

Approximate Voronoi cells for lattices, revisited

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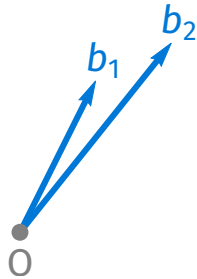
Lattices

Basics

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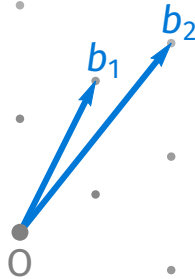
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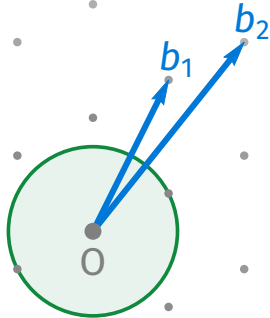
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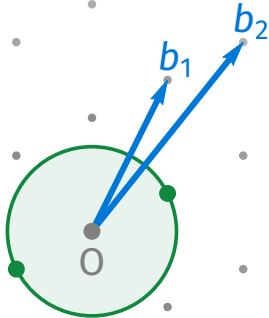
Lattice problems

Shortest Vector Problem (SVP)



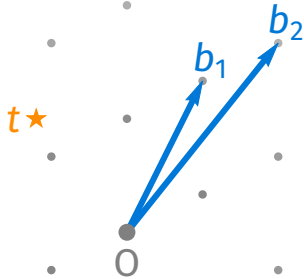
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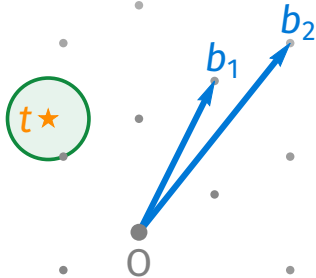
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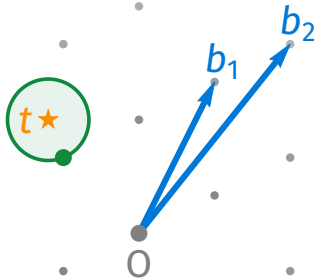
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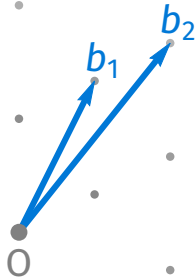
Lattice problems

Asymptotics for SVP and CVP

	Algorithm	$\log_2(\text{Time})$	$\log_2(\text{Space})$	Experiments
Worst-case SVP	Enumeration [Poh81, Kan83, FP85, ...]	$O(n \log n)$	$O(\log n)$	–
	AKS-sieve [AKS01, NV08, MV10, HPS11]	$3.398n$	$1.985n$	–
	Birthday sieves [PS09, HPS11]	$2.465n$	$1.233n$	–
	Enumeration/DGS hybrid [CCL17]	$2.048n$	$0.500n$	–
	Voronoi cell algorithm [AEVZ02, MV10b, BD15]	$2.000n$	$1.000n$	40
	Quantum sieve [LMP13, LMP15]	$1.799n$	$1.286n$	–
	Quantum enumeration/DGS hybrid [CCL17]	$1.256n$	0.500n	–
	Discrete Gaussian sampling [ADRS15, ADS15, AS18]	1.000n	$1.000n$	–
Average-case SVP	Enumeration [GNR10, MW15, AN17]	$O(n \log n)$	$O(\log n)$	152
	The Nguyen–Vidick sieve [NV08]	$0.415n$	$0.208n$	50
	GaussSieve [MV10, ..., IKMT14, BNvdP16, YKYC17]	$0.415n$	$0.208n$	130*
	Triple sieve [BLS16, HK17]	$0.396n$	$0.189n$	80
	Kleinjung sieve [Kle14]	$0.379n$	$0.189n$	116
	Leveled sieving [WLTB11, ZPH13]	$0.378n$	$0.283n$	–
	Overlattice sieve [BGJ14]	$0.377n$	$0.293n$	90
	Triple sieve with NNS [HK17, HKL18]	$0.359n$	0.189n	76
	Single filters [DL17, ADH+19]	$0.349n$	$0.246n$	155
	Hyperplane LSH [Cha02, FBB+14, Laa15, ..., LM18]	$0.337n$	$0.337n$	107
	May–Ozerov NNS [MO15, BGJ15]	$0.311n$	$0.311n$	–
	Quantum sieve [LMP13]	$0.311n$	$0.208n$	–
	Spherical LSH [AINR14, LdW15]	$0.297n$	$0.297n$	–
	Cross-polytope LSH [TT07, AILRS15, BL16, KW17]	$0.297n$	$0.297n$	80
	Spherical LSF [BDGL16, MLB17, ALRW17, DSvW19]	0.292n	$0.292n$	157
	Quantum NNS sieve [LMP15, Laa16]	0.265n	$0.265n$	–

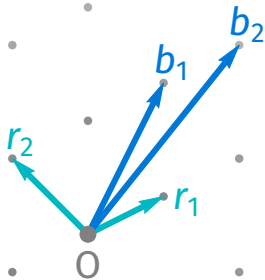
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Closest Vector Problem with Preprocessing (CVPP)



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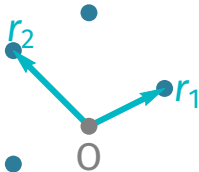
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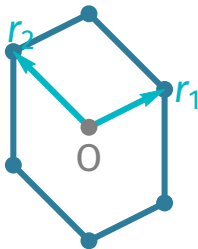
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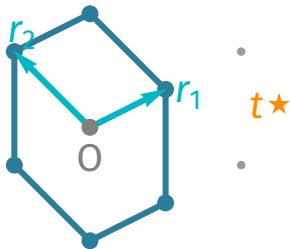
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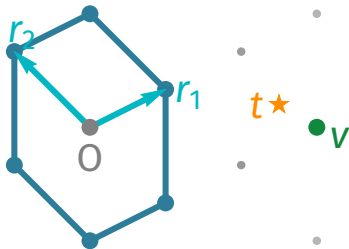
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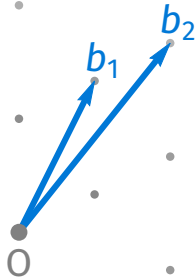
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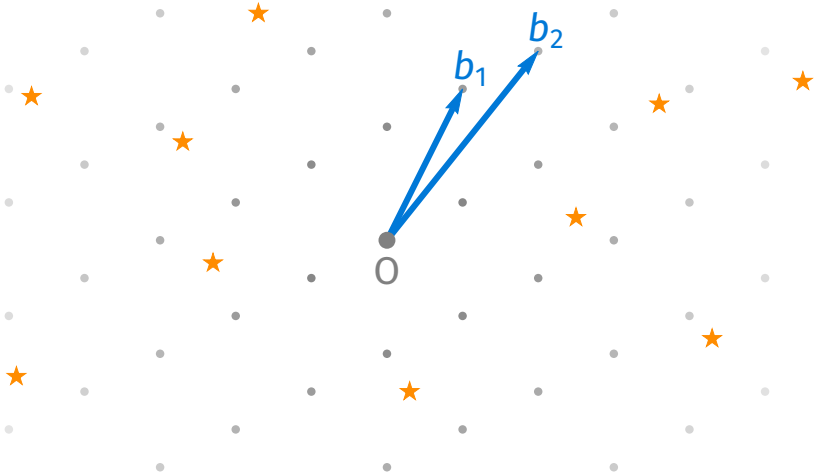
Lattice problems

Batch Closest Vector Problem



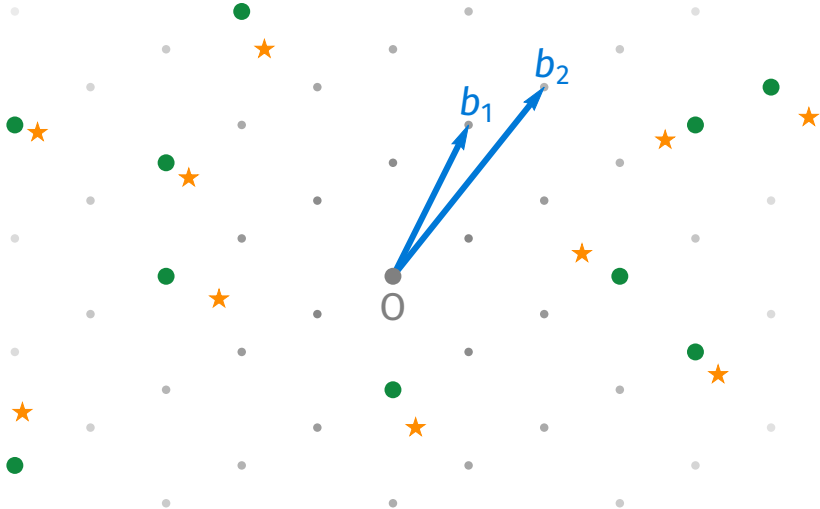
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Lattice problems

Why study CVPP?

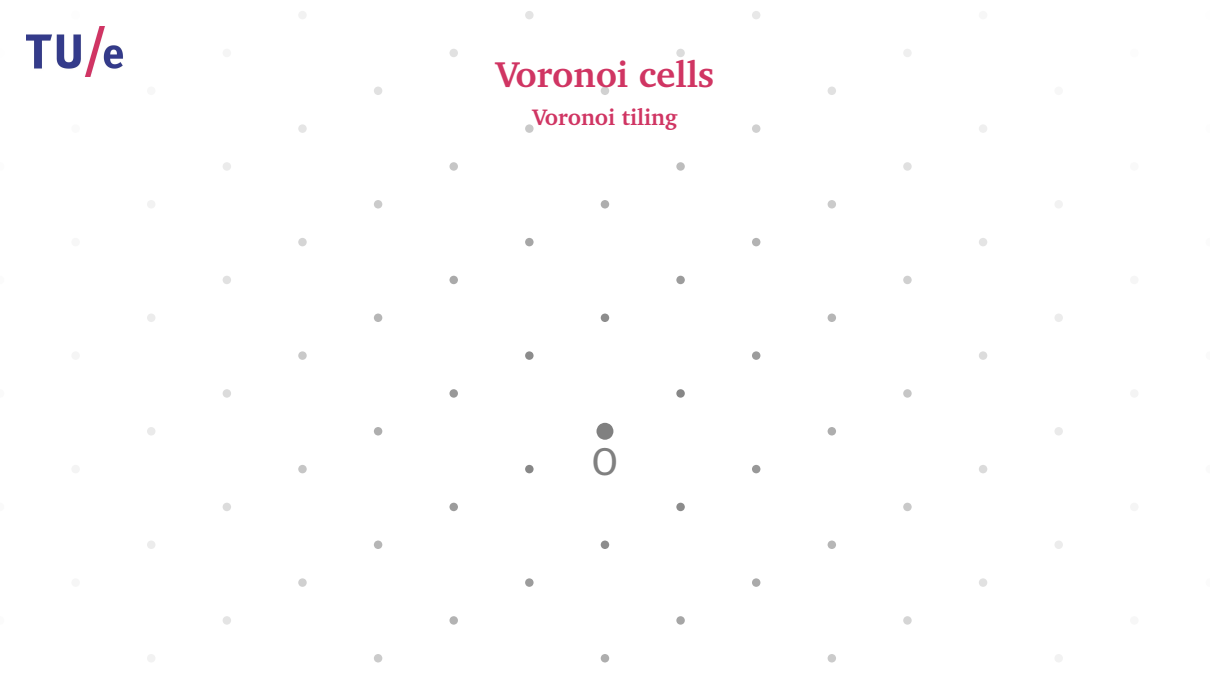
Concrete applications

- Speeding up lattice enumeration for SVP or CVP [GNR10]
- Solving approximate SVP on ideal lattices [PHS19]
- Computing class group actions in a relation lattice [BKV19]

Commonly a lattice basis (public key) is known long before the target vectors (encryptions, signatures)

Voronoi cells

Voronoi tiling

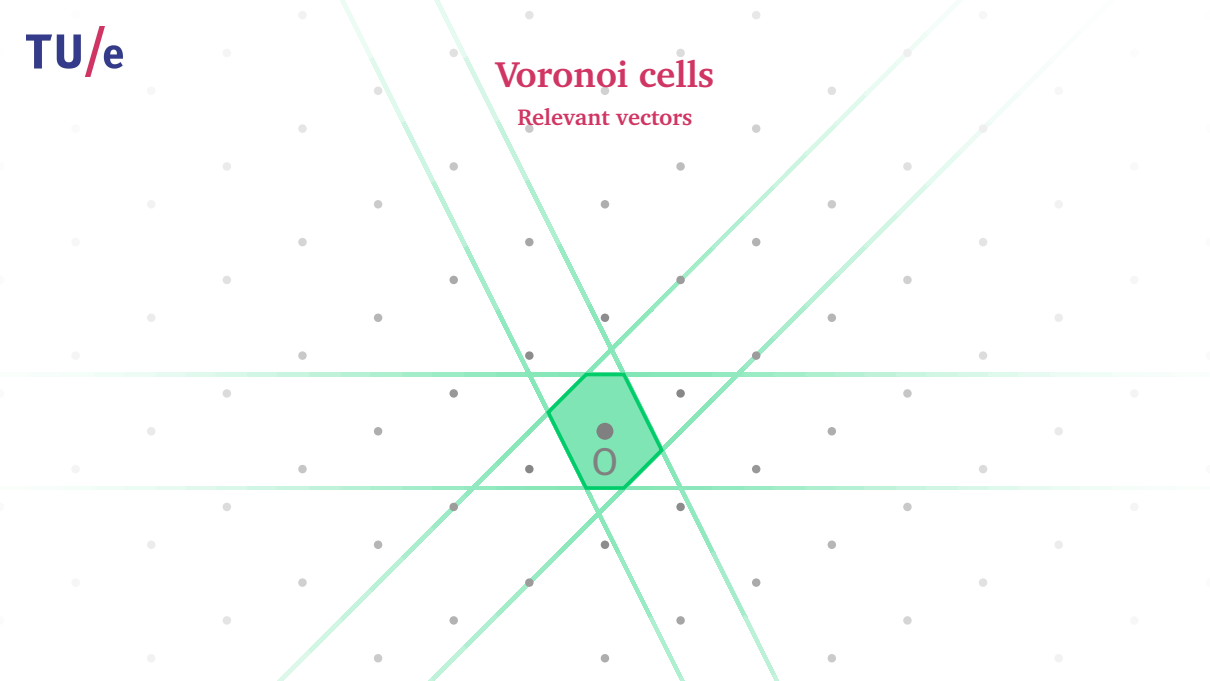


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Voronoi tiling

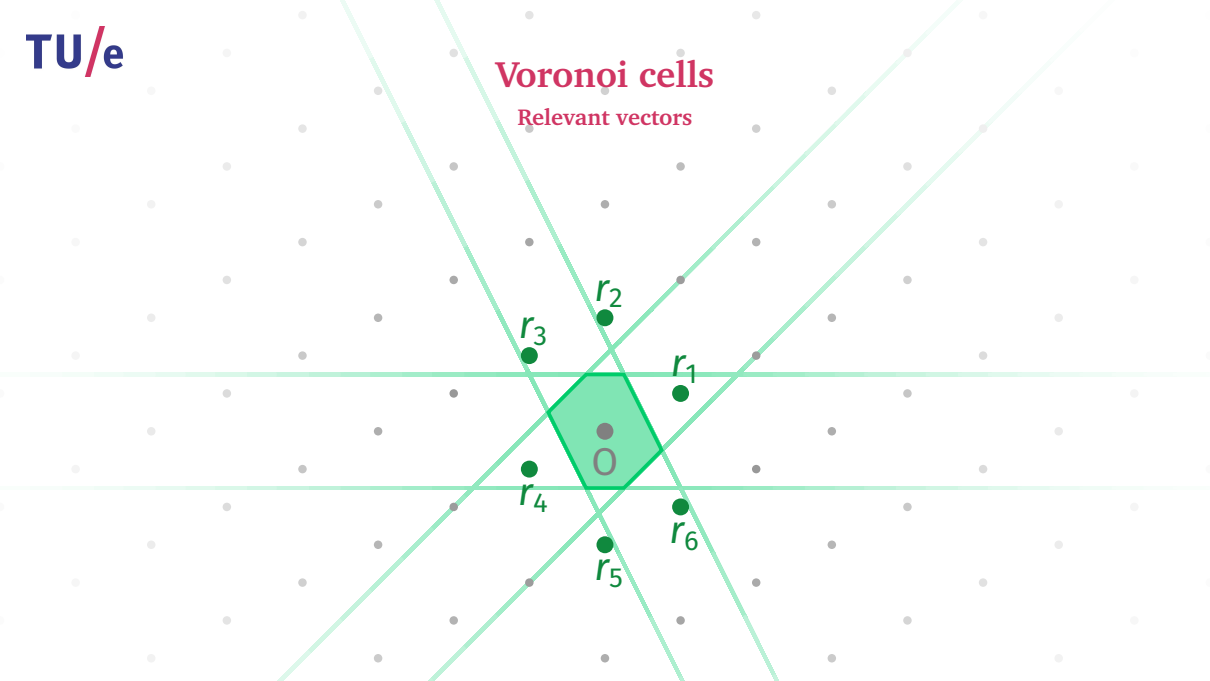


Voronoi cells
Relevant vectors



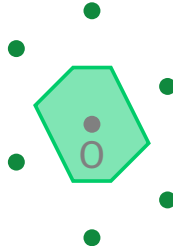
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Relevant vectors

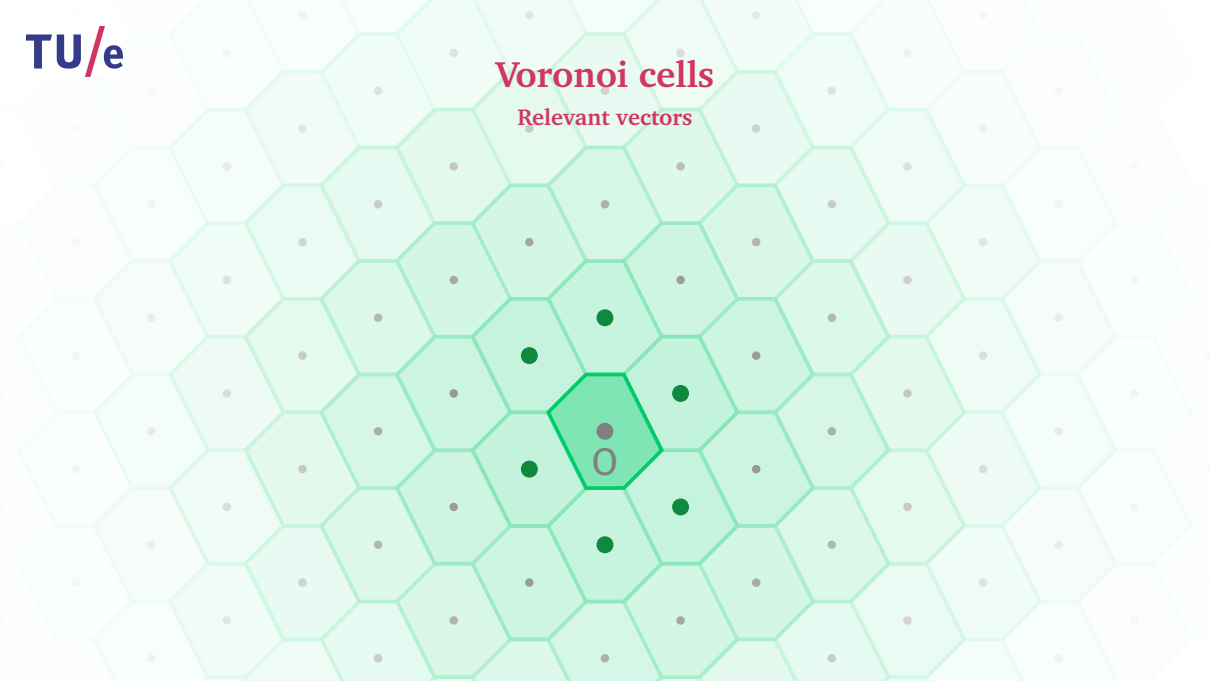


Voronoi cells

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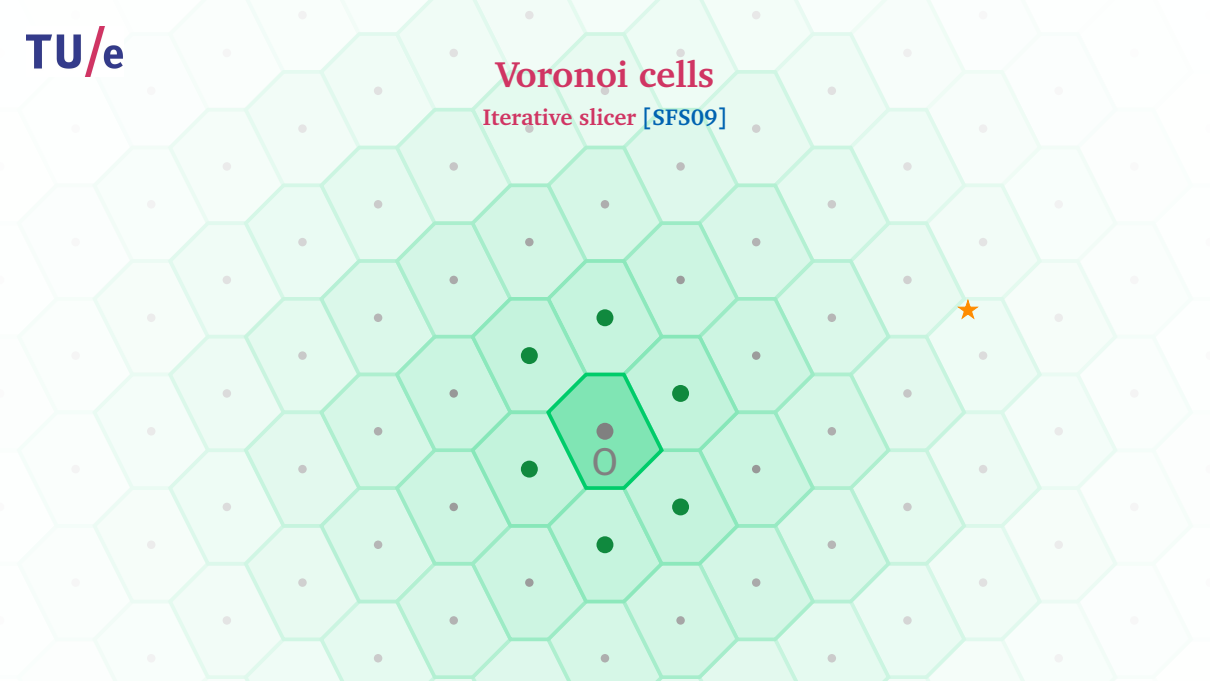


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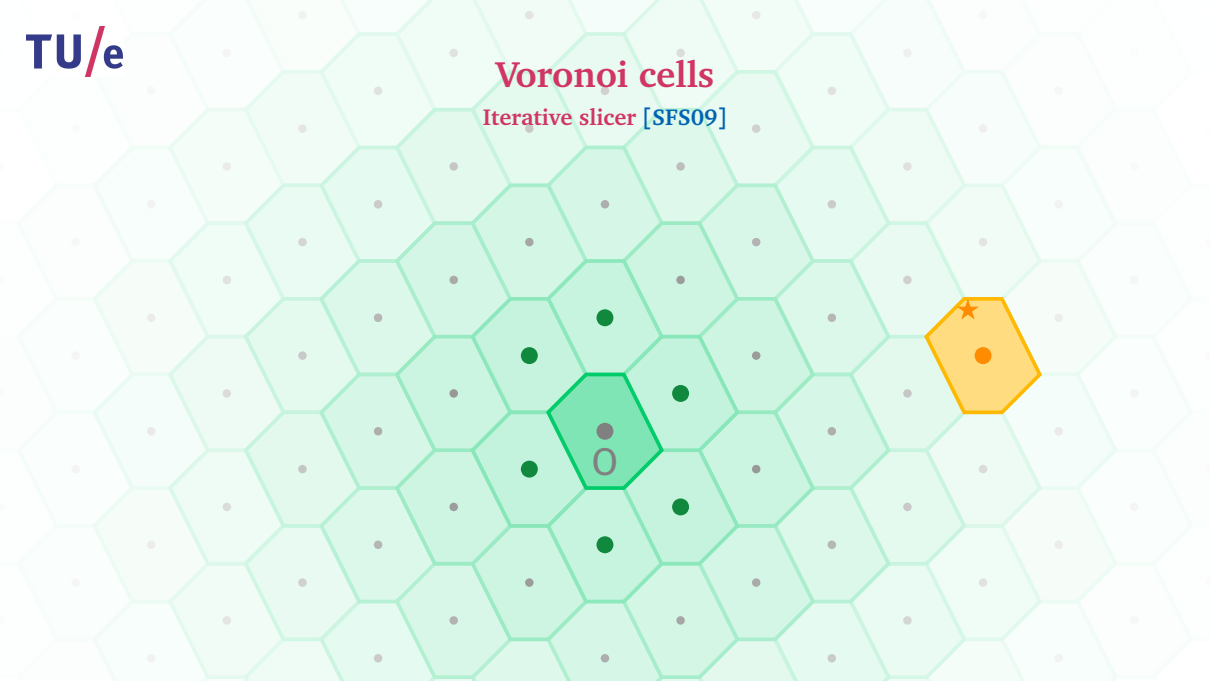
Voronoi cells

Iterative slicer [SFS09]



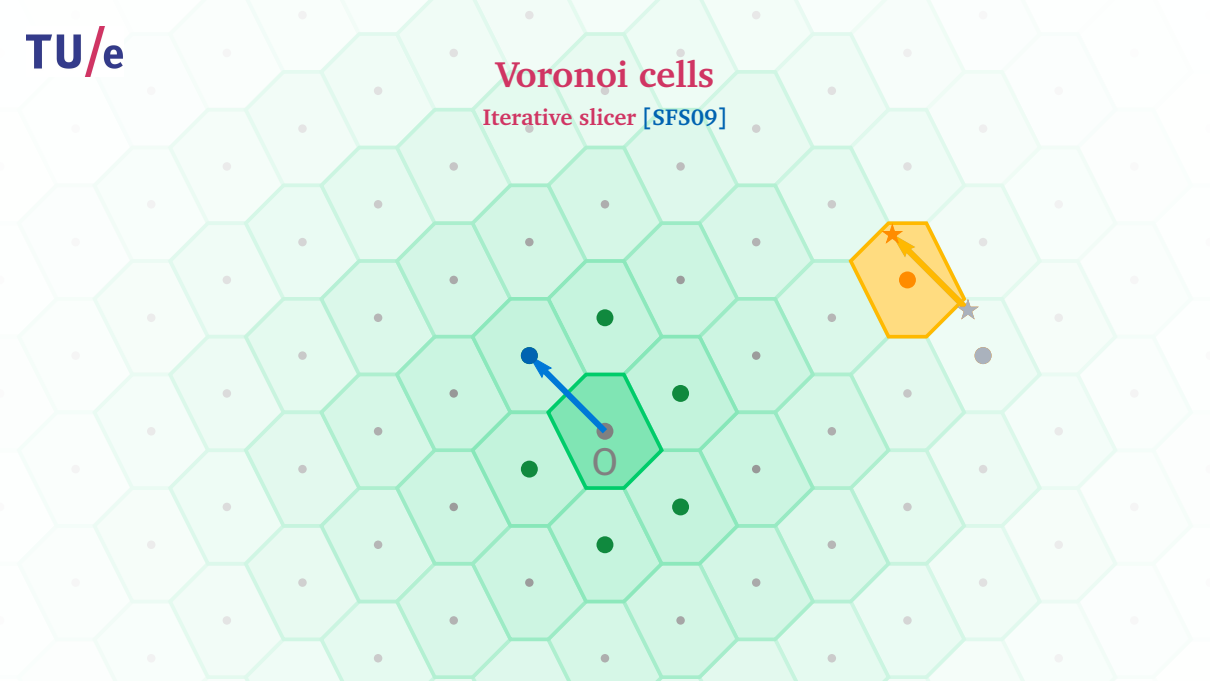
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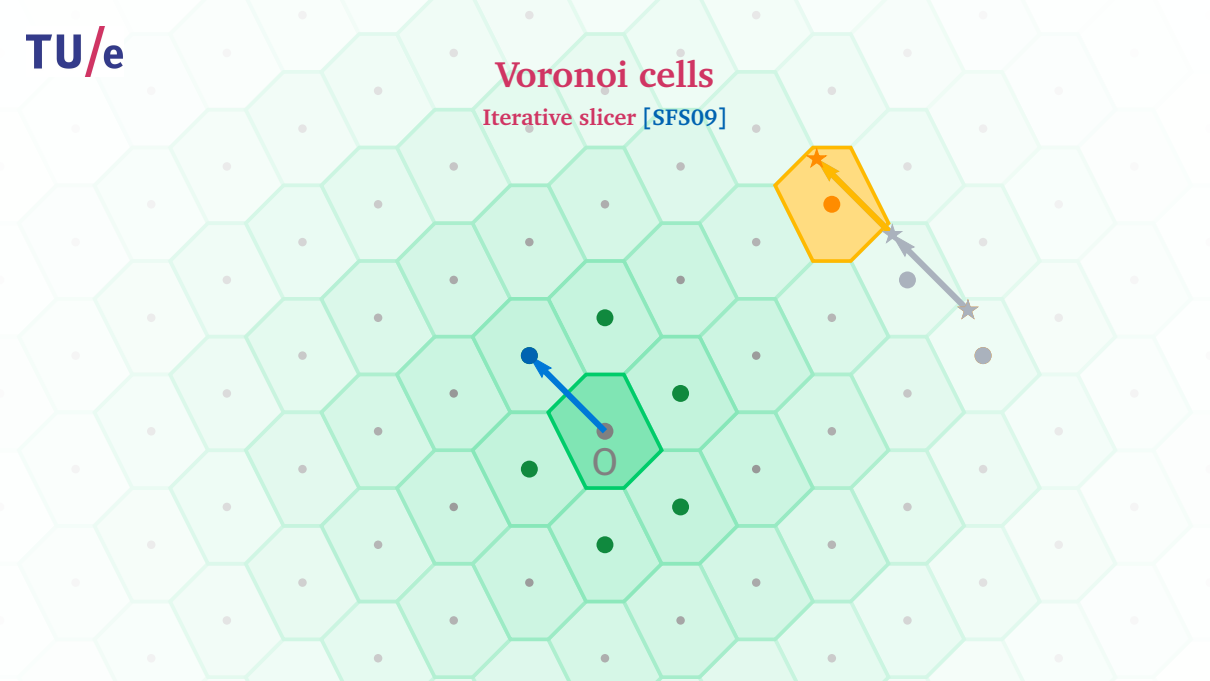
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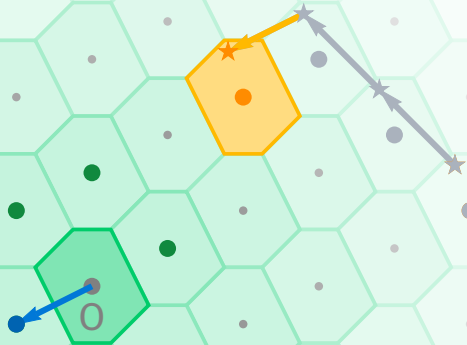
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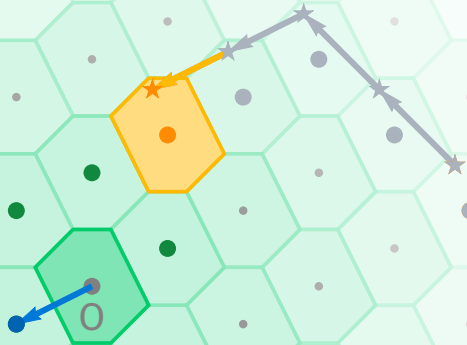
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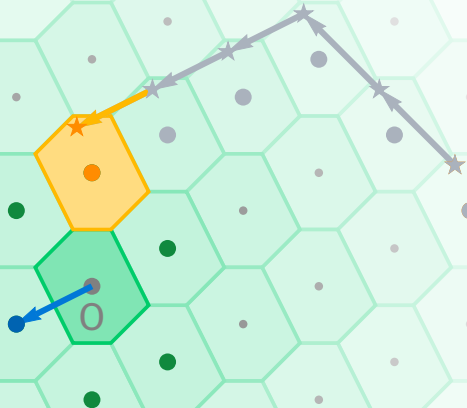
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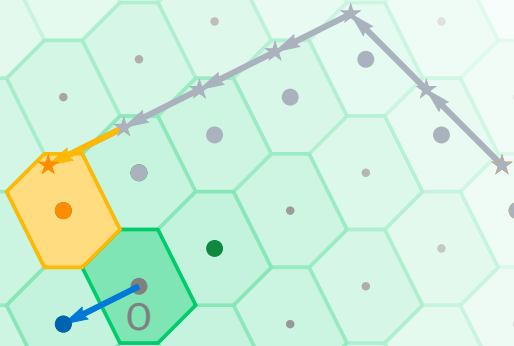
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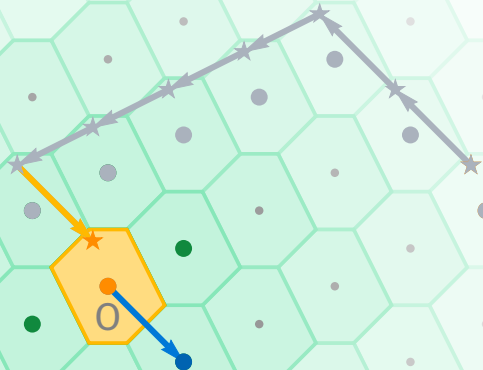
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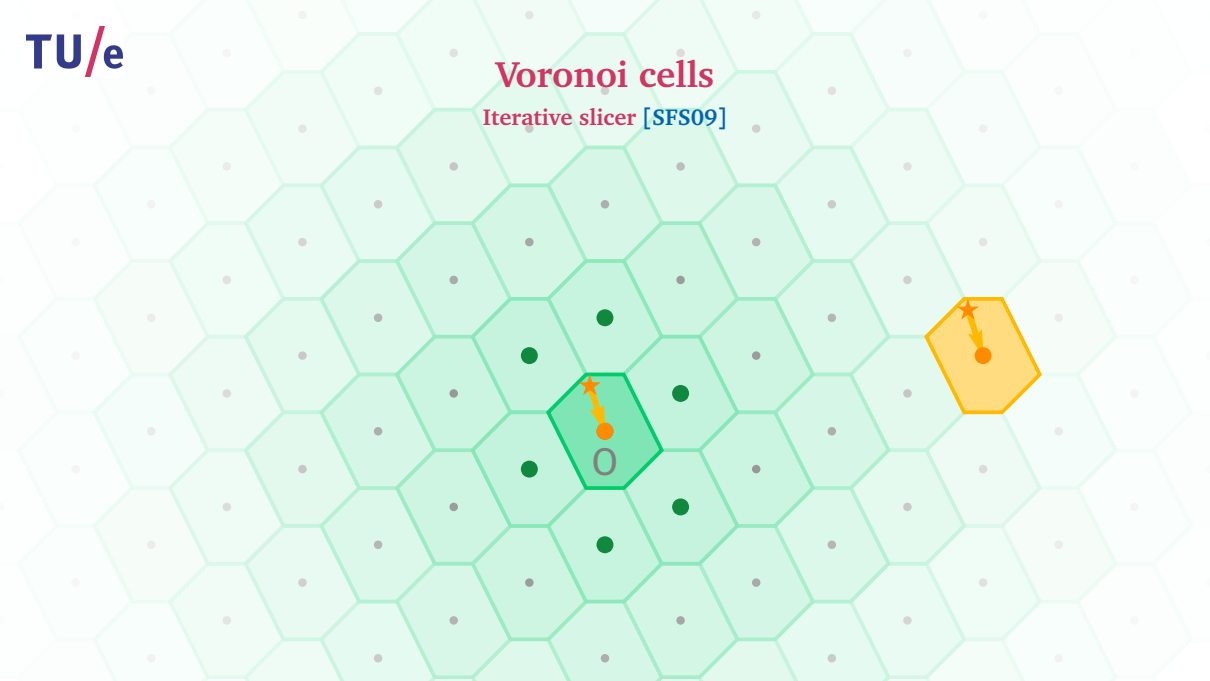
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Approximate Voronoi cells

Decrease list size



0

Approximate Voronoi cells

Decrease list size

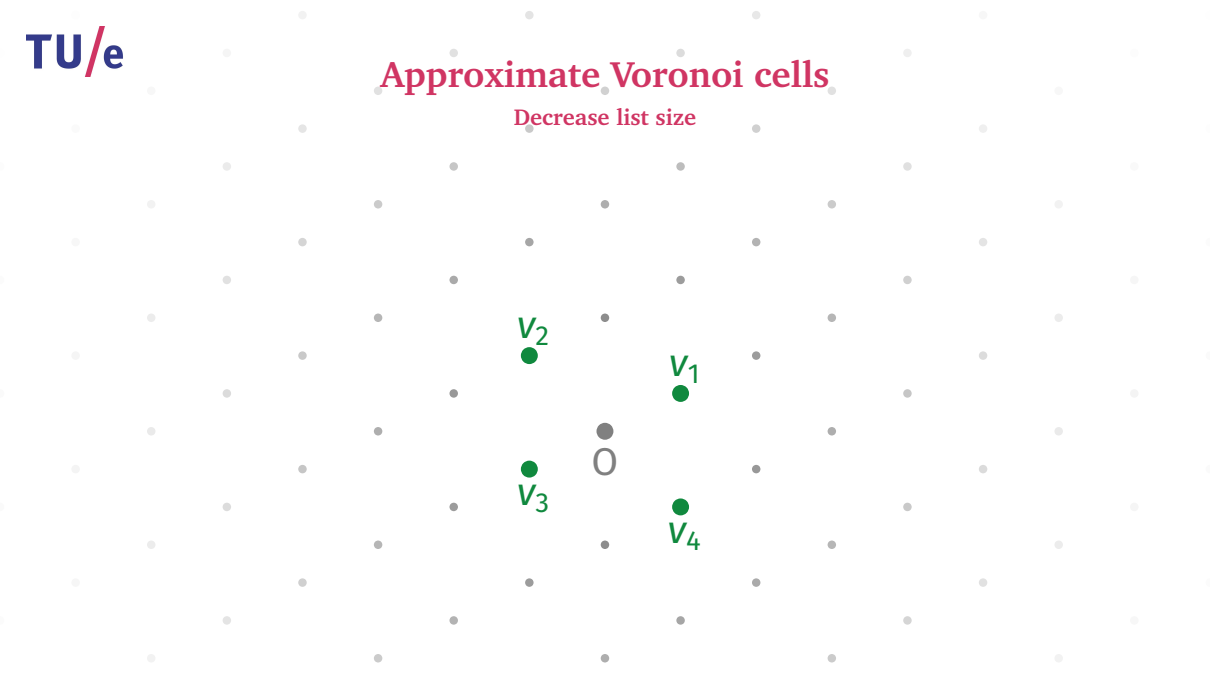
V_2

V_1

0

V_3

V_4



Approximate Voronoi cells

Decrease list size

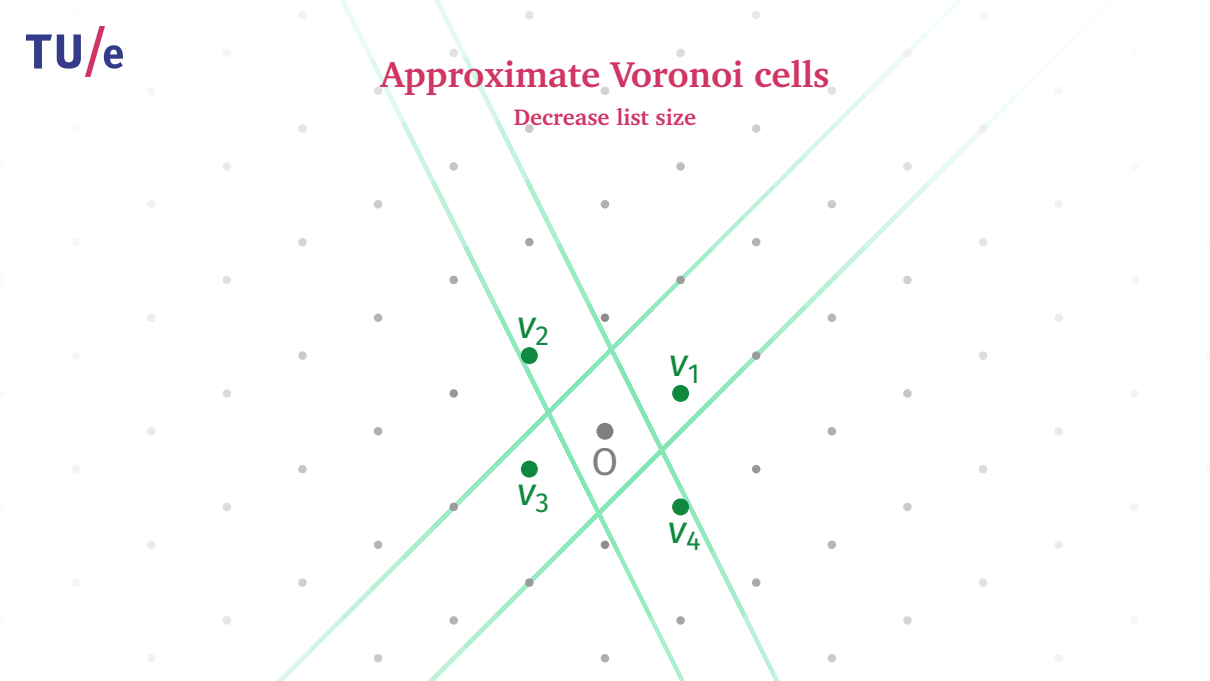
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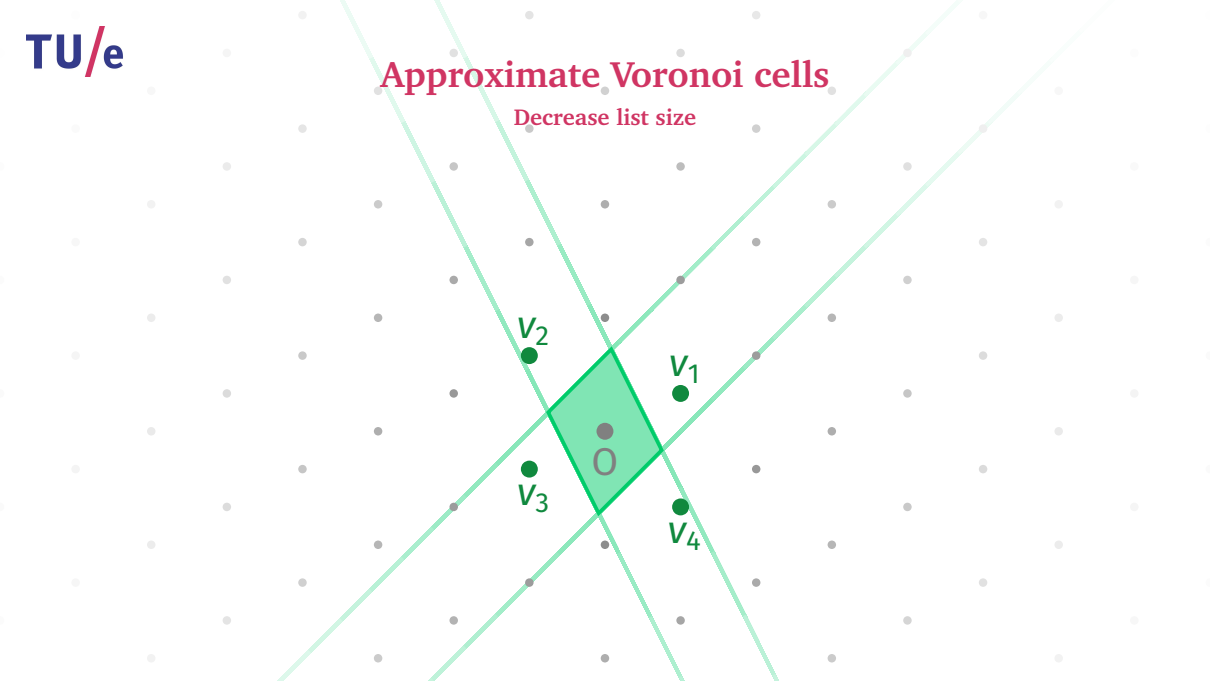
V_2

V_1

V_3

V_4

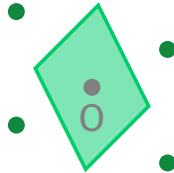
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The diagram illustrates the process of updating a Voronoi cell. A central point, labeled '0', is surrounded by four green lines that intersect at it, forming a green quadrilateral cell. This cell is labeled with V_1 , V_2 , V_3 , and V_4 at its vertices. The background is filled with a grid of small gray dots, representing a set of points. The text 'Approximate Voronoi cells' and 'Decrease list size' is displayed in red at the top. The TU/e logo is in the top left corner.

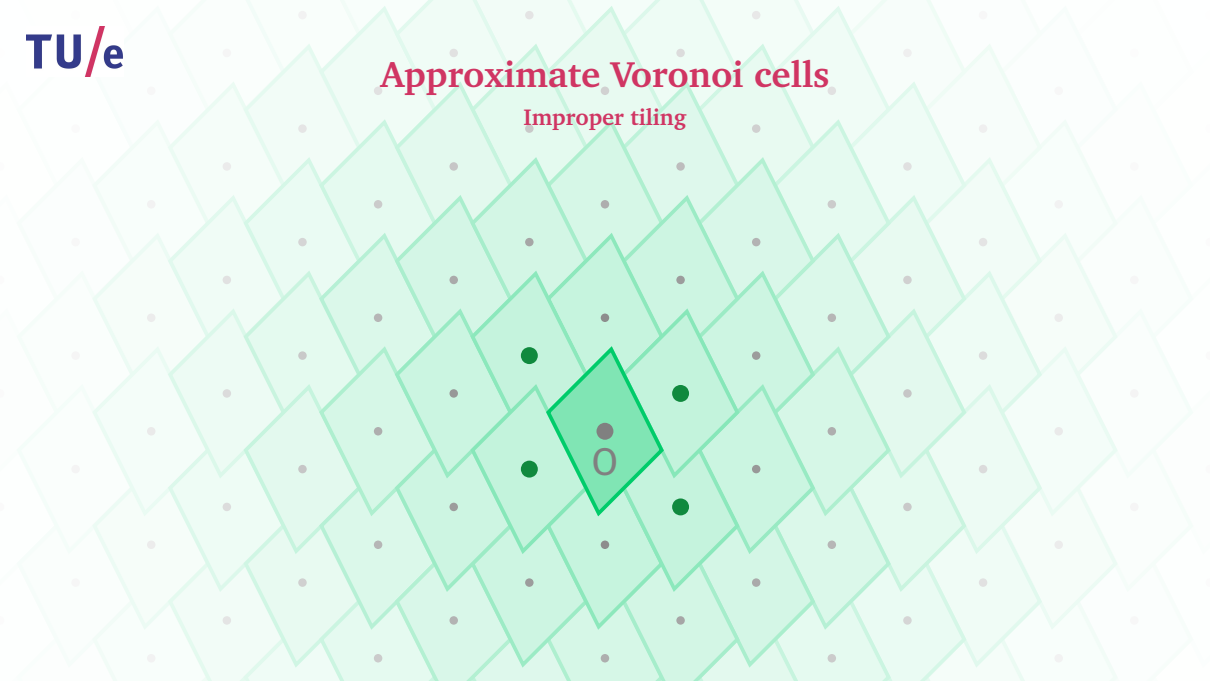
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Decrease list size



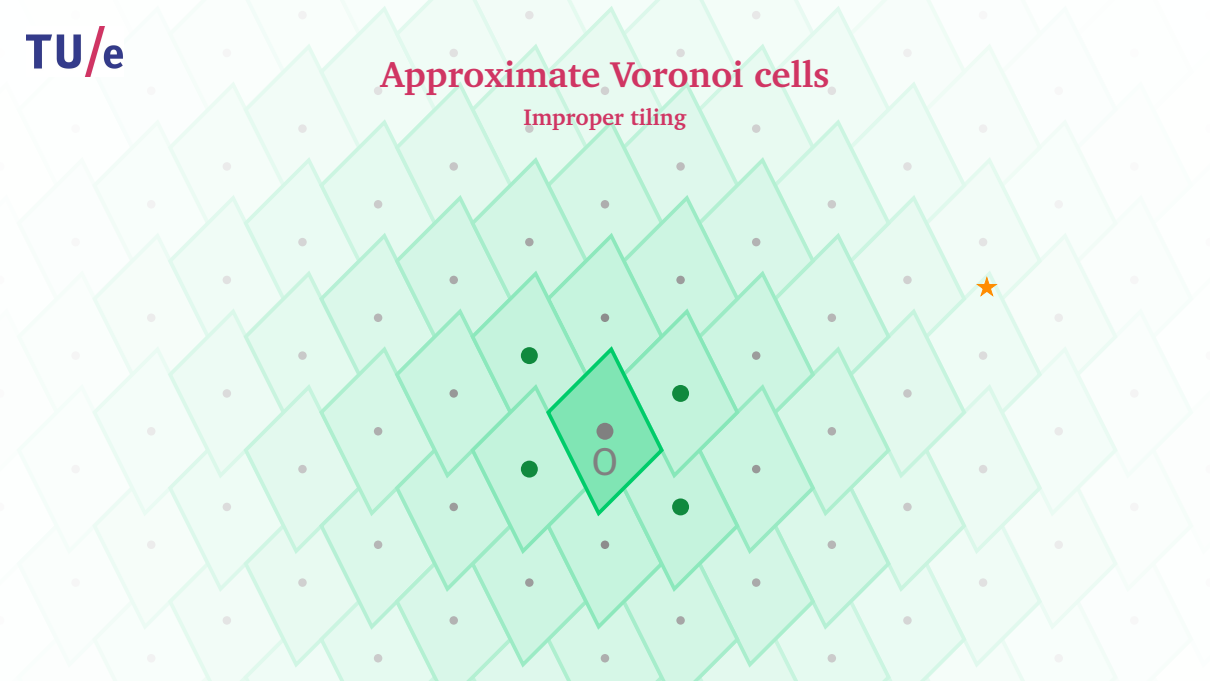
Approximate Voronoi cells

Improper tiling



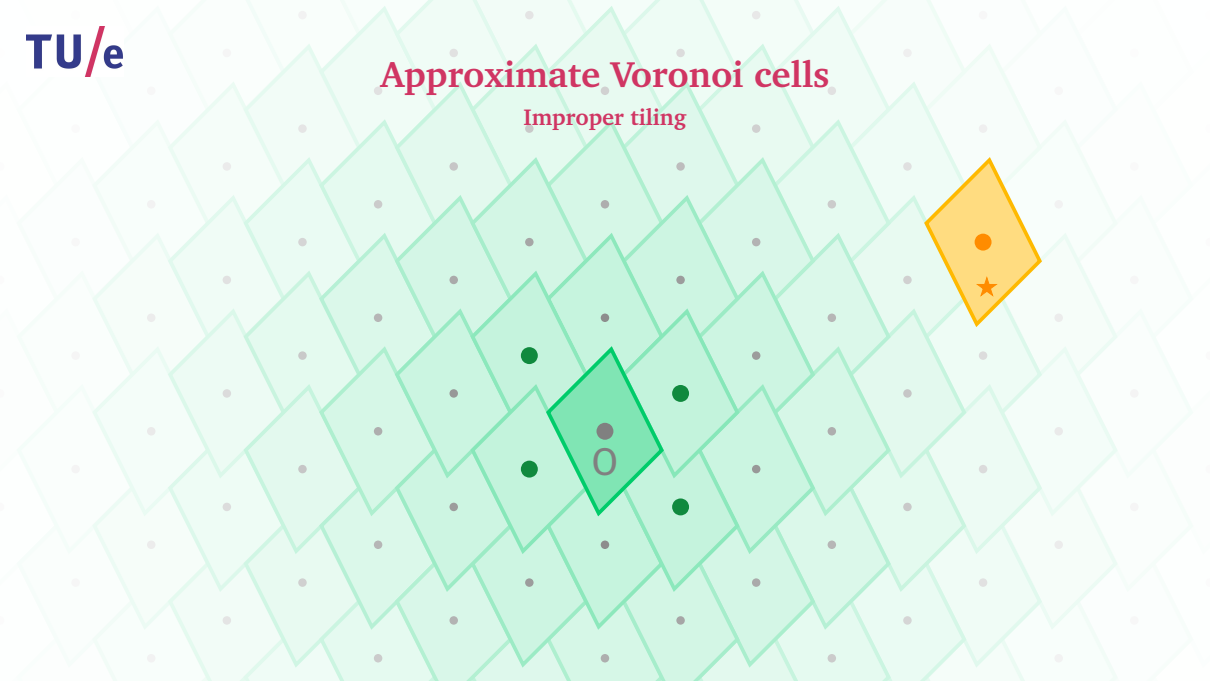
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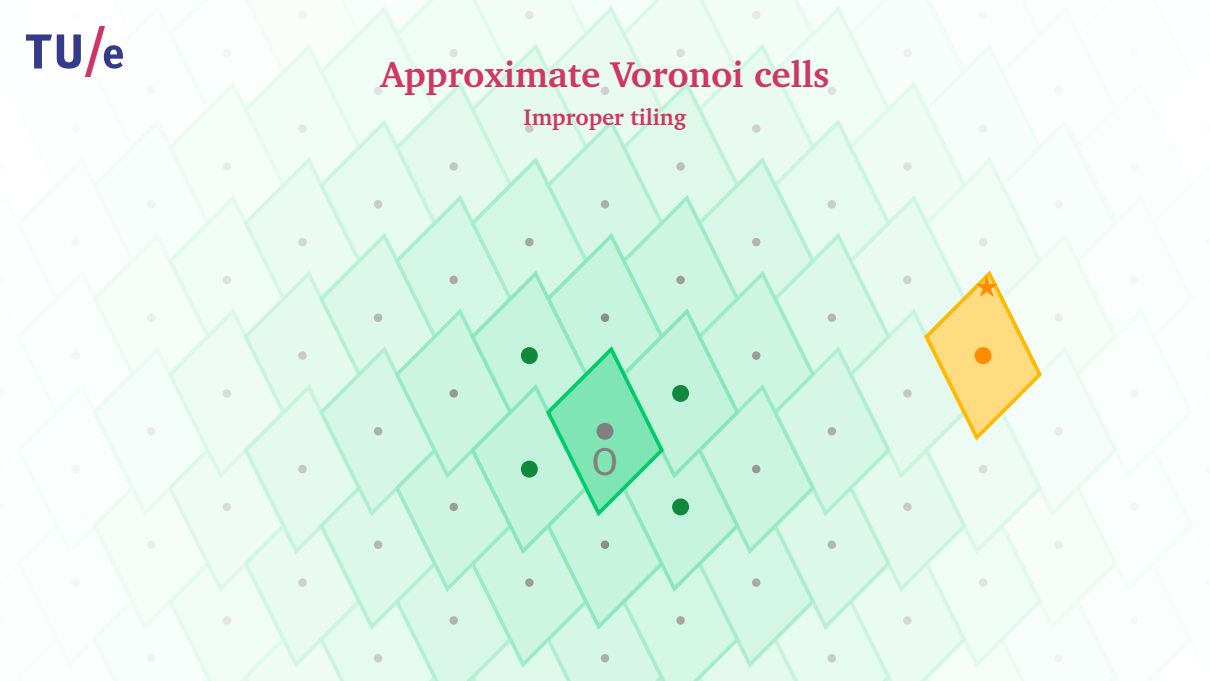
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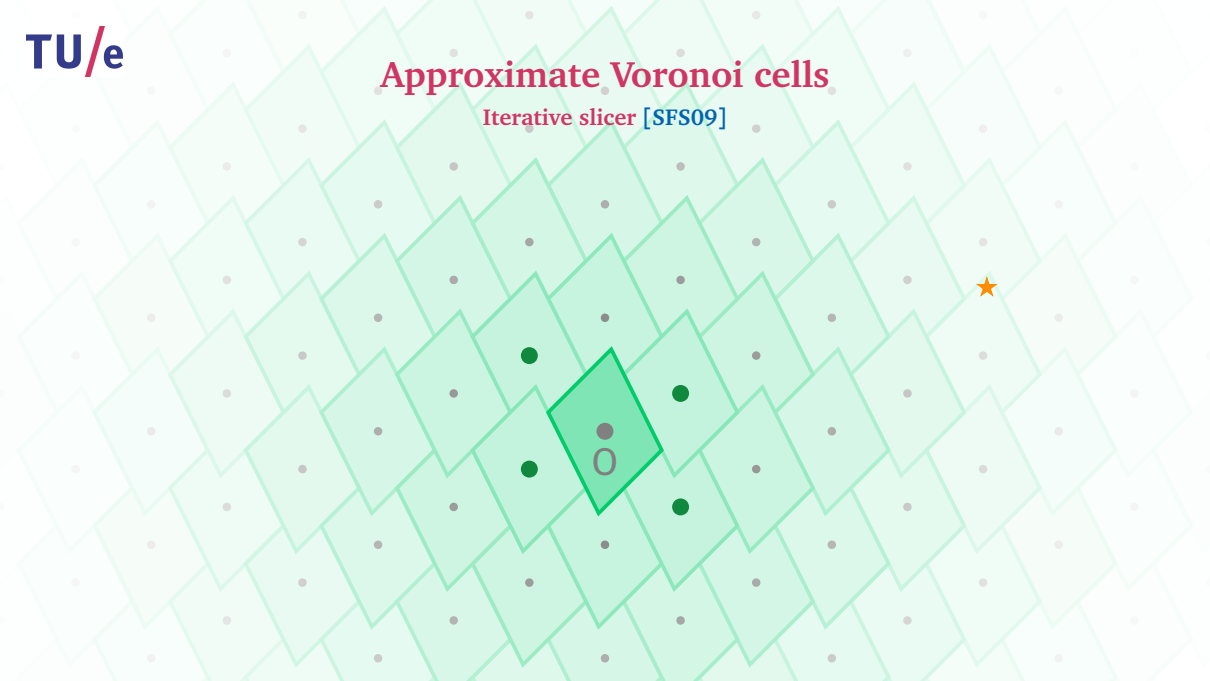
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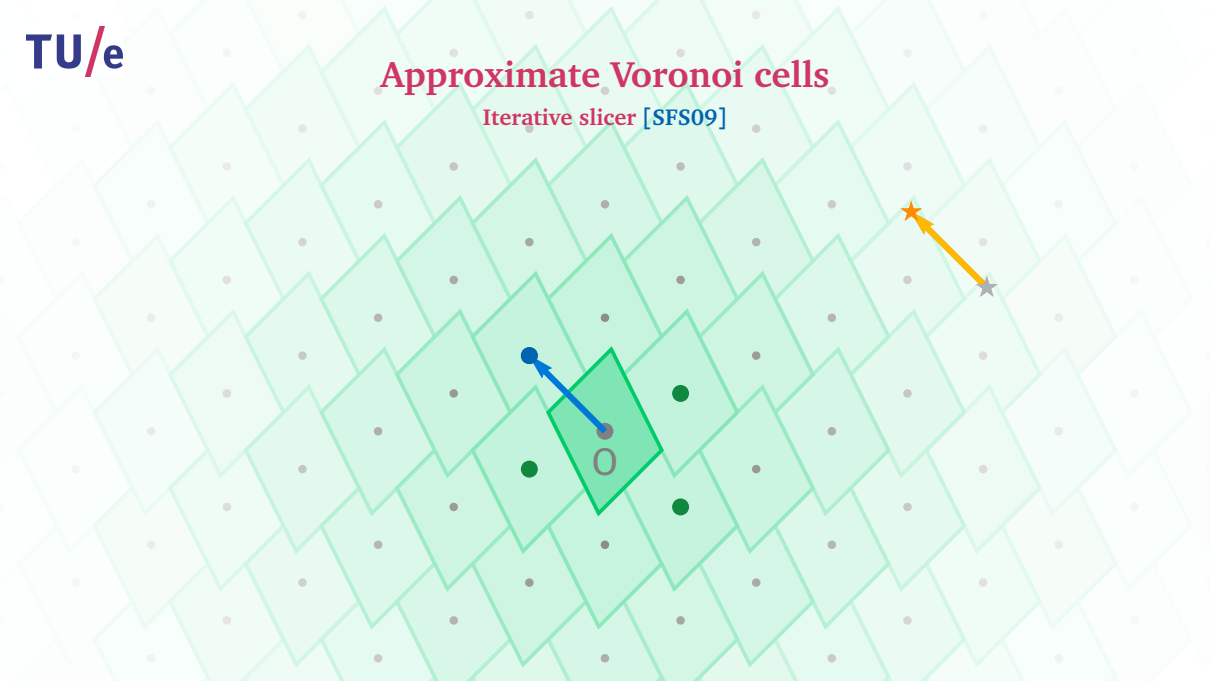
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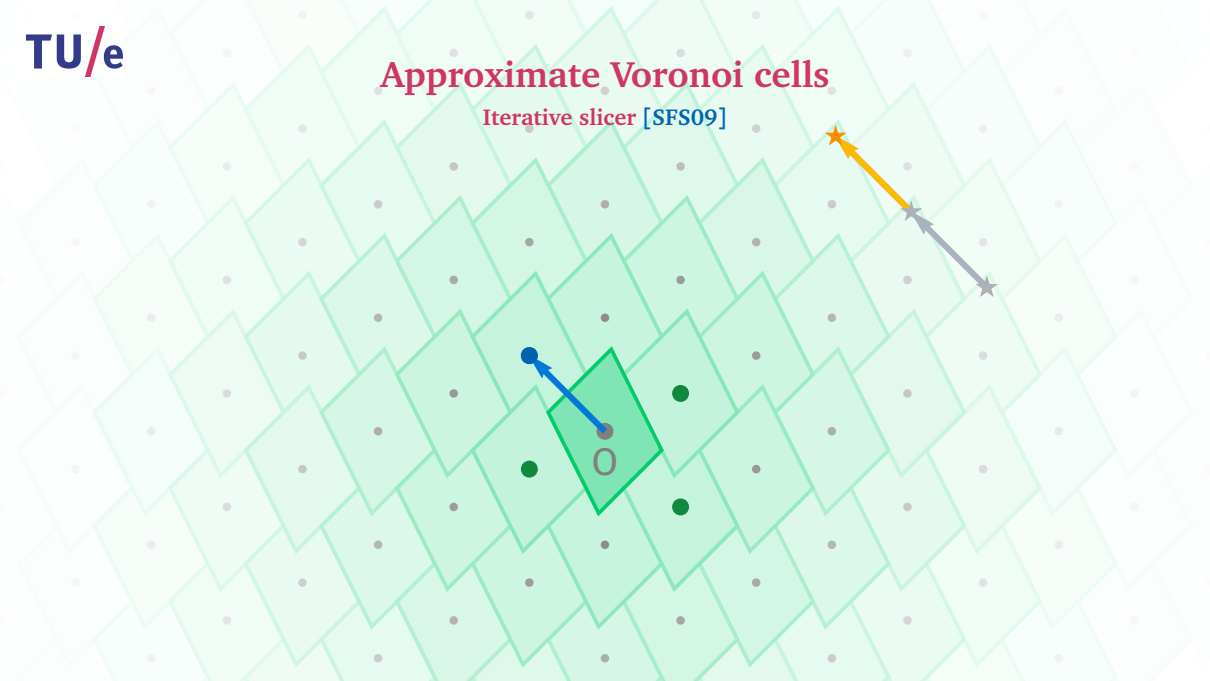
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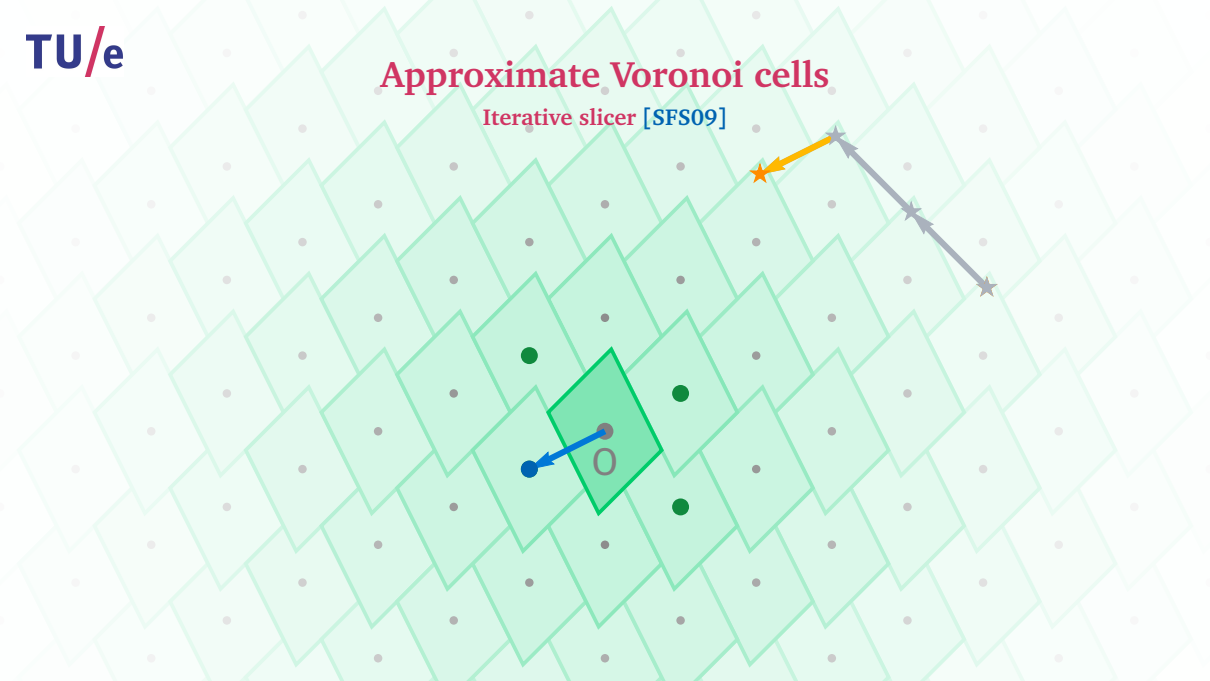
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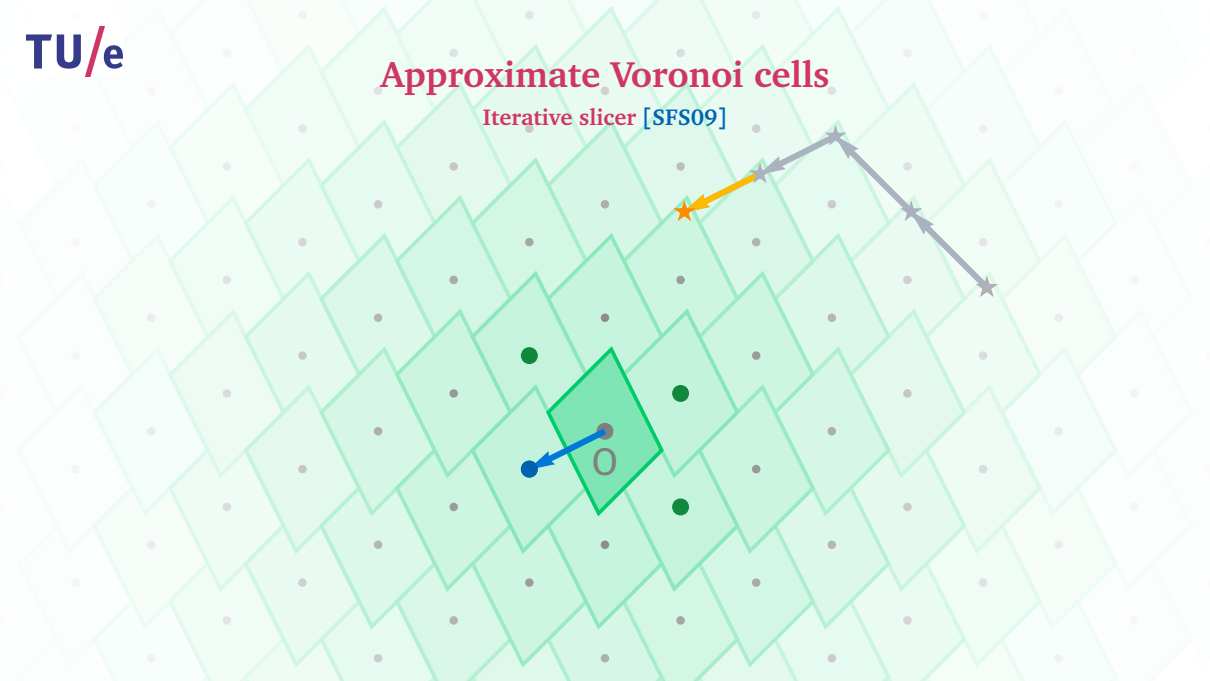
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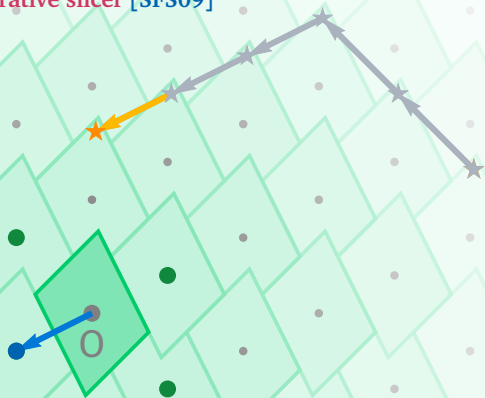
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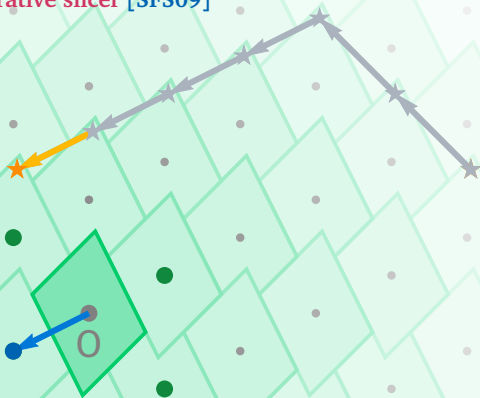
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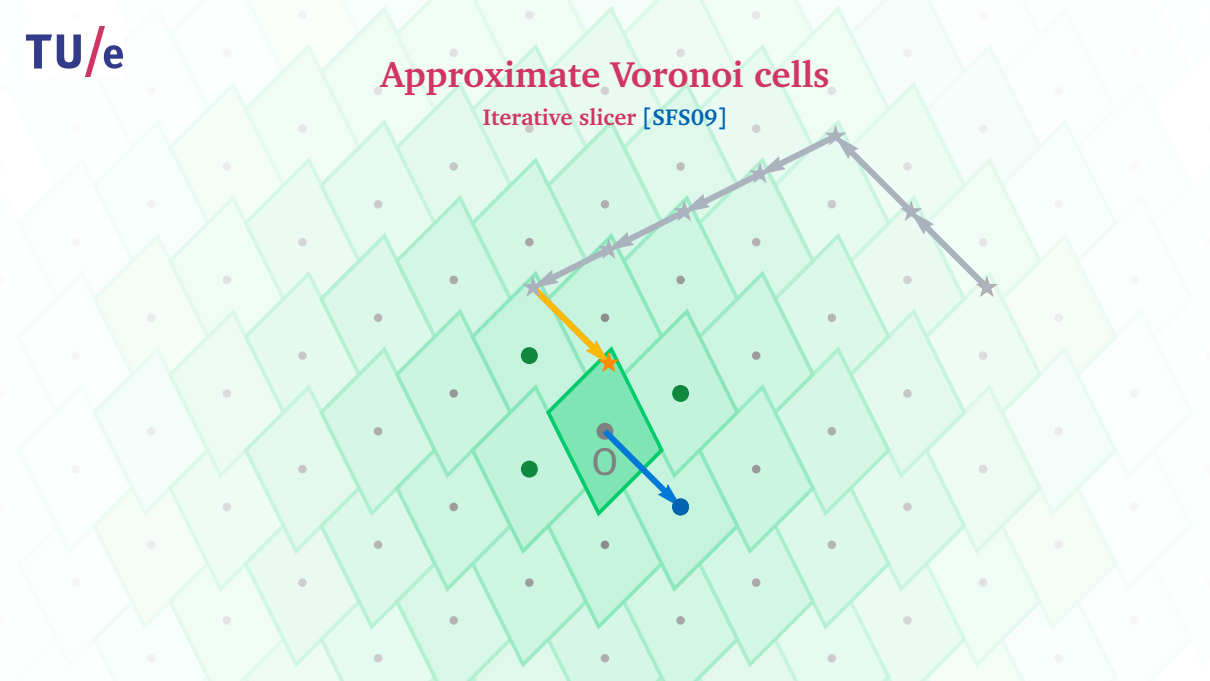
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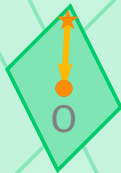
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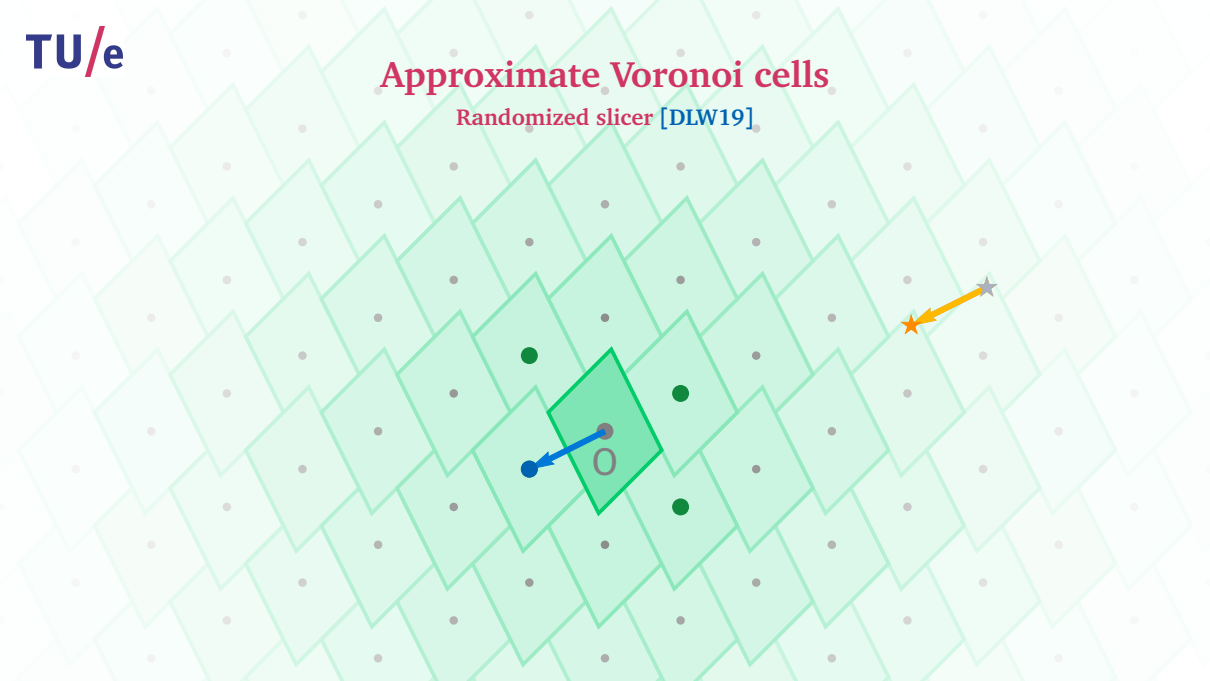
Approximate Voronoi cells

Randomized slicer [DLW19]



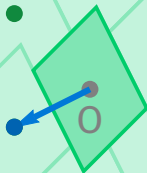
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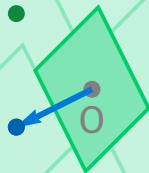
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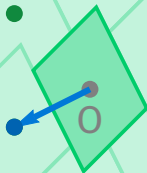
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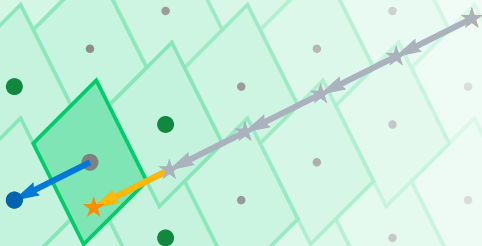
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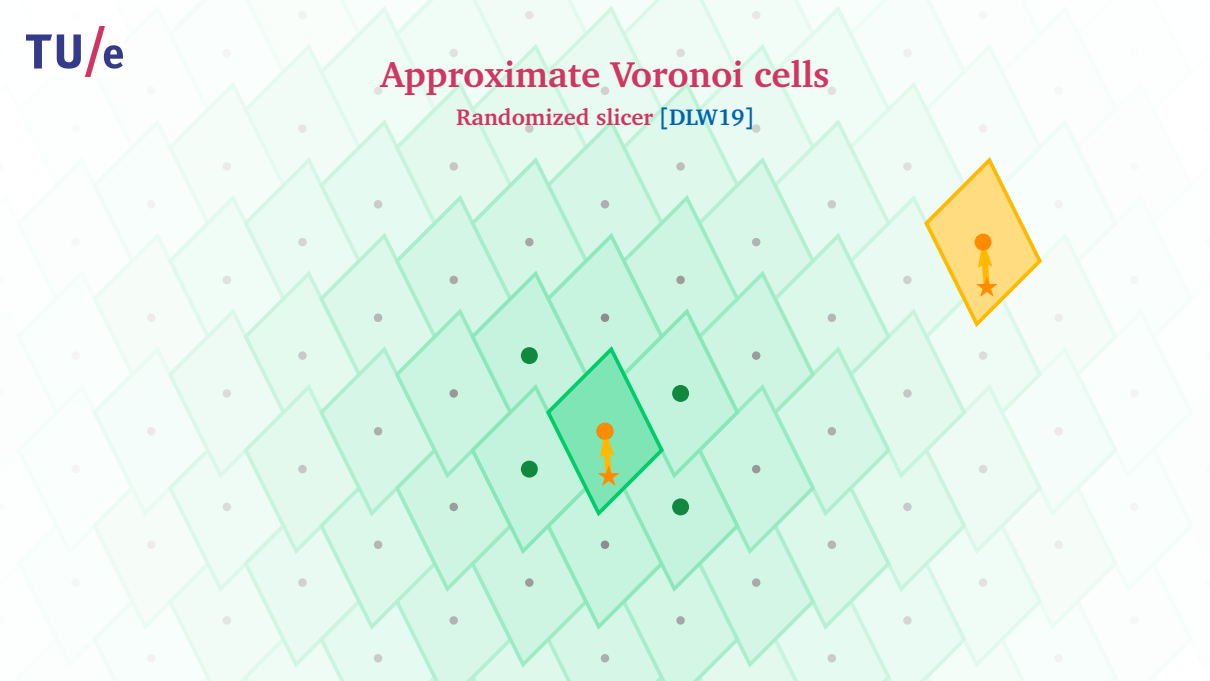
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Approximate Voronoi cells

Success probability estimation

Main problem: Success probability p of the iterative slicer?

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Attempt 1: [DLW19]

- Directly obtained a lower bound on p via the slicer
- Conjectured that p is exactly proportional to $\text{vol}(\mathcal{V})/\text{vol}(\mathcal{V}_L)$
- Open problem: obtain a tight analysis, perhaps via $\text{vol}(\mathcal{V})/\text{vol}(\mathcal{V}_L)$

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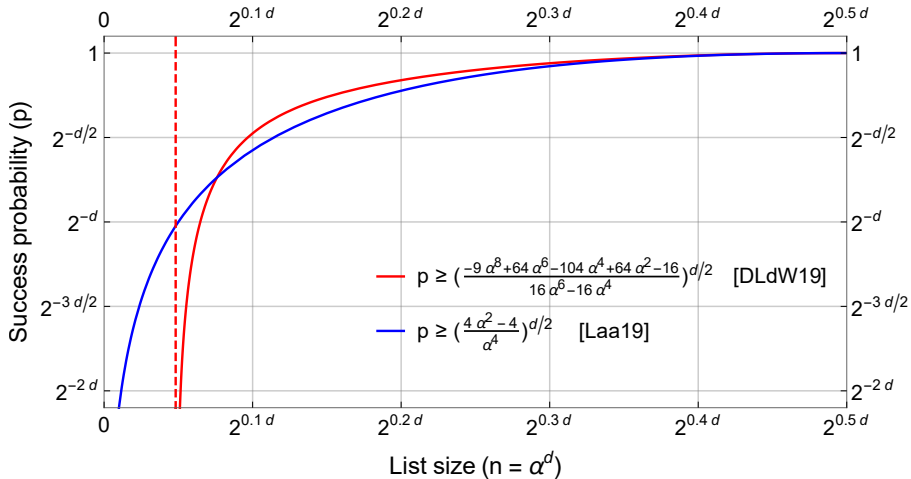
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Attempt 2: [Laa19]

- Proved tight bounds on $\text{vol}(\mathcal{V})/\text{vol}(\mathcal{V}_L)$ under the Gaussian heuristic
- Results show that p **cannot** be (exactly) proportional to $\text{vol}(\mathcal{V})/\text{vol}(\mathcal{V}_L)$
- From $p \geq \text{vol}(\mathcal{V})/\text{vol}(\mathcal{V}_L)$ we obtain new lower bounds on p
- No nonsensical asymptote at $2^{0.05d}$ memory anymore

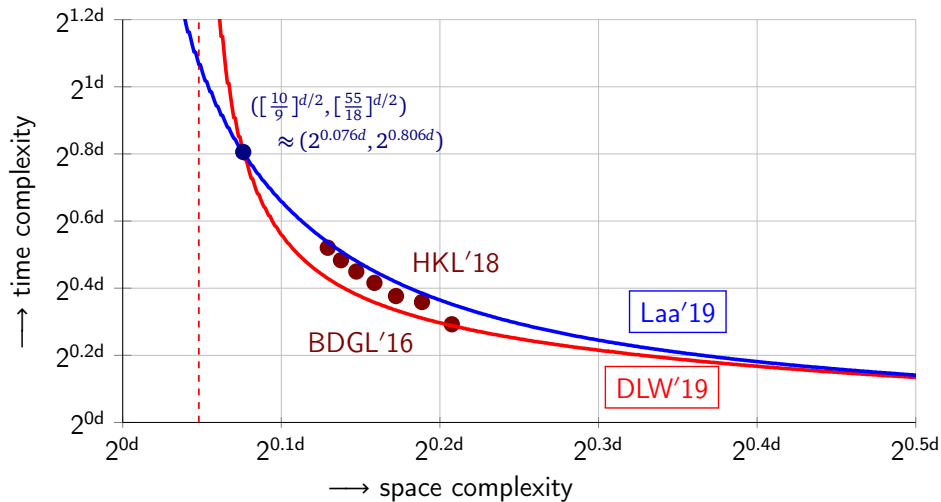
Approximate Voronoi cells

Lower bounds on success probability



Approximate Voronoi cells

Time-space trade-offs for CVPP



Conclusion

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 - ▶ Original analysis did not appear to be tight
 - ▶ Conjectured that tighter bounds may be obtained via $\text{vol}(\mathcal{V})/\text{vol}(\mathcal{V}_L)$
 - ▶ New results obtained tight bounds on the ratio $\text{vol}(\mathcal{V})/\text{vol}(\mathcal{V}_L)$
 - ▶ Resulted in better CVPP complexities for low-memory regime
 - ▶ Unfortunately, approach via $\text{vol}(\mathcal{V})/\text{vol}(\mathcal{V}_L)$ is not tight either

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Open problems

- Obtain truly tight bounds (ongoing work with Leo Ducas, Wessel van Woerden)
- Find an efficient BDDP-version of this CVPP algorithm