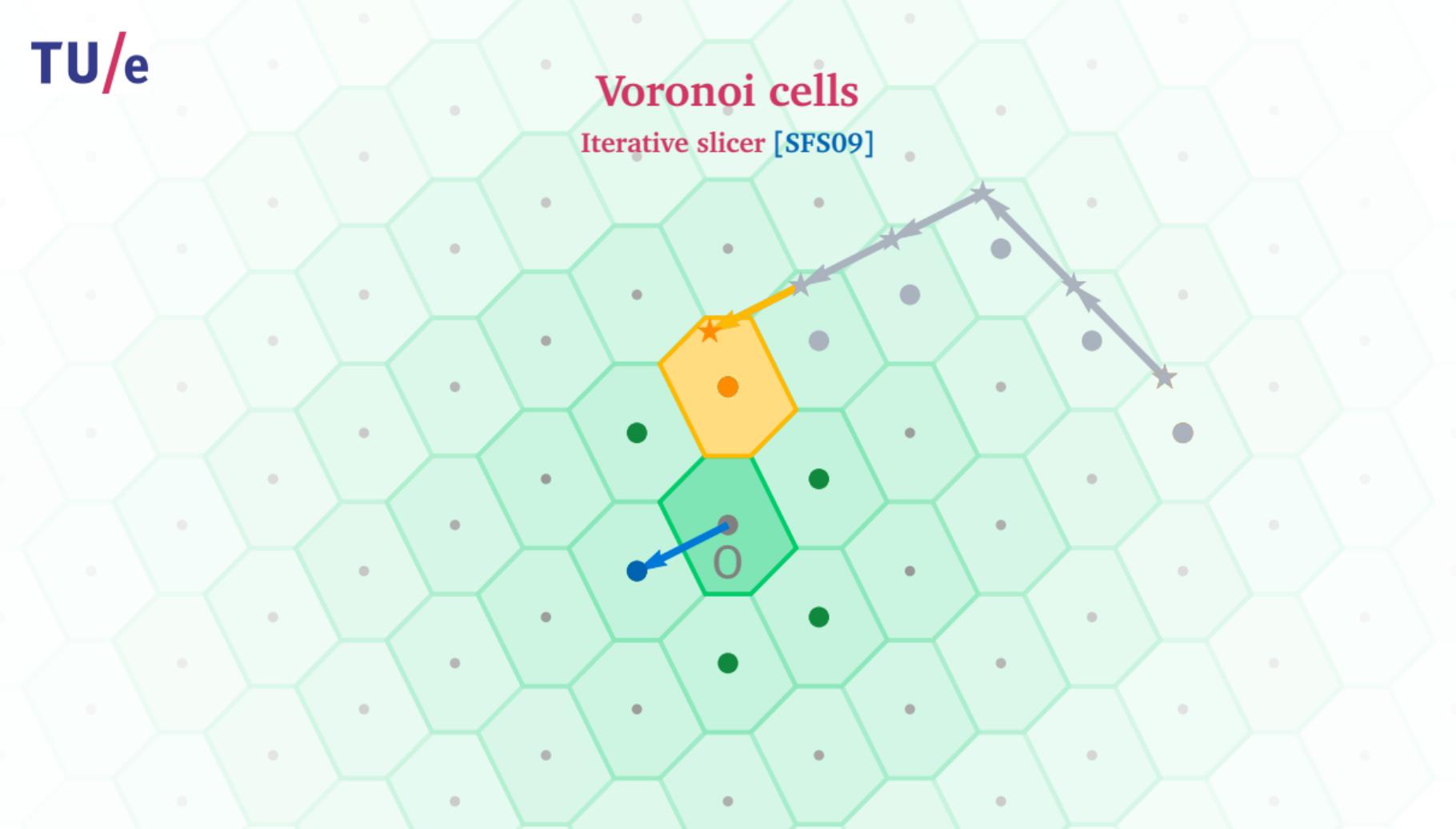
Voronoi cells

Iterative slicer [SFS09]



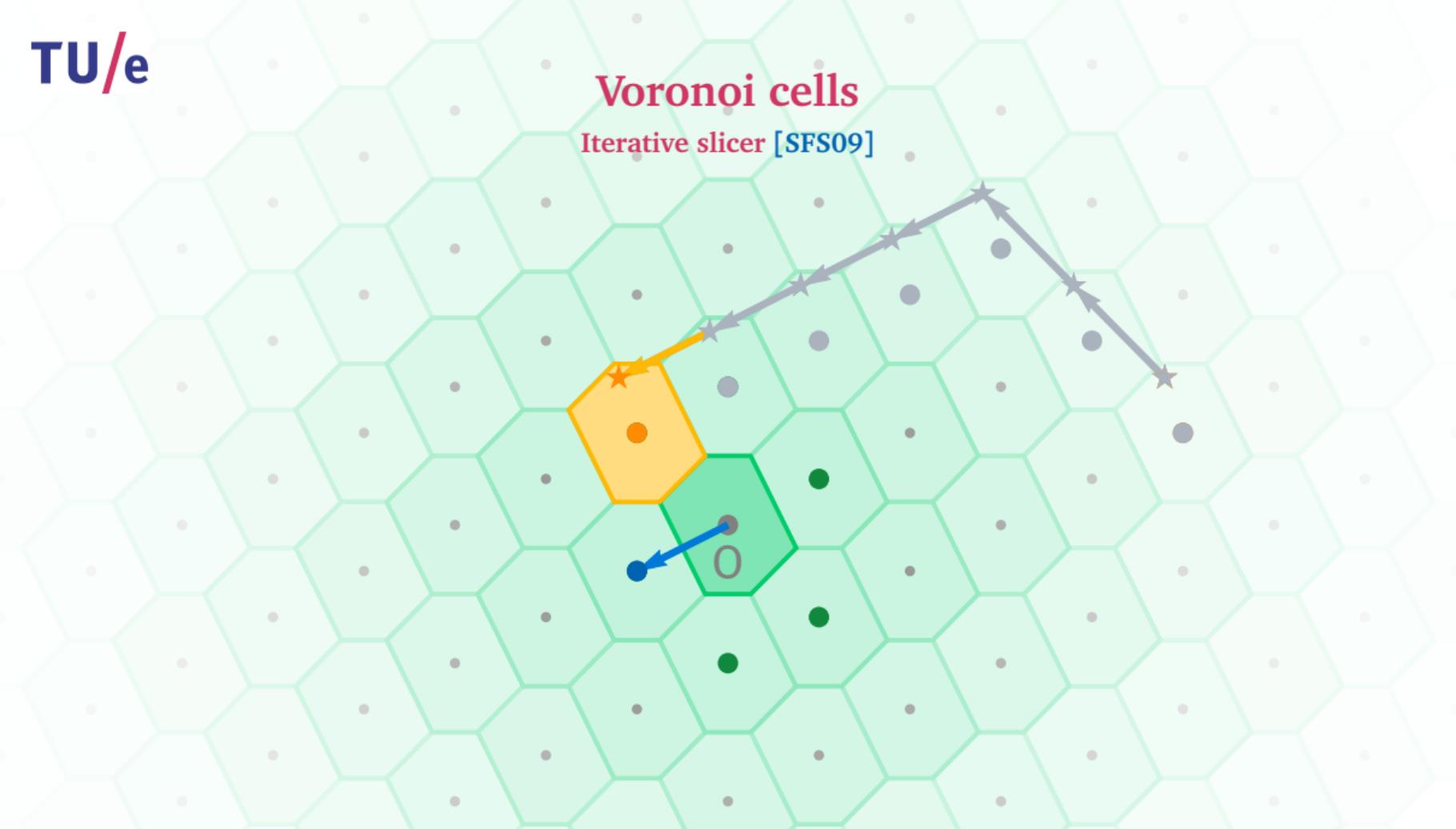
Voronoi cells

Iterative slicer [SFS09]



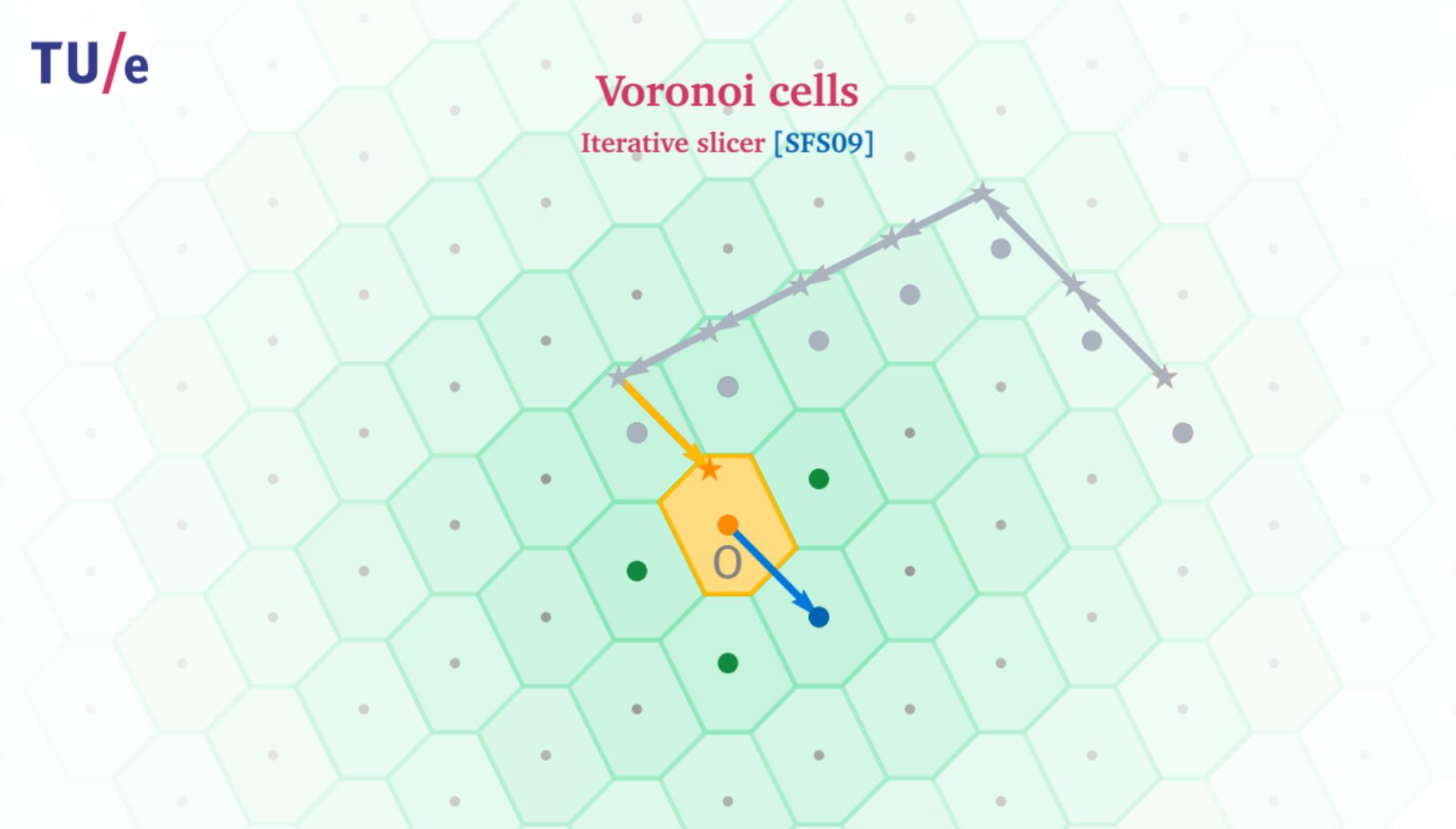
Voronoi cells

Iterative slicer [SFS09]



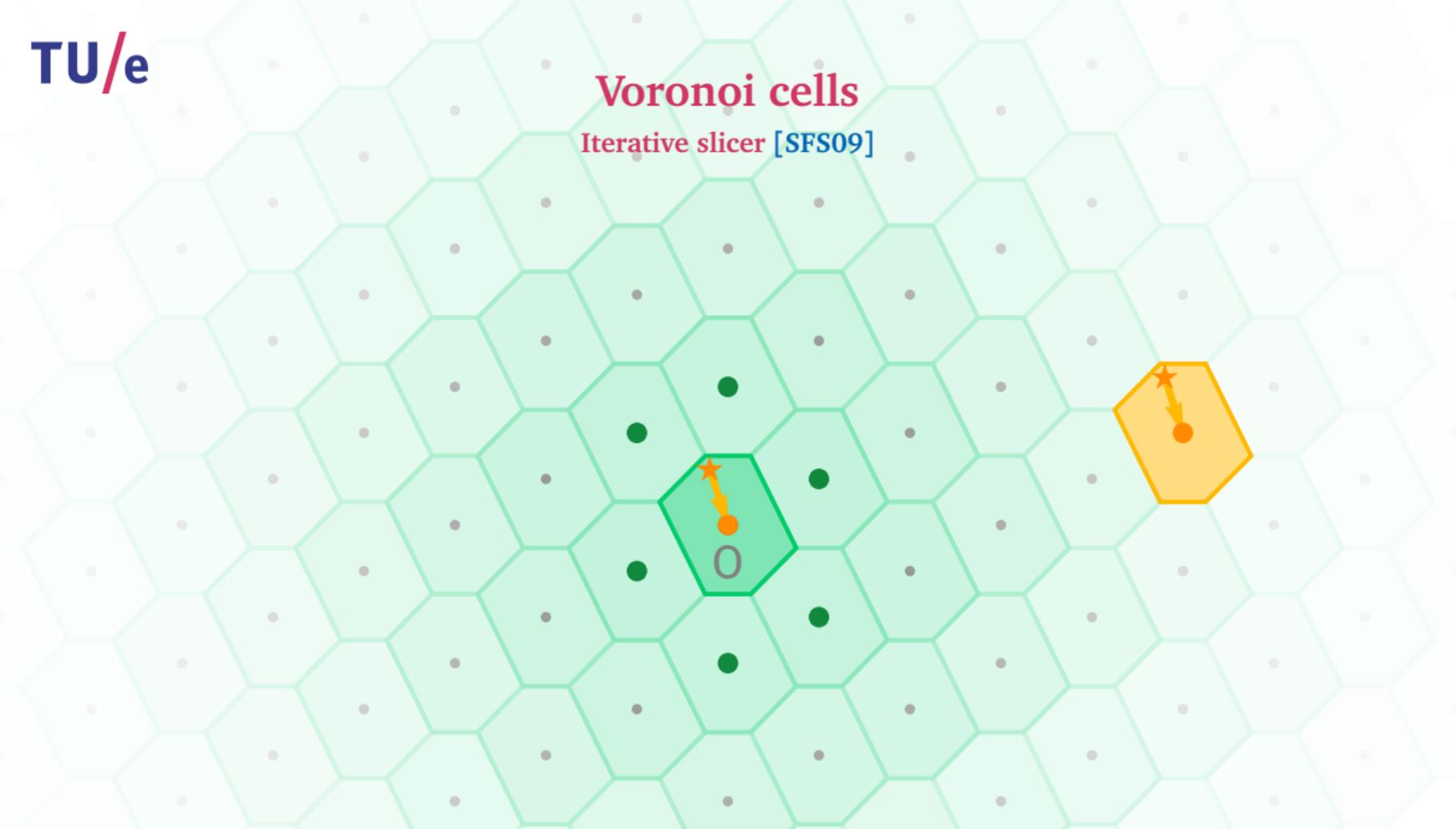
Voronoi cells

Iterative slicer [SFS09]



Voronoi cells

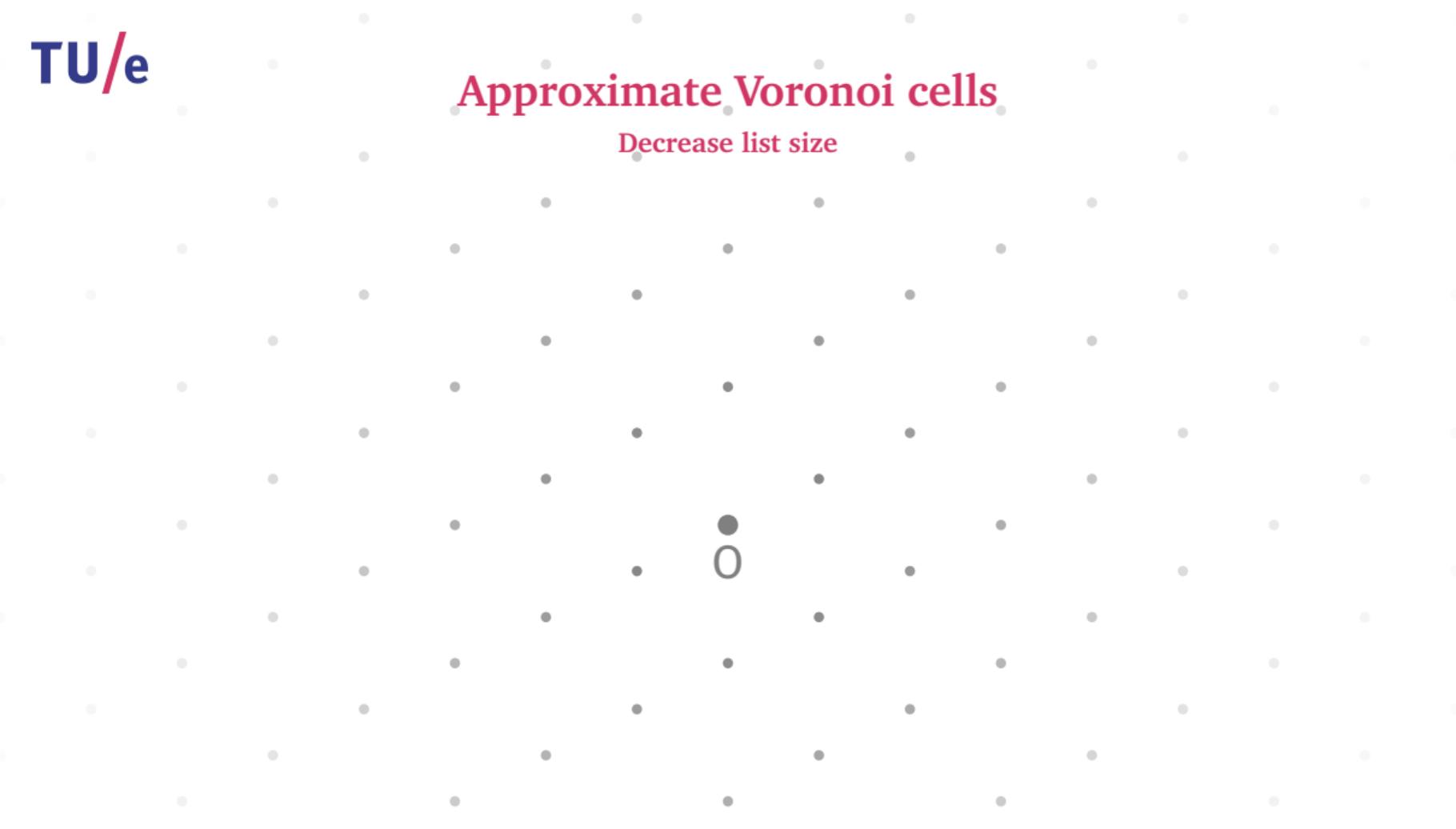
Iterative slicer [SFS09]



Approximate Voronoi cells

Decrease list size

0

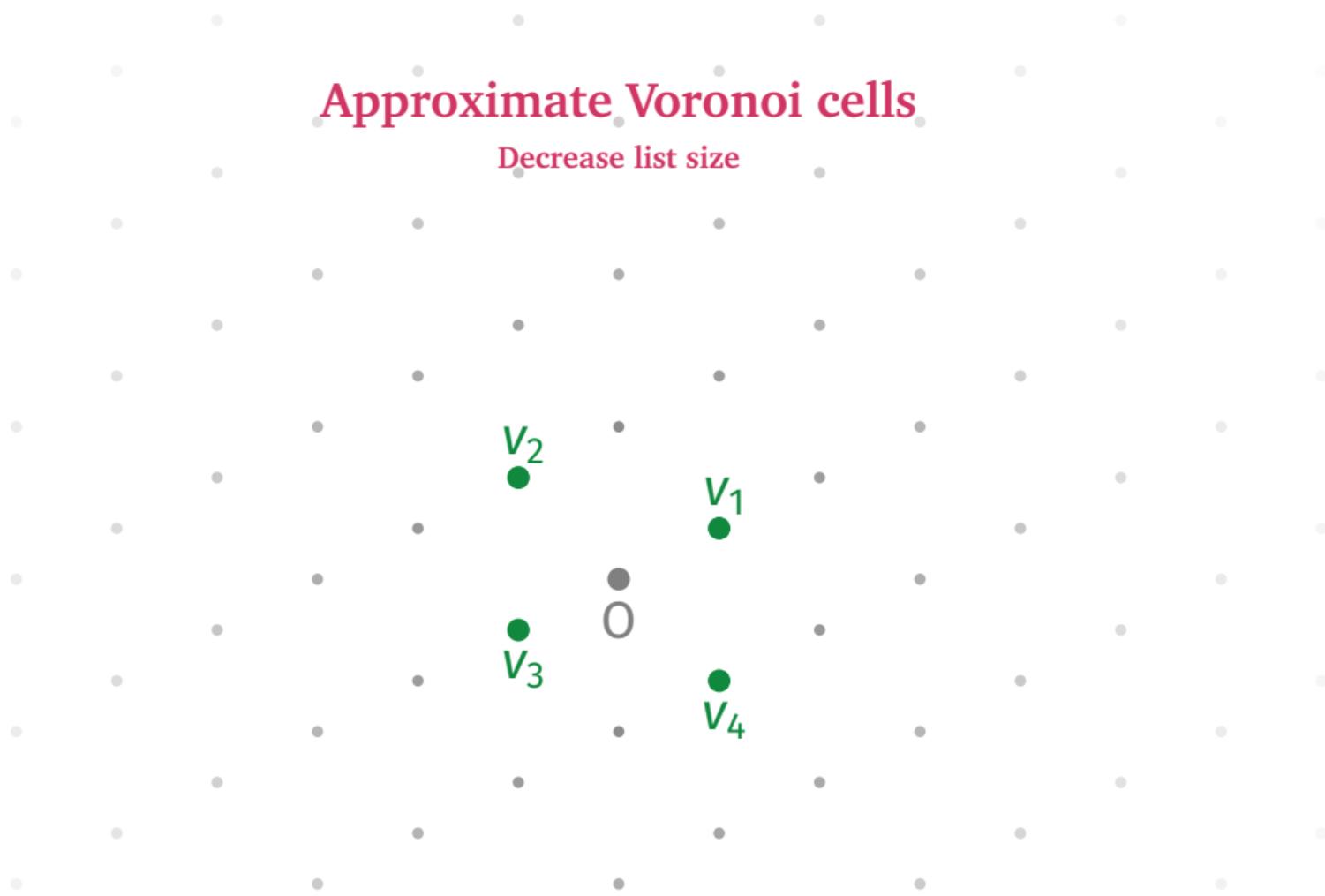
The diagram shows a grid of points on a white background. A central point is labeled '0'. Above it is a larger black dot. The text 'Approximate Voronoi cells' is at the top in red, and 'Decrease list size' is below it in red. The grid of points is composed of small grey dots, with some dots being slightly larger or darker than others, indicating a process of refinement or selection.

Approximate Voronoi cells

Decrease list size

 V_2 V_1 V_3 V_4

0



Approximate Voronoi cells

Decrease list size

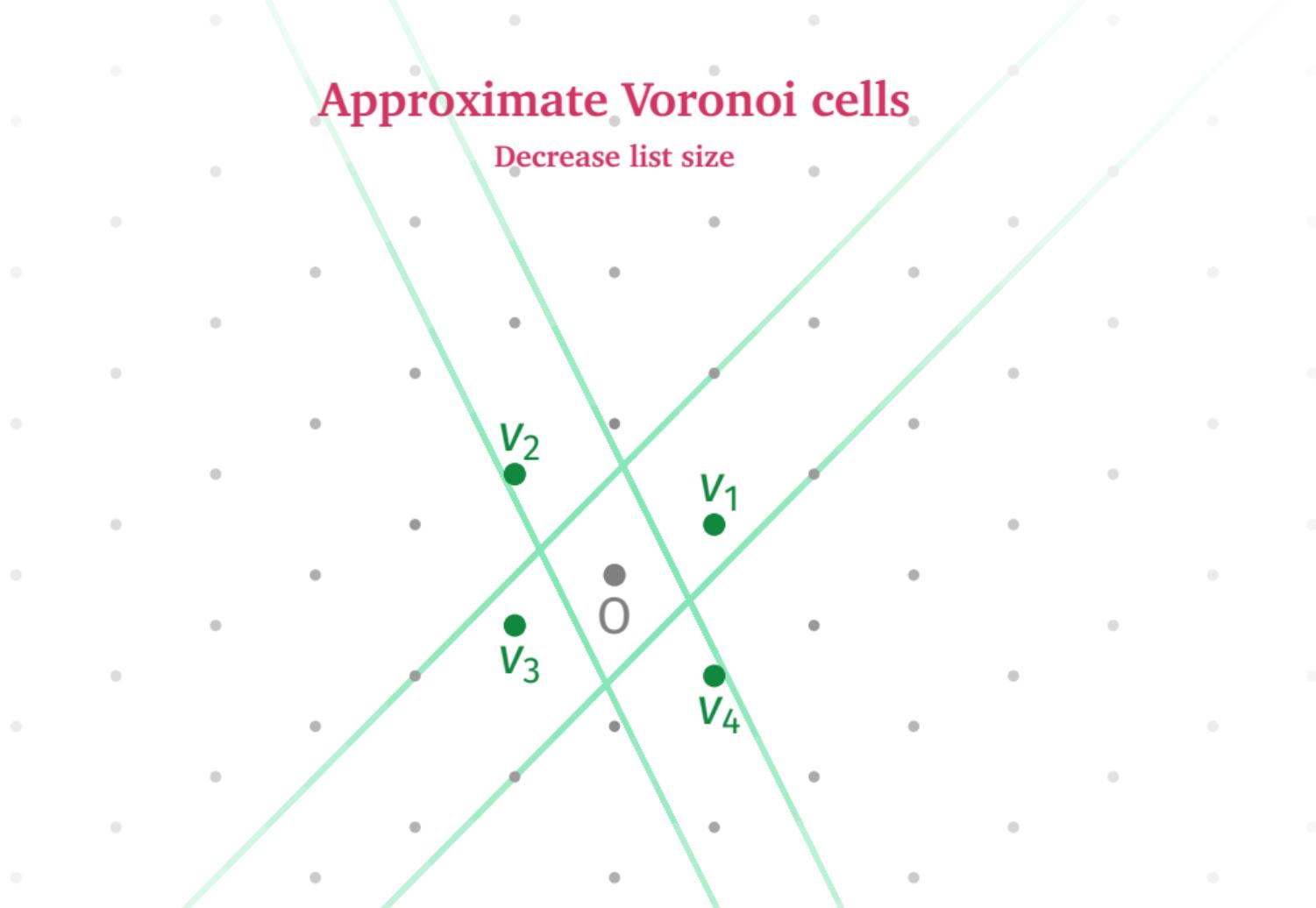
V_2

V_1

0

V_3

V_4



Approximate Voronoi cells

Decrease list size

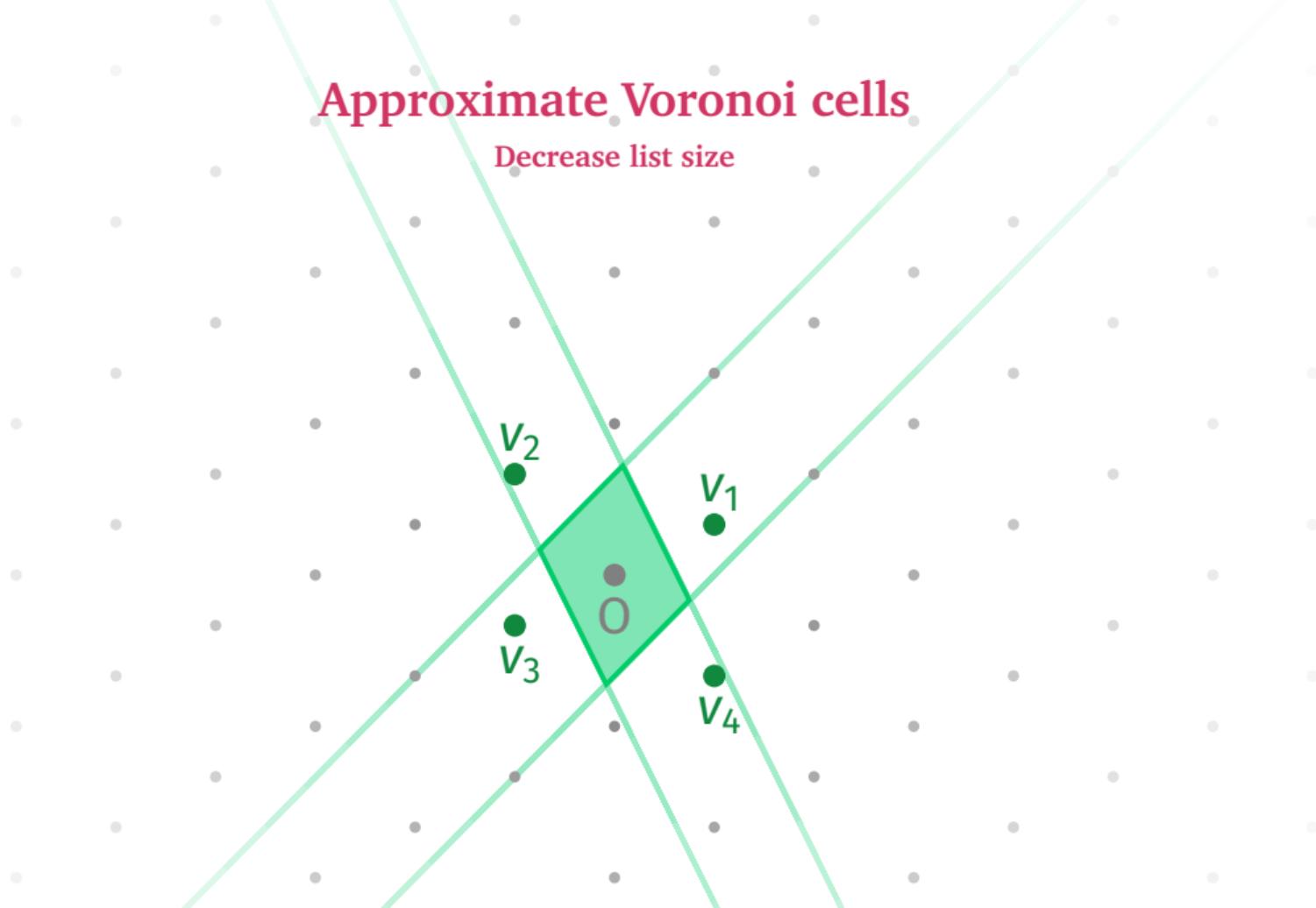
V_2

V_1

O

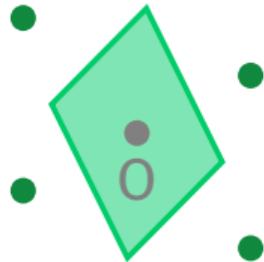
V_3

V_4



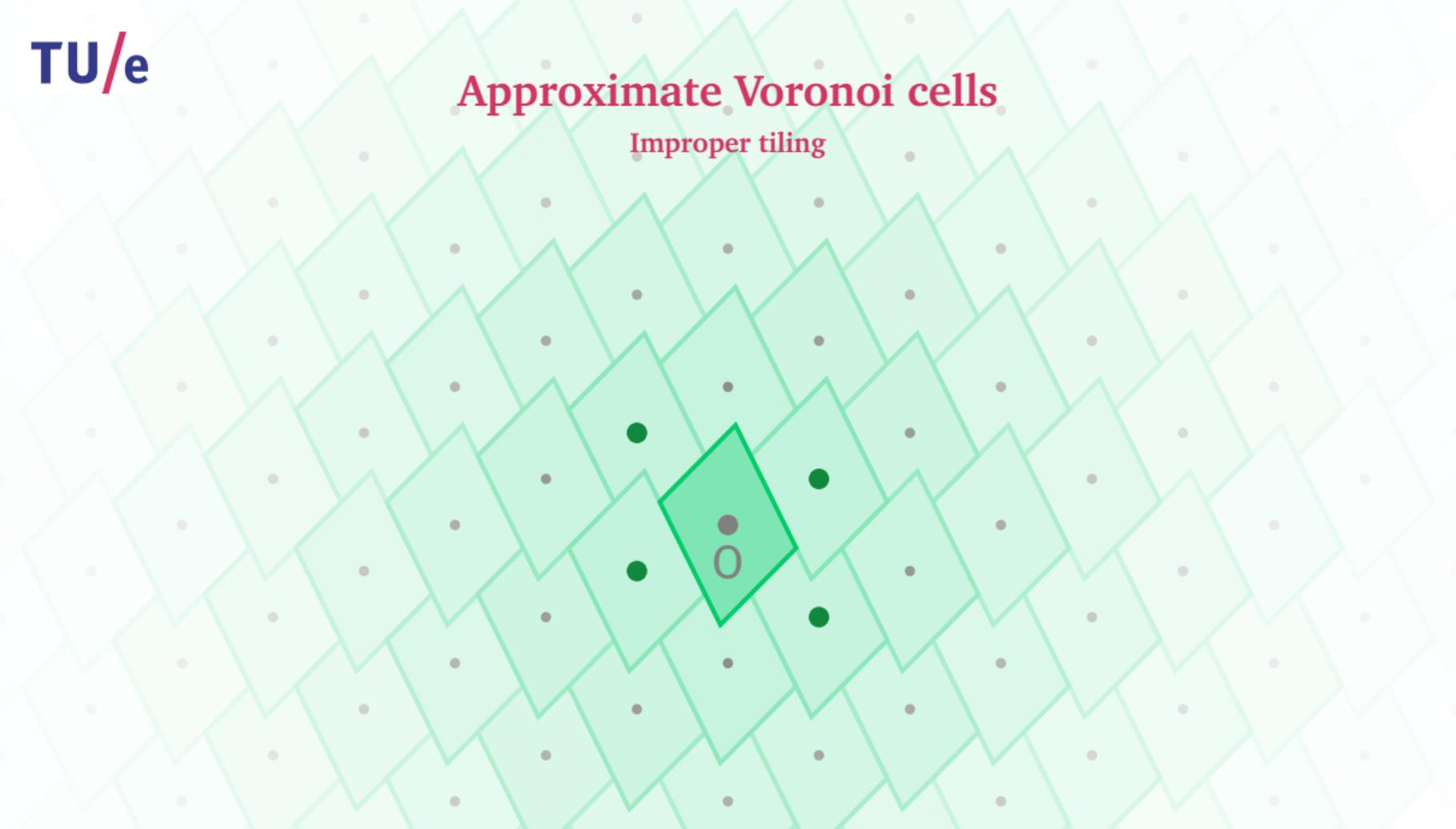
Approximate Voronoi cells

Decrease list size



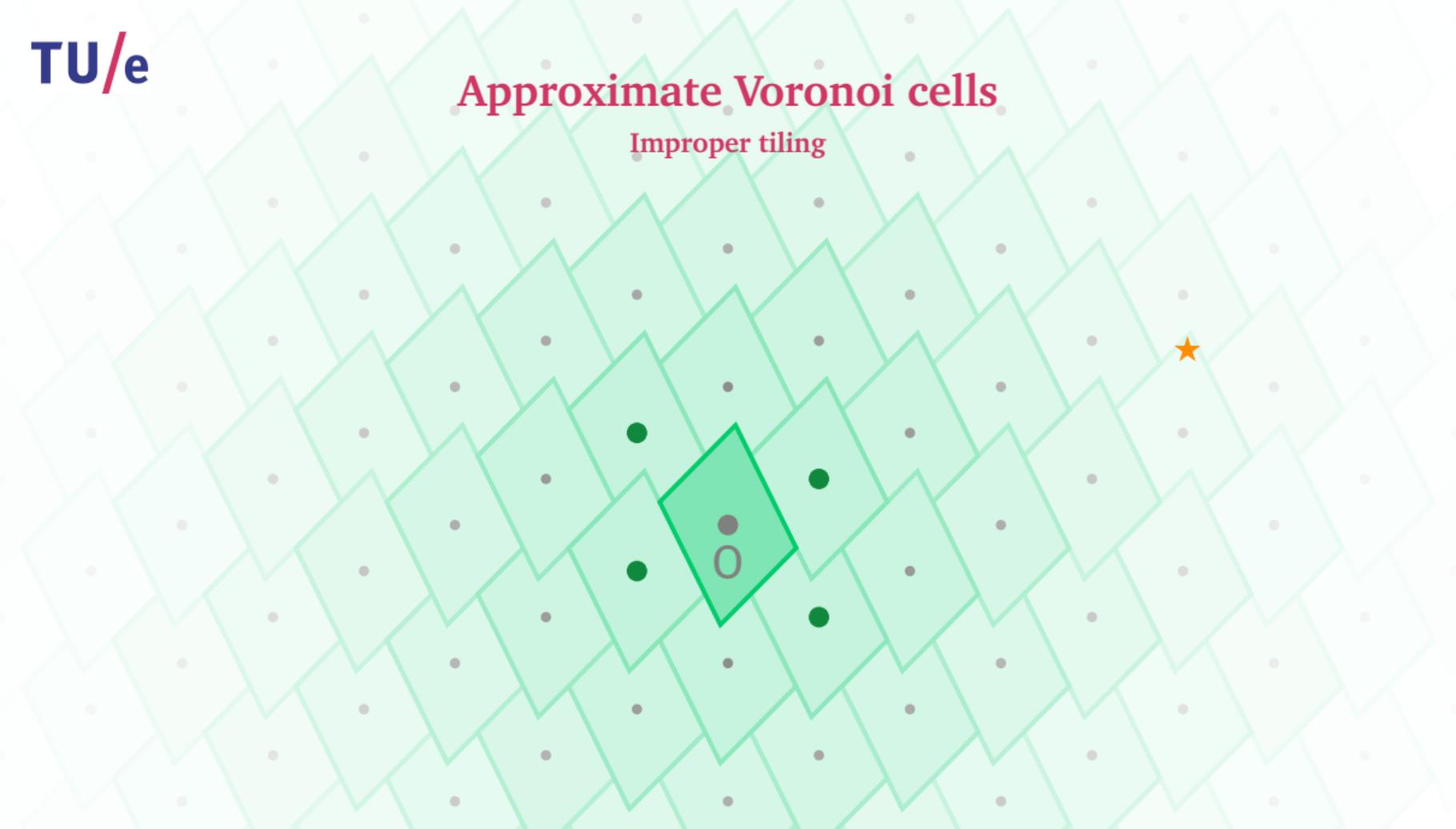
Approximate Voronoi cells

Improper tiling



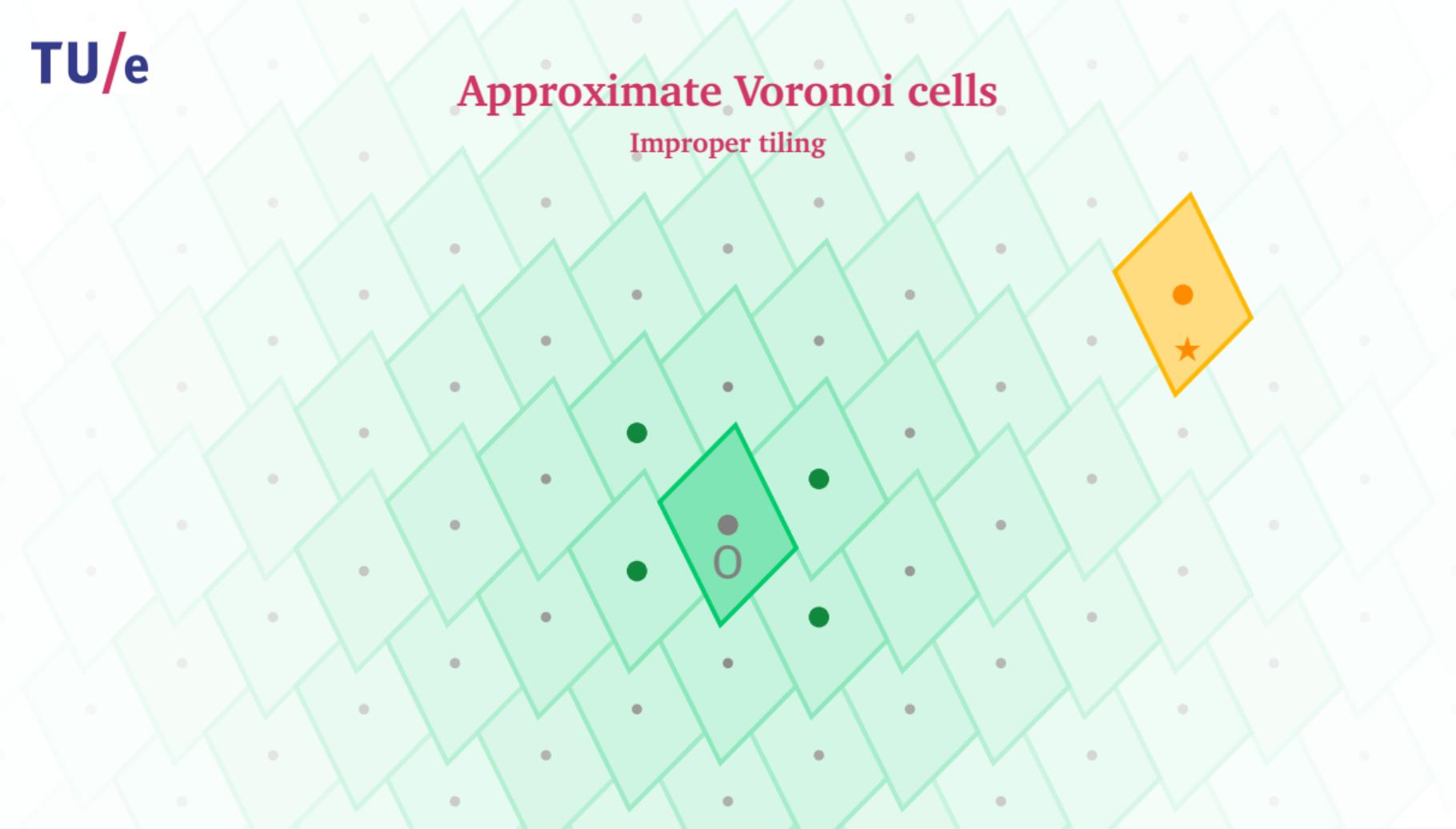
Approximate Voronoi cells

Improper tiling



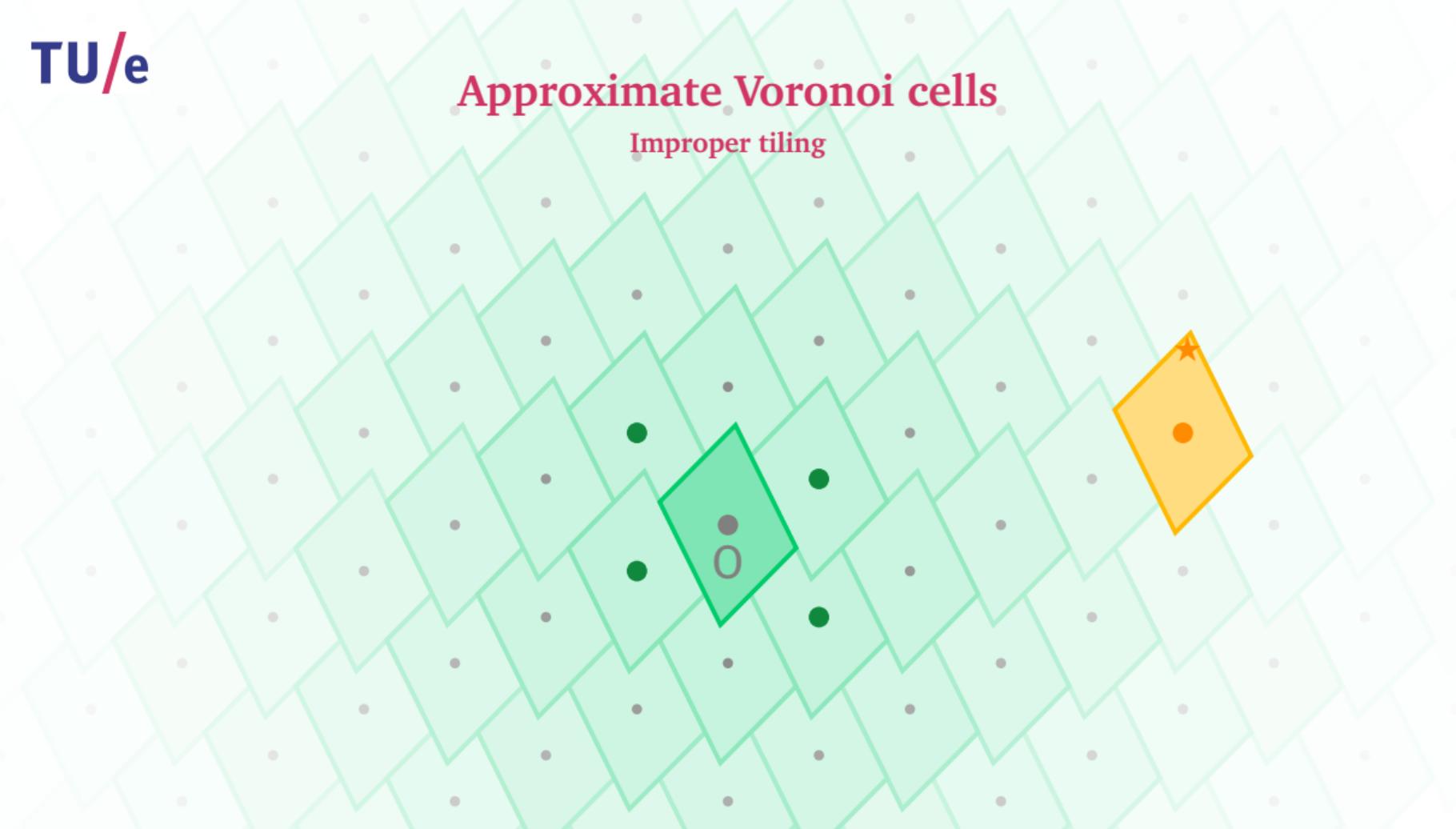
Approximate Voronoi cells

Improper tiling



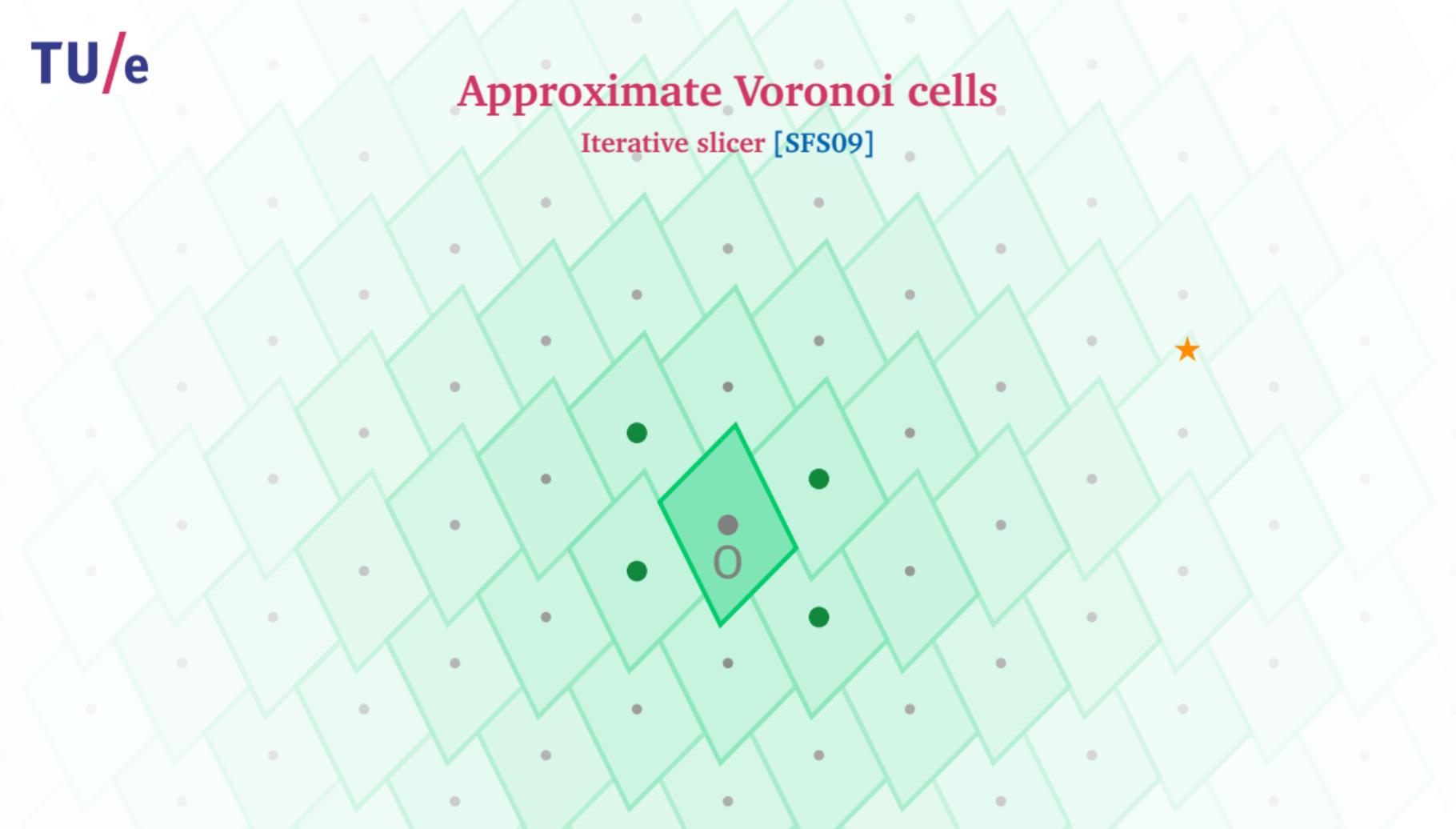
Approximate Voronoi cells

Improper tiling



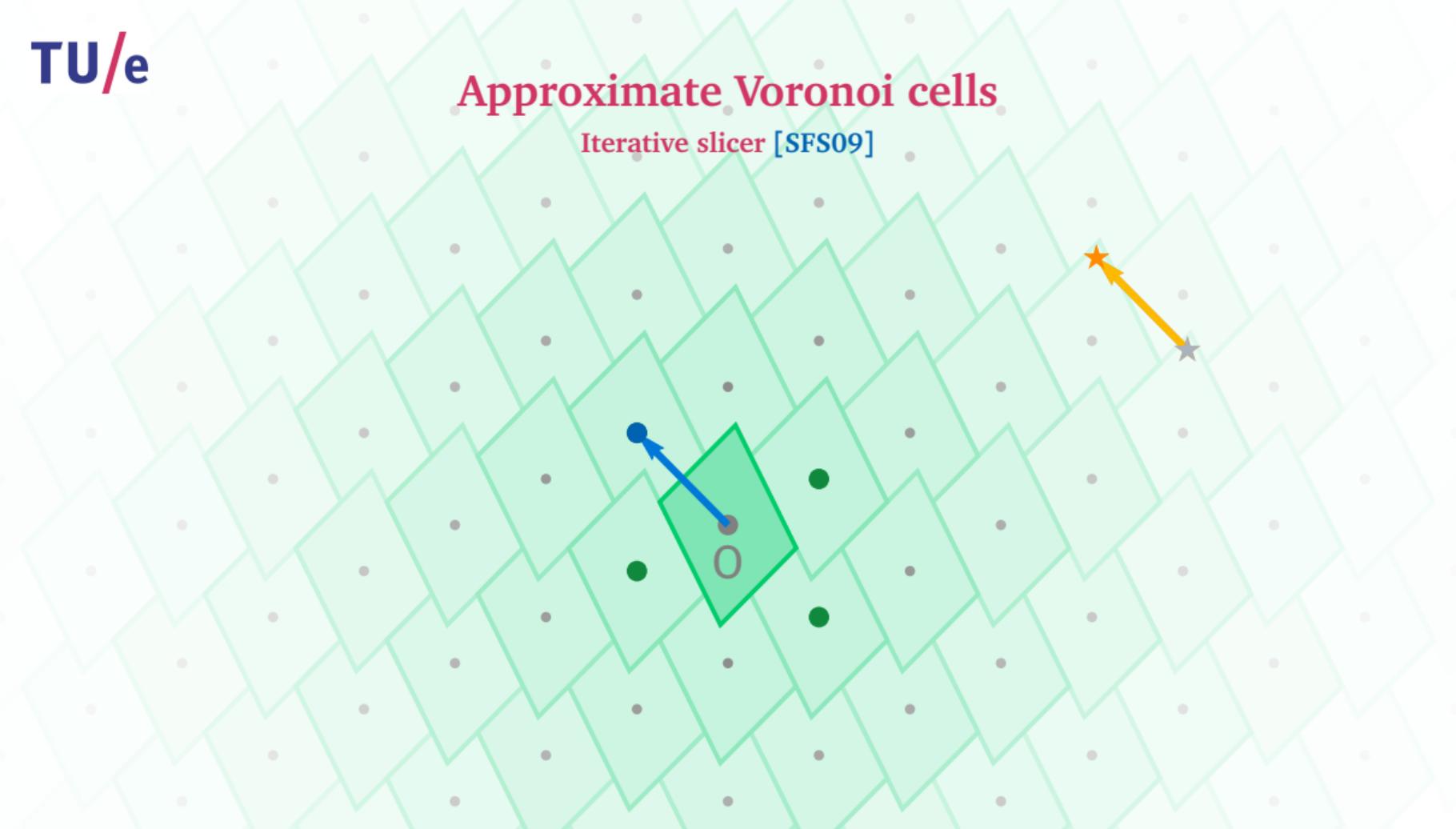
Approximate Voronoi cells

Iterative slicer [SFS09]



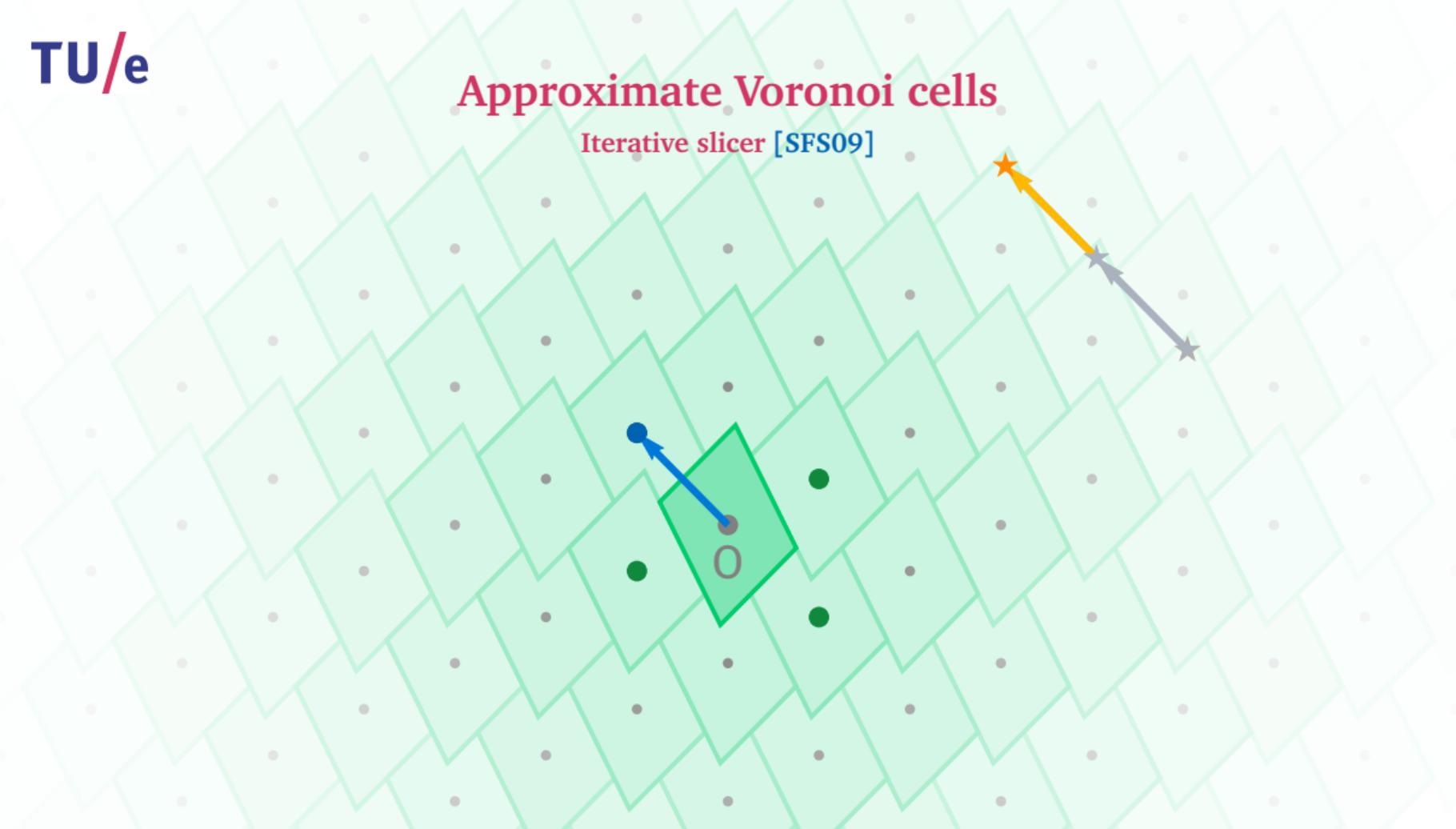
Approximate Voronoi cells

Iterative slicer [SFS09]



Approximate Voronoi cells

Iterative slicer [SFS09]



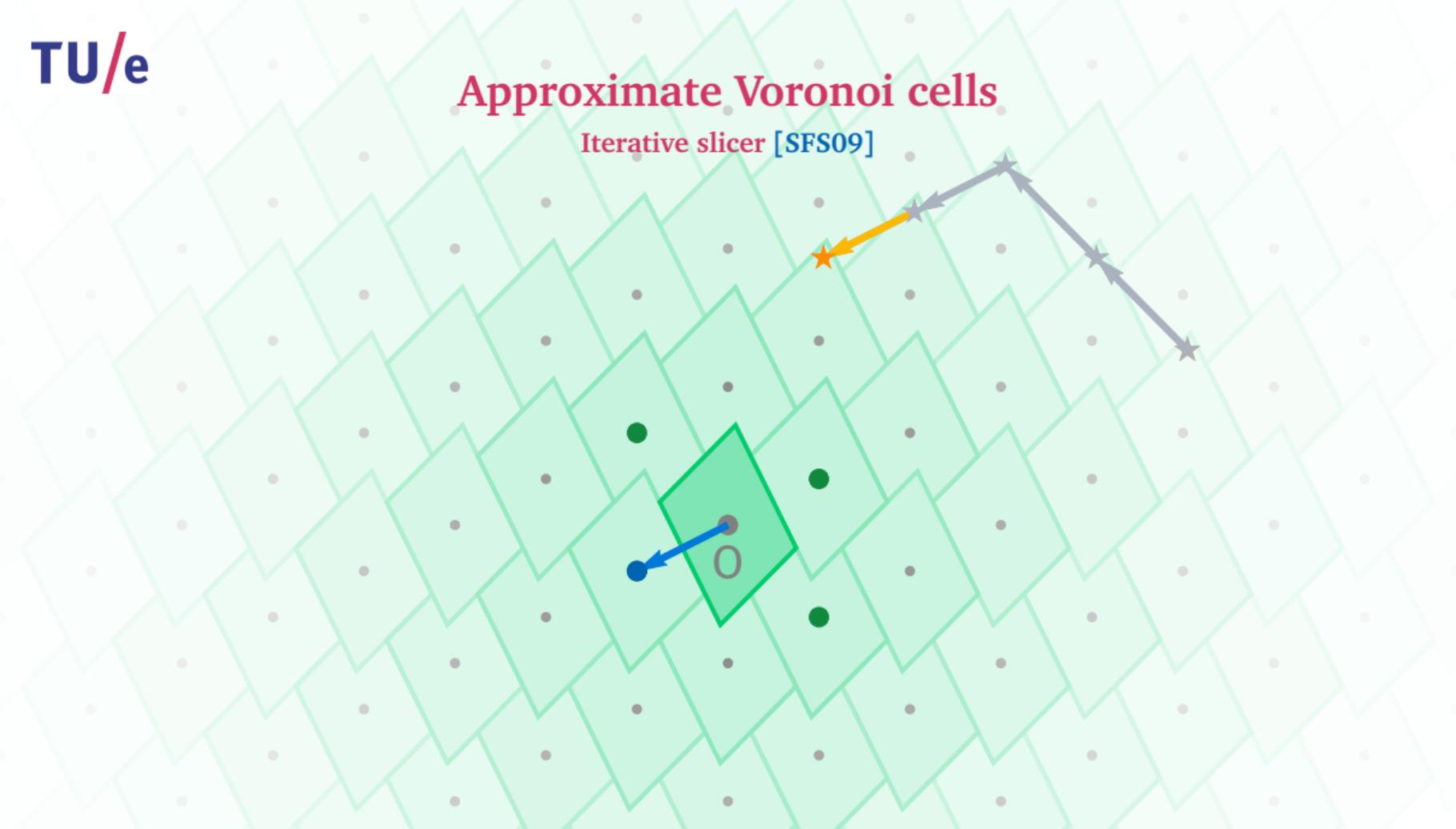
Approximate Voronoi cells

Iterative slicer [SFS09]



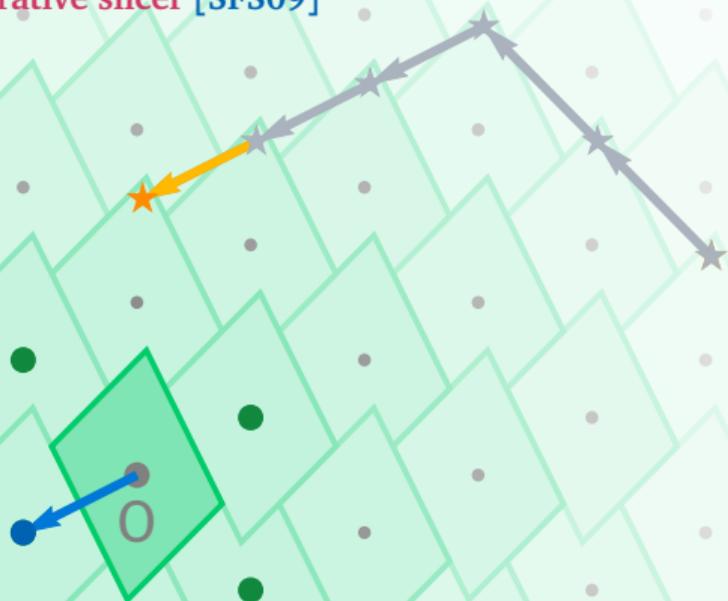
Approximate Voronoi cells

Iterative slicer [SFS09]



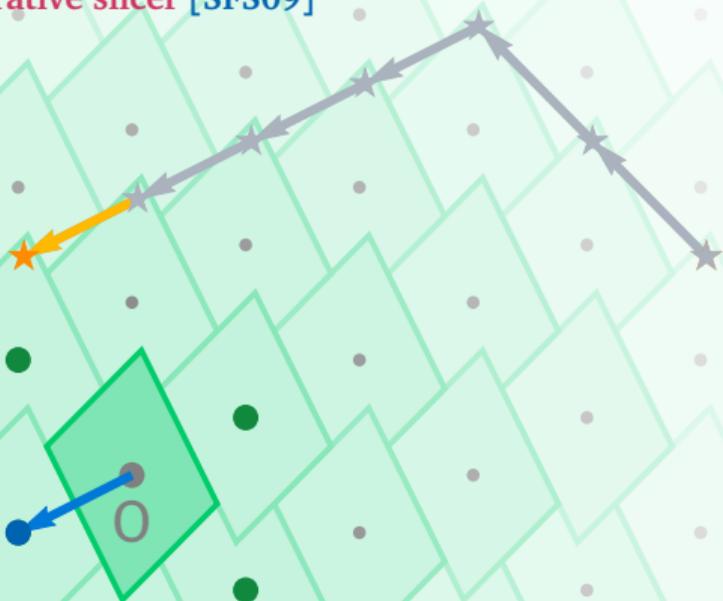
Approximate Voronoi cells

Iterative slicer [SFS09]



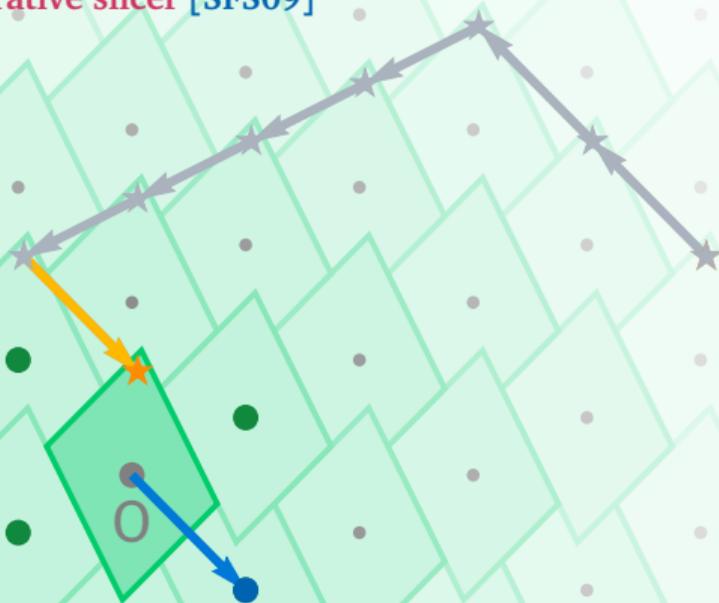
Approximate Voronoi cells

Iterative slicer [SFS09]



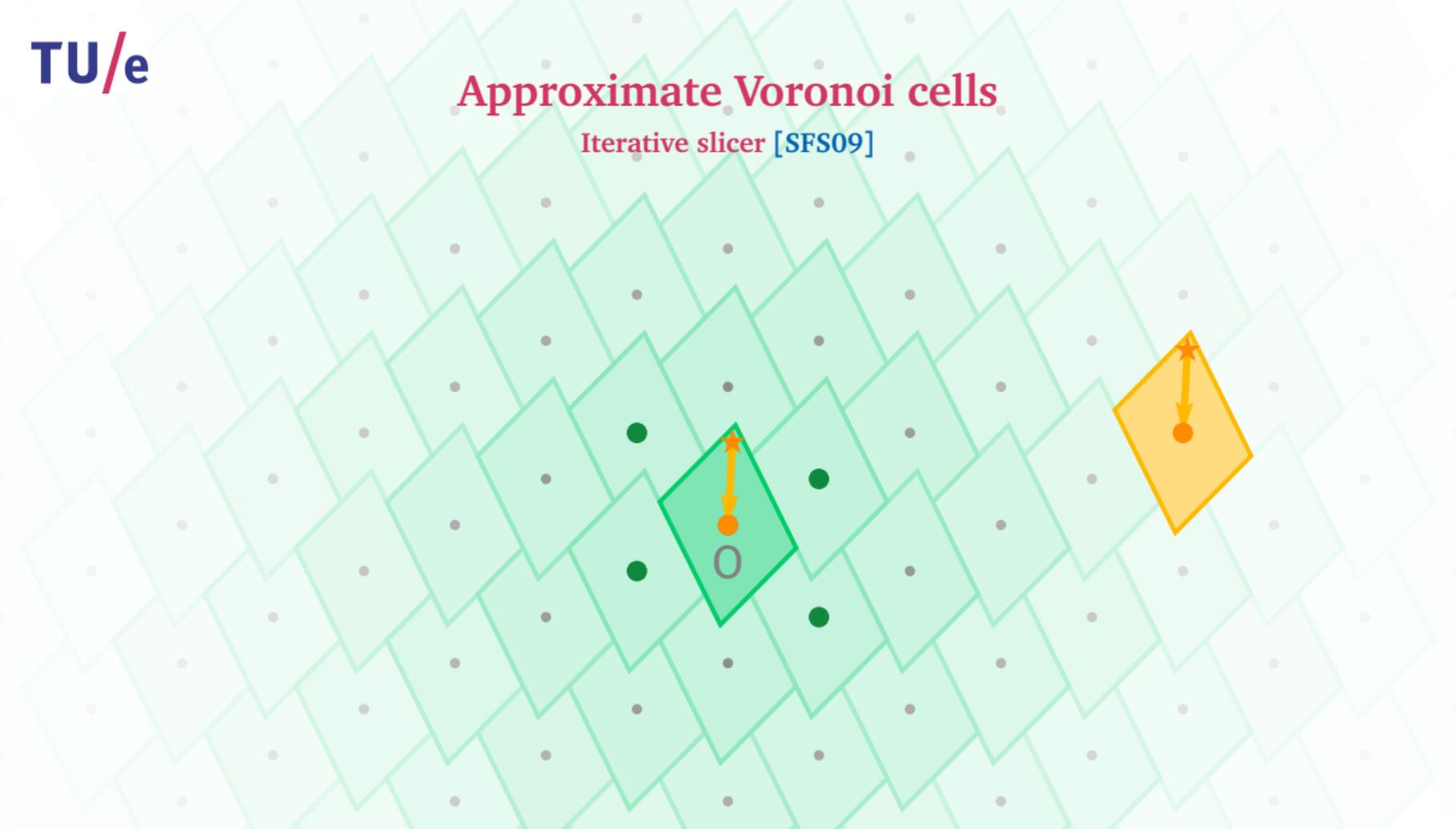
Approximate Voronoi cells

Iterative slicer [SFS09]



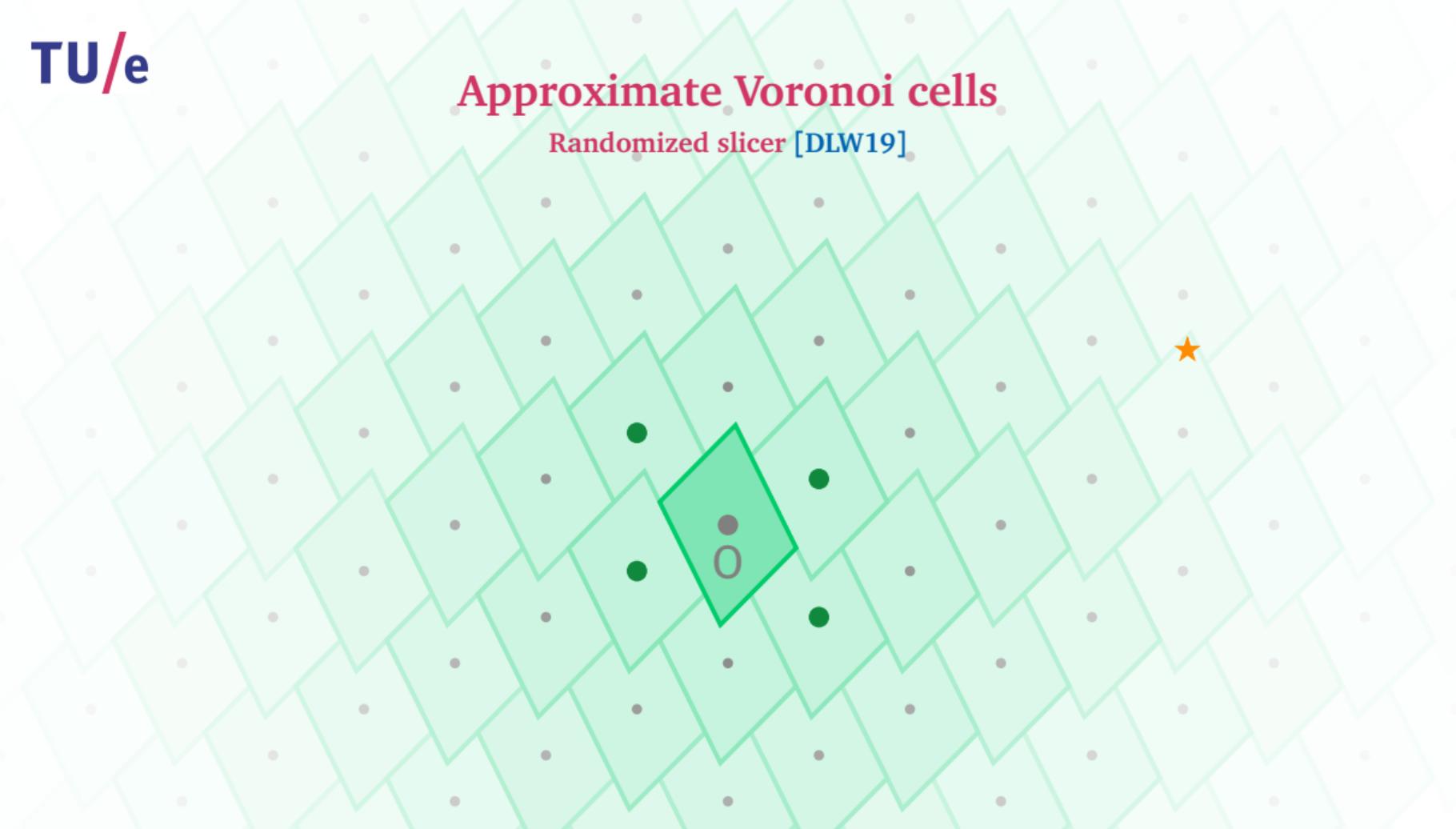
Approximate Voronoi cells

Iterative slicer [SFS09]



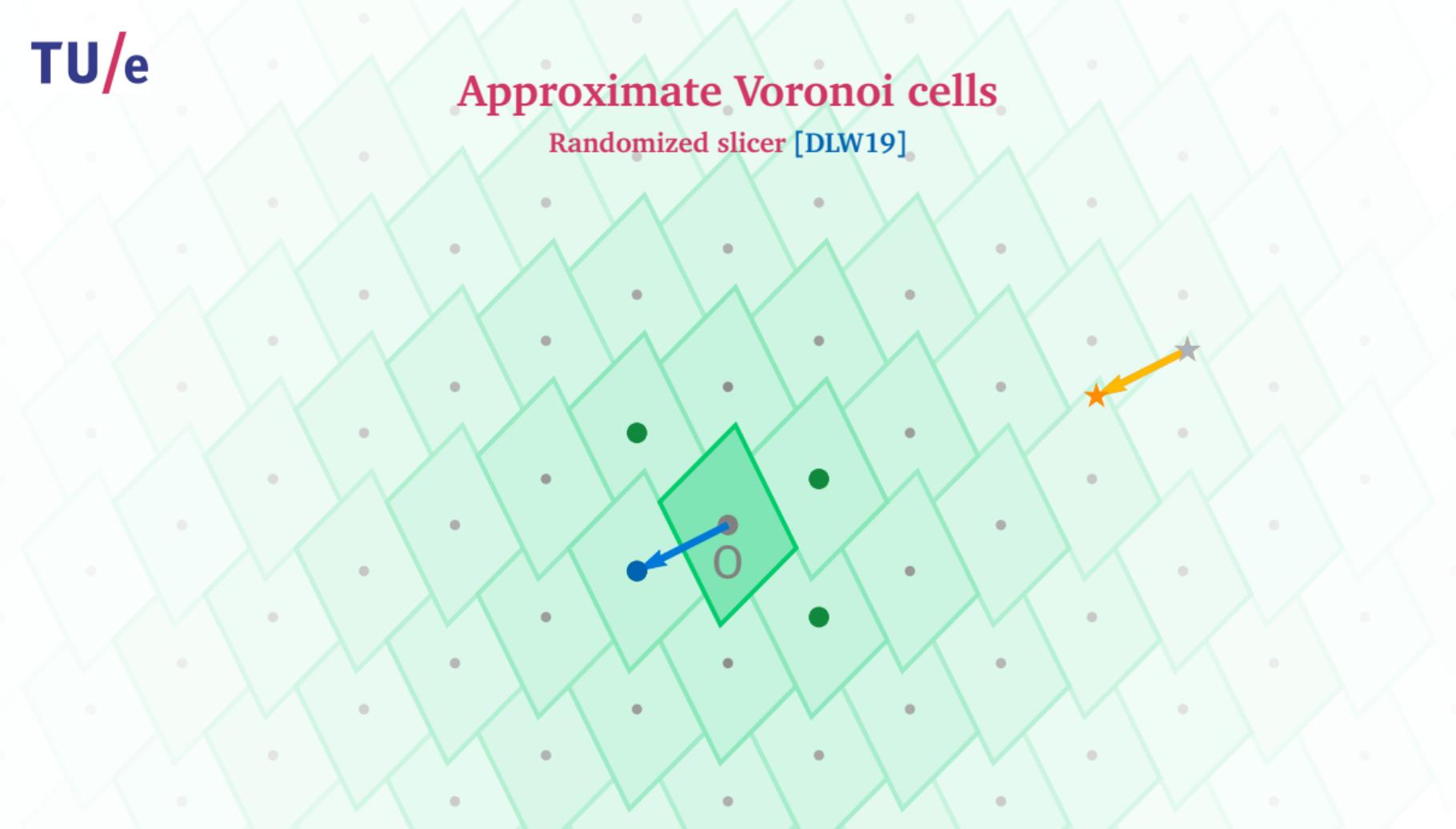
Approximate Voronoi cells

Randomized slicer [DLW19]



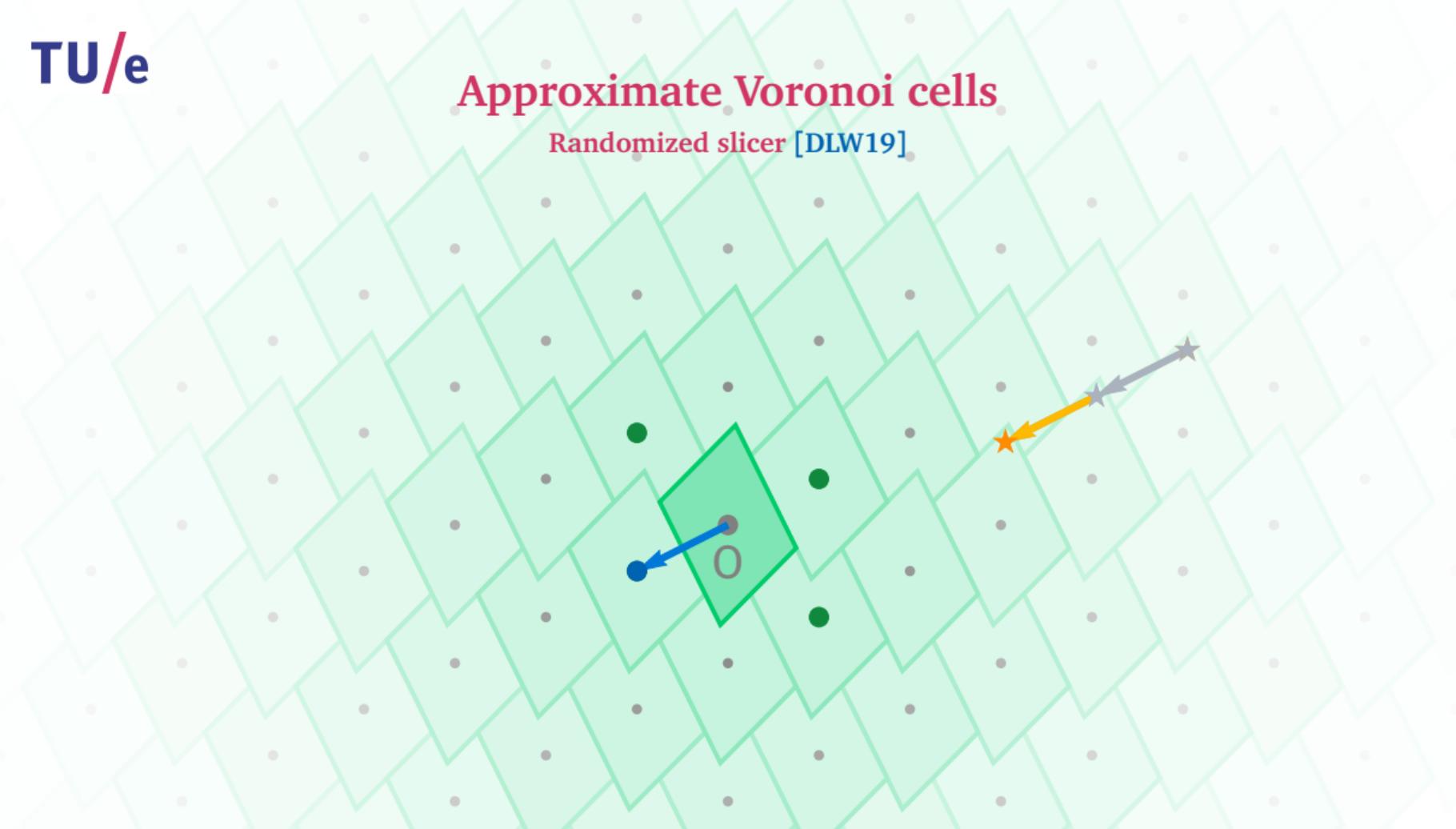
Approximate Voronoi cells

Randomized slicer [DLW19]



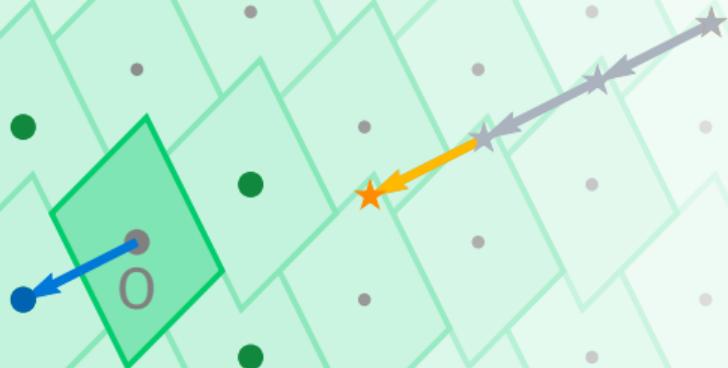
Approximate Voronoi cells

Randomized slicer [DLW19]



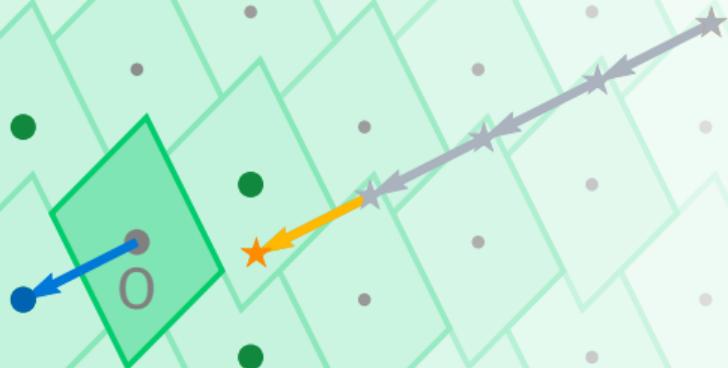
Approximate Voronoi cells

Randomized slicer [DLW19]



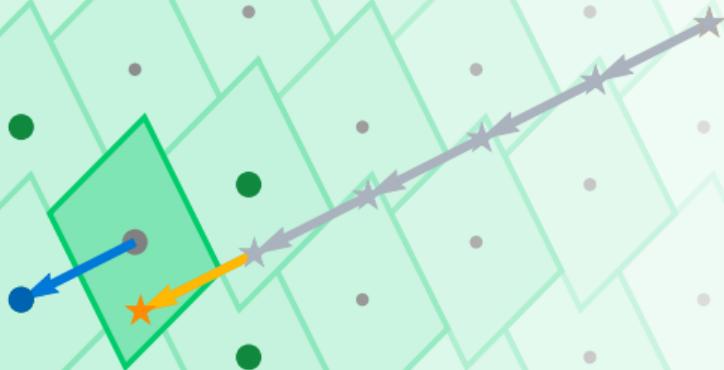
Approximate Voronoi cells

Randomized slicer [DLW19]



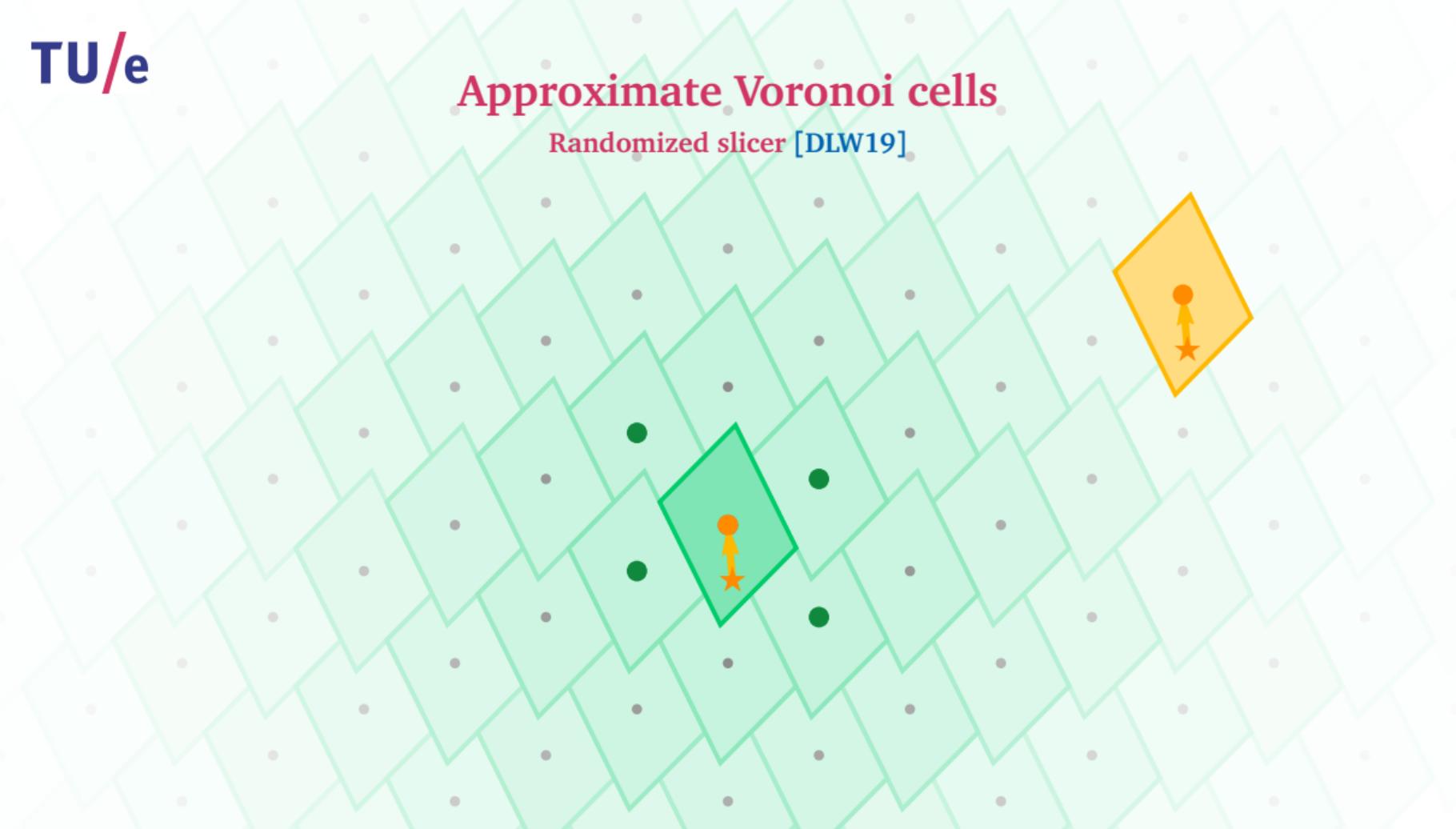
Approximate Voronoi cells

Randomized slicer [DLW19]



Approximate Voronoi cells

Randomized slicer [DLW19]



Approximate Voronoi cells

Success probability estimation

Main problem: Success probability p of the iterative slicer?

Approximate Voronoi cells

Success probability estimation

Main problem: Success probability p of the iterative slicer?

Previous results: [DLW19]

- Directly obtained a lower bound on p via the slicer
- Conjectured that p is exactly proportional to $\text{vol}(\mathcal{V})/\text{vol}(\mathcal{V}_L)$
- Open problem: obtain a tight analysis, perhaps via $\text{vol}(\mathcal{V})/\text{vol}(\mathcal{V}_L)$

Approximate Voronoi cells

Success probability estimation

Main problem: Success probability p of the iterative slicer?

Previous results: [DLW19]

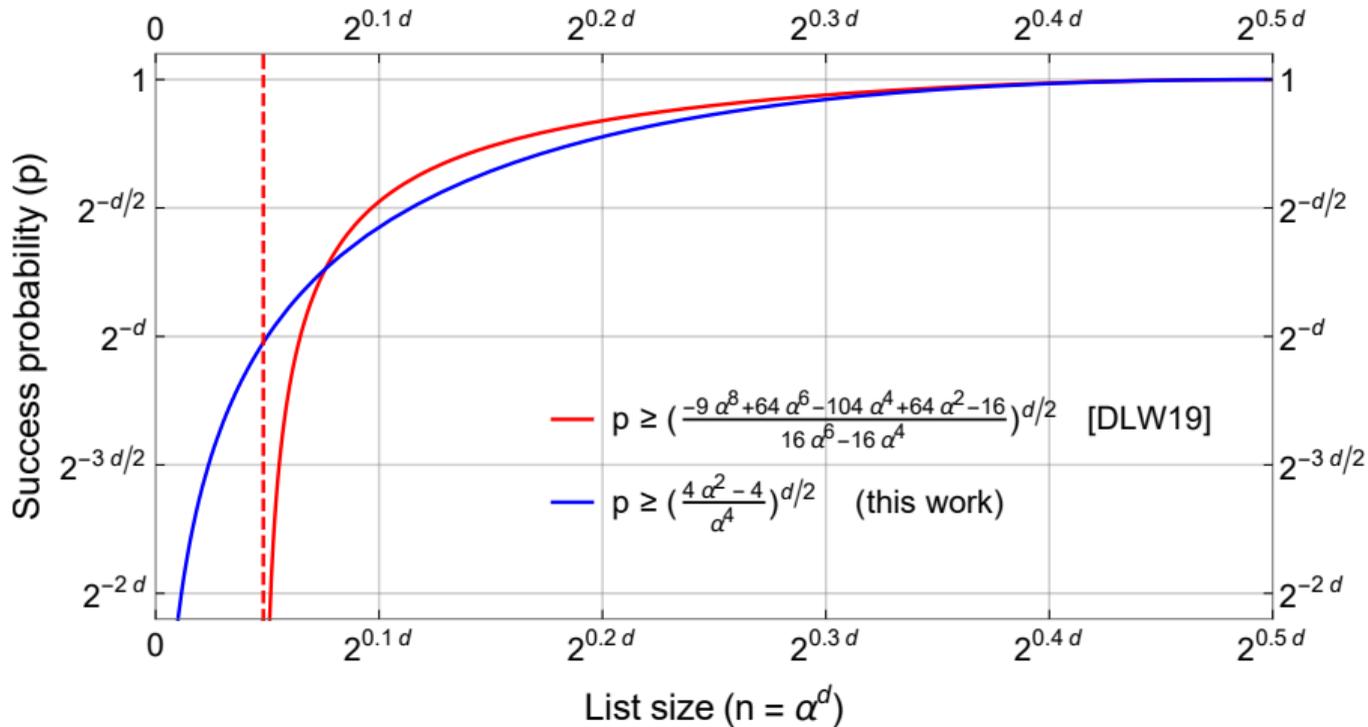
- Directly obtained a lower bound on p via the slicer
- Conjectured that p is exactly proportional to $\text{vol}(\mathcal{V})/\text{vol}(\mathcal{V}_L)$
- Open problem: obtain a tight analysis, perhaps via $\text{vol}(\mathcal{V})/\text{vol}(\mathcal{V}_L)$

This work:

- Proved tight bounds on $\text{vol}(\mathcal{V})/\text{vol}(\mathcal{V}_L)$ under the Gaussian heuristic
- Results show that p **cannot** be (exactly) proportional to $\text{vol}(\mathcal{V})/\text{vol}(\mathcal{V}_L)$
- From $p \geq \text{vol}(\mathcal{V})/\text{vol}(\mathcal{V}_L)$ we obtain new lower bounds on p
- No nonsensical asymptote at $2^{0.05d}$ memory anymore

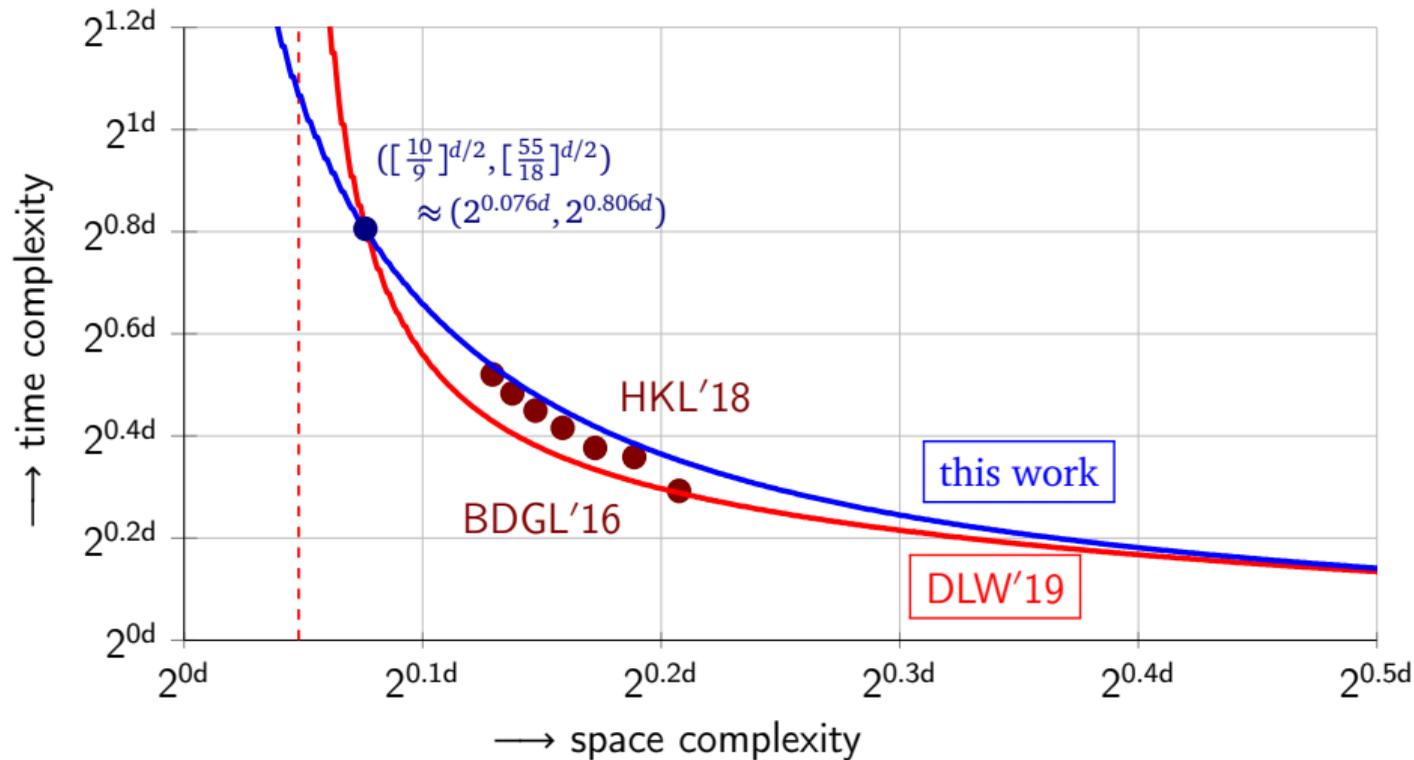
Approximate Voronoi cells

Lower bounds on success probability



Approximate Voronoi cells

Time-space trade-offs for CVPP



Conclusion

Voronoi cells

- Solves CVPP exactly in the worst case for all lattices
- Requires too much space (and time) to be useful

Voronoi cells

- Solves CVPP exactly in the worst case for all lattices
- Requires too much space (and time) to be useful

Approximate Voronoi cells

- Offers heuristic alternative to exact Voronoi cells
- Success probability analysis:
 - ▶ Original analysis did not appear to be tight
 - ▶ Conjectured that tighter bounds may be obtained via $\text{vol}(\mathcal{V})/\text{vol}(\mathcal{V}_L)$
 - ▶ This work: obtained tight bounds on the ratio $\text{vol}(\mathcal{V})/\text{vol}(\mathcal{V}_L)$
 - ▶ Results in better CVPP complexities for low-memory regime
 - ▶ Unfortunately, approach via $\text{vol}(\mathcal{V})/\text{vol}(\mathcal{V}_L)$ is not tight either

Voronoi cells

- Solves CVPP exactly in the worst case for all lattices
- Requires too much space (and time) to be useful

Approximate Voronoi cells

- Offers heuristic alternative to exact Voronoi cells
- Success probability analysis:
 - ▶ Original analysis did not appear to be tight
 - ▶ Conjectured that tighter bounds may be obtained via $\text{vol}(\mathcal{V})/\text{vol}(\mathcal{V}_L)$
 - ▶ This work: obtained tight bounds on the ratio $\text{vol}(\mathcal{V})/\text{vol}(\mathcal{V}_L)$
 - ▶ Results in better CVPP complexities for low-memory regime
 - ▶ Unfortunately, approach via $\text{vol}(\mathcal{V})/\text{vol}(\mathcal{V}_L)$ is not tight either

Open problems

- Obtain truly tight bounds (ongoing work with Leo Ducas, Wessel van Woerden)
- Find an efficient BDDP-version of this CVPP algorithm