

Sieving for shortest lattice vectors using fast search algorithms

Thijs Laarhoven

(joint work with Michele Mosca, Joop van de Pol, Benne de Weger)

`mail@thijs.com`

`http://www.thijs.com/`

Rutgers University, New Brunswick, USA
(January 13, 2015)

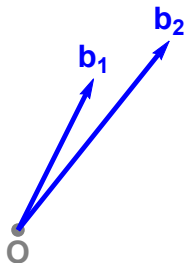
Lattices

What is a lattice?



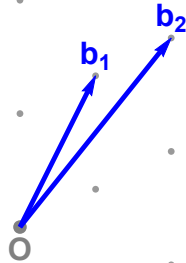
Lattices

What is a lattice?



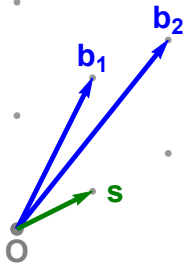
Lattices

What is a lattice?



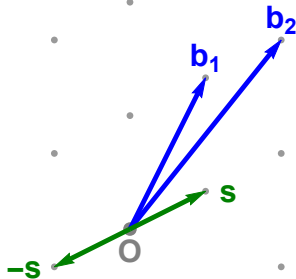
Lattices

Shortest Vector Problem (SVP)



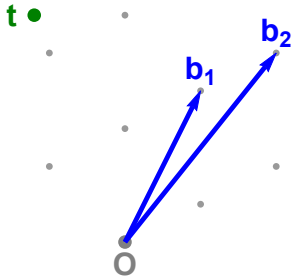
Lattices

Shortest Vector Problem (SVP)



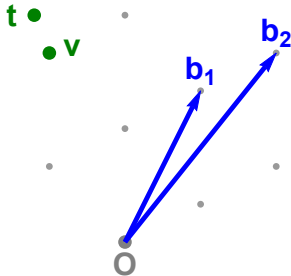
Lattices

Closest Vector Problem (CVP)



Lattices

Closest Vector Problem (CVP)



Lattices

Applications

- “Constructive cryptography”: Lattice-based cryptosystems
 - ▶ Based on hard lattice problems (SVP, CVP, LWE, SIS)
 - ▶ NTRU cryptosystem [HPS98, ...]
 - ▶ Fully Homomorphic Encryption [Gen09, ...]
 - ▶ Candidate for post-quantum cryptography

Lattices

Applications

- “Constructive cryptography”: Lattice-based cryptosystems
 - ▶ Based on hard lattice problems (SVP, CVP, LWE, SIS)
 - ▶ NTRU cryptosystem [HPS98, ...]
 - ▶ Fully Homomorphic Encryption [Gen09, ...]
 - ▶ Candidate for post-quantum cryptography
- “Destructive cryptography”: Lattice cryptanalysis
 - ▶ Attack knapsack-based cryptosystems [Sha82, LO85]
 - ▶ Attack RSA with Coppersmith’s method [Cop97]
 - ▶ Attack lattice-based cryptosystems [Ngu99, JJ00]

Lattices

Applications

- “Constructive cryptography”: Lattice-based cryptosystems
 - ▶ Based on hard lattice problems (SVP, CVP, LWE, SIS)
 - ▶ NTRU cryptosystem [HPS98, ...]
 - ▶ Fully Homomorphic Encryption [Gen09, ...]
 - ▶ Candidate for post-quantum cryptography
- “Destructive cryptography”: Lattice cryptanalysis
 - ▶ Attack knapsack-based cryptosystems [Sha82, LO85]
 - ▶ Attack RSA with Coppersmith’s method [Cop97]
 - ▶ Attack lattice-based cryptosystems [Ngu99, JJ00]

How hard are hard lattice problems such as SVP?

Nguyen-Vidick sieve

0

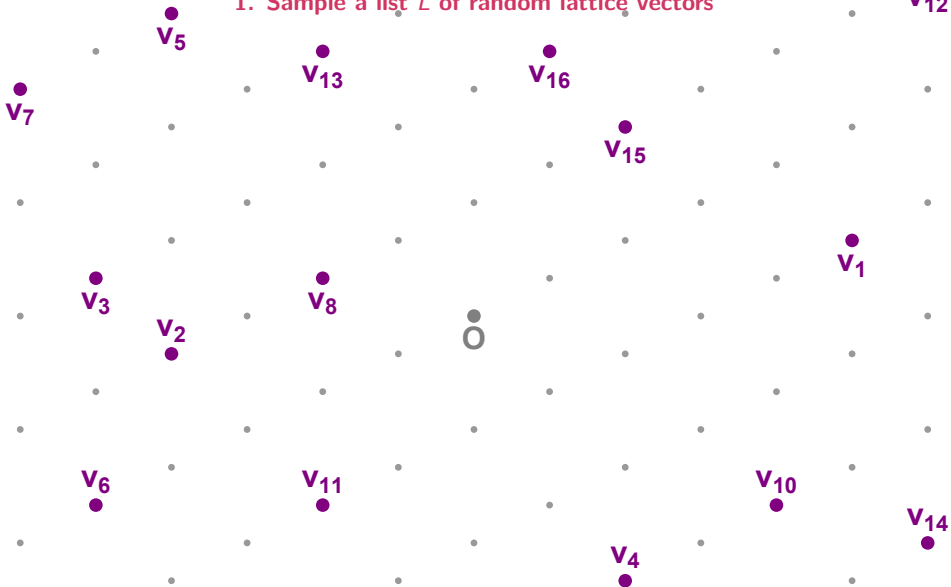
Nguyen-Vidick sieve

1. Sample a list L of random lattice vectors



Nguyen-Vidick sieve

1. Sample a list L of random lattice vectors



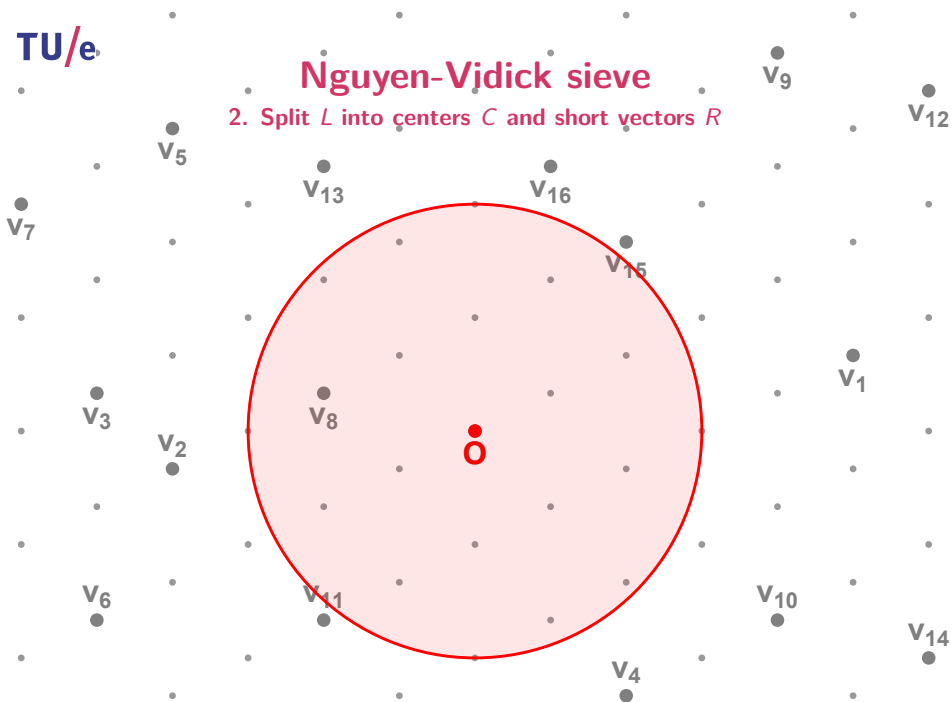
Nguyen-Vidick sieve

2. Split L into centers C and short vectors R



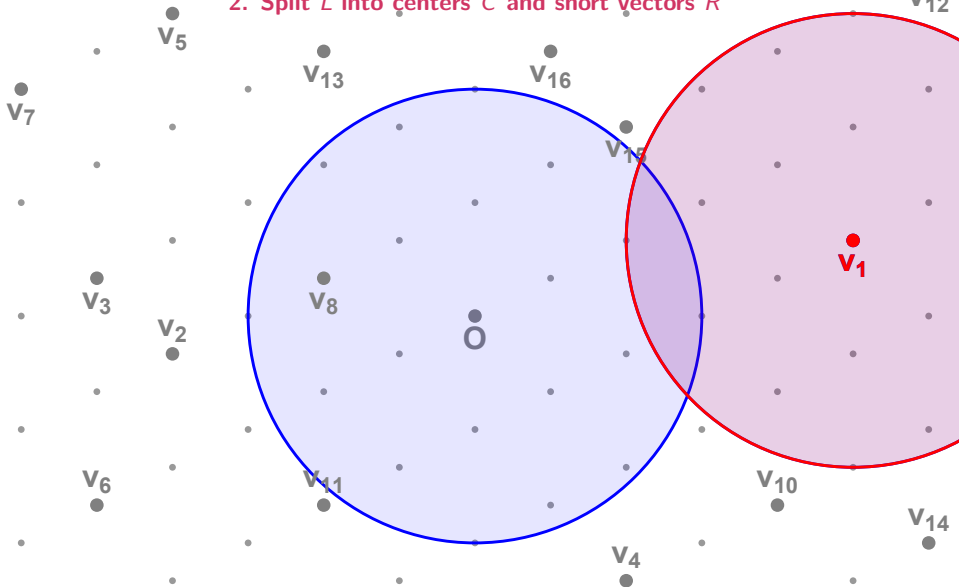
Nguyen-Vidick sieve

2. Split L into centers C and short vectors R



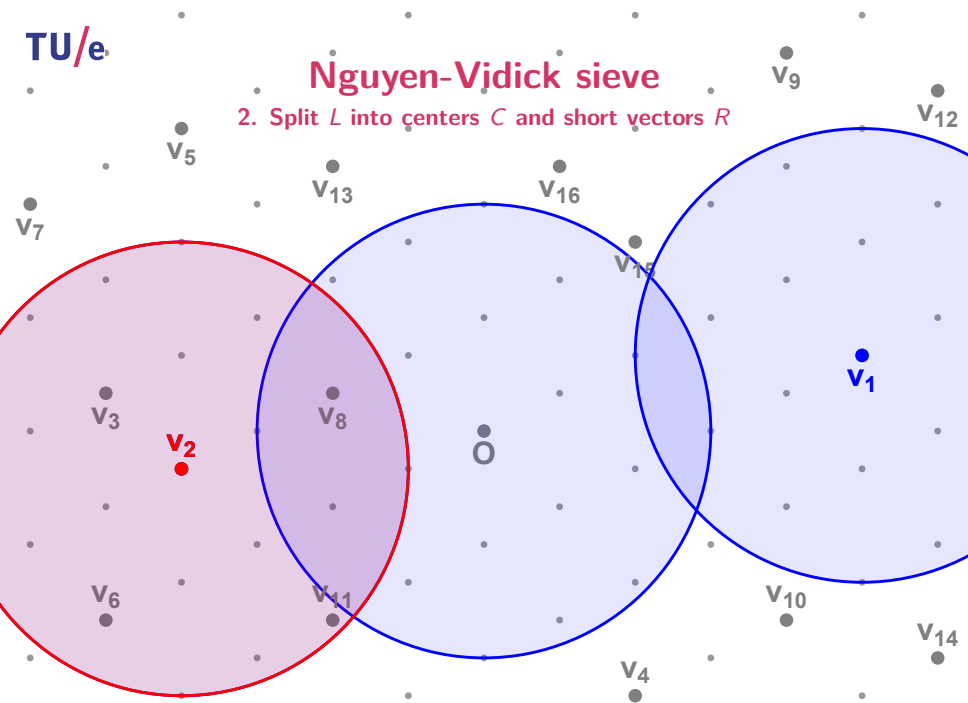
Nguyen-Vidick sieve

2. Split L into centers C and short vectors R



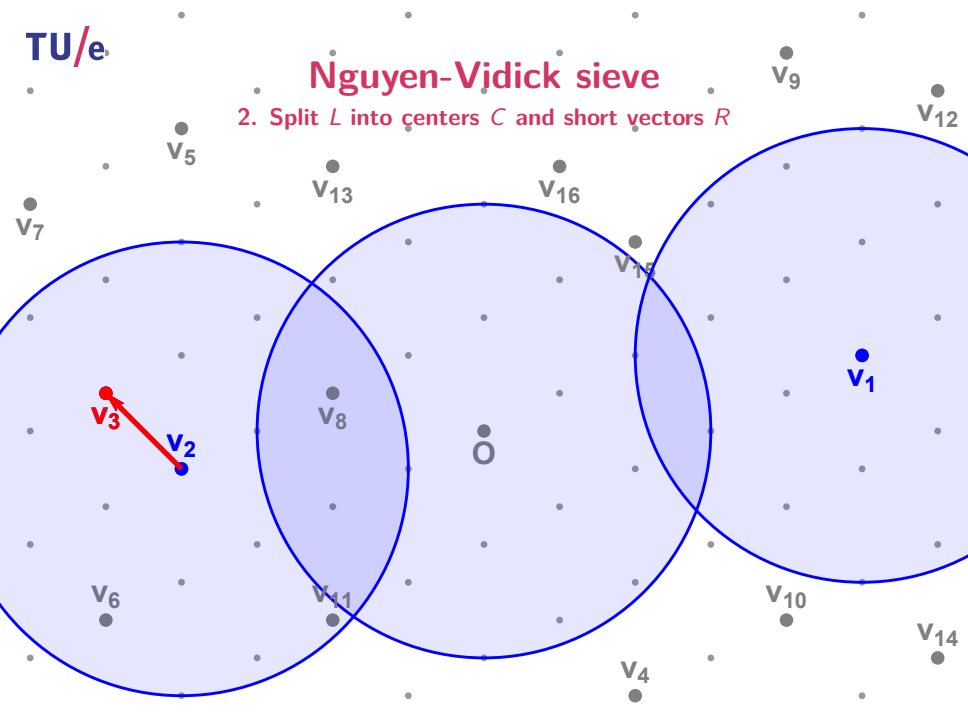
Nguyen-Vidick sieve

2. Split L into centers C and short vectors R



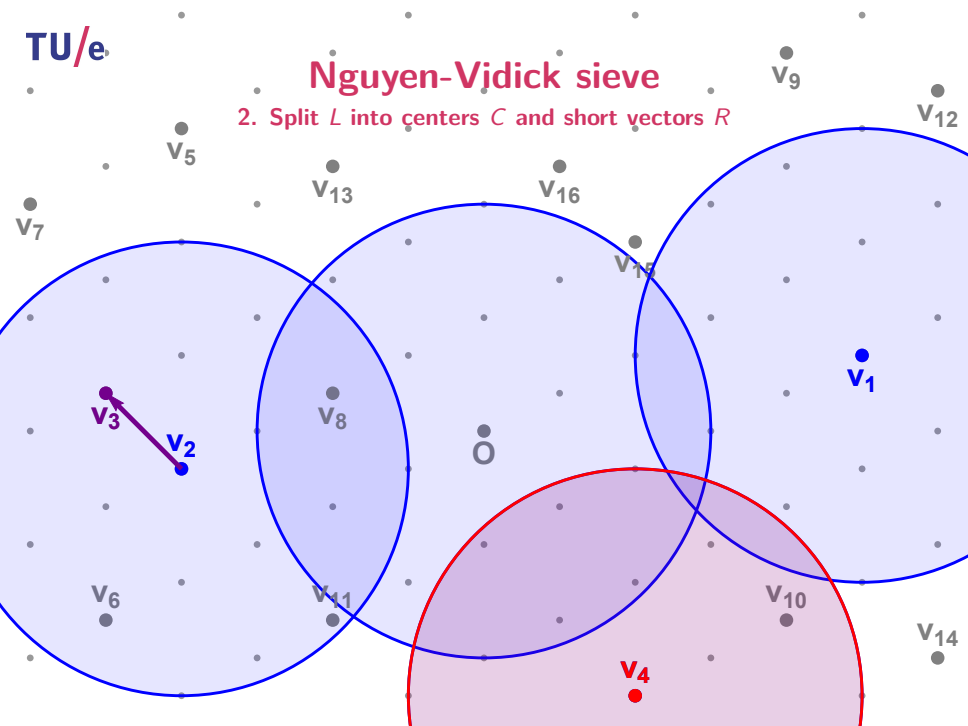
Nguyen-Vidick sieve

2. Split L into centers C and short vectors R



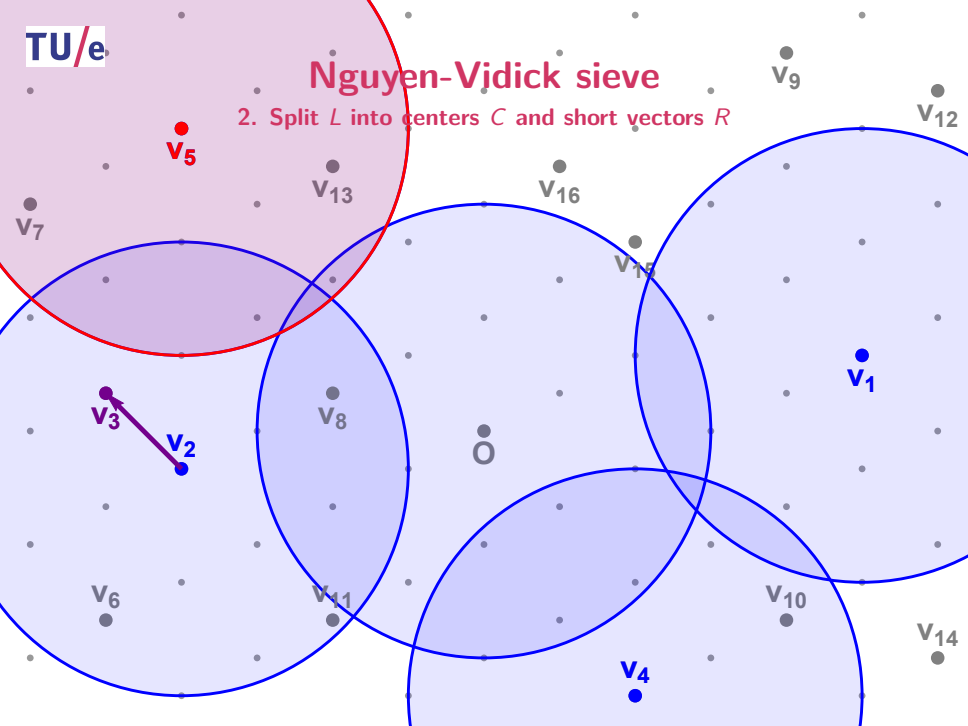
Nguyen-Vidick sieve

2. Split L into centers C and short vectors R



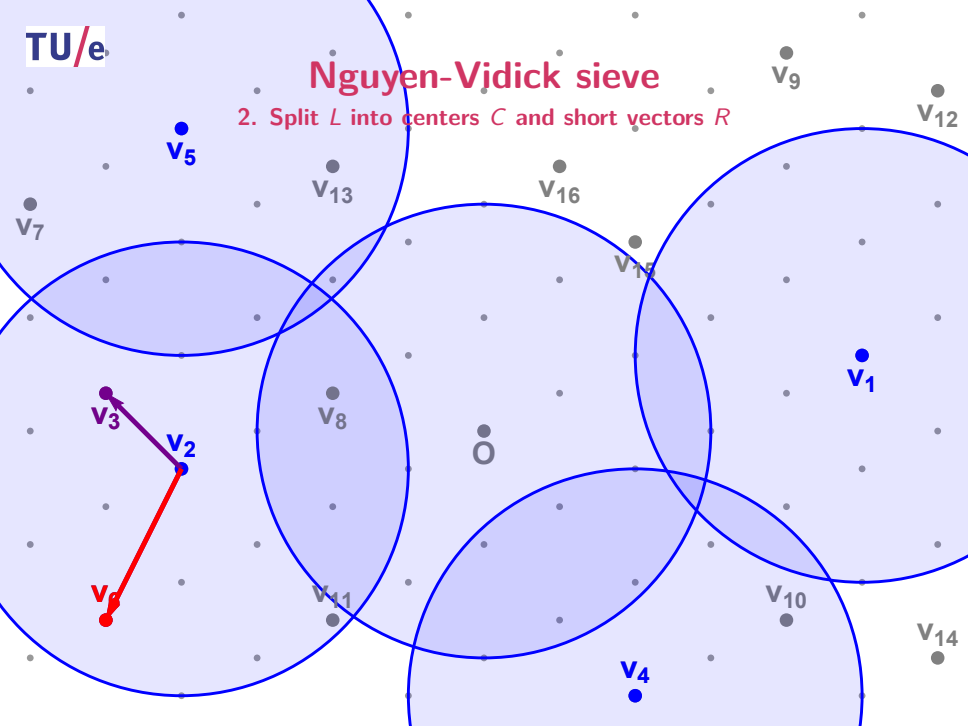
Nguyen-Vidick sieve

2. Split L into centers C and short vectors R



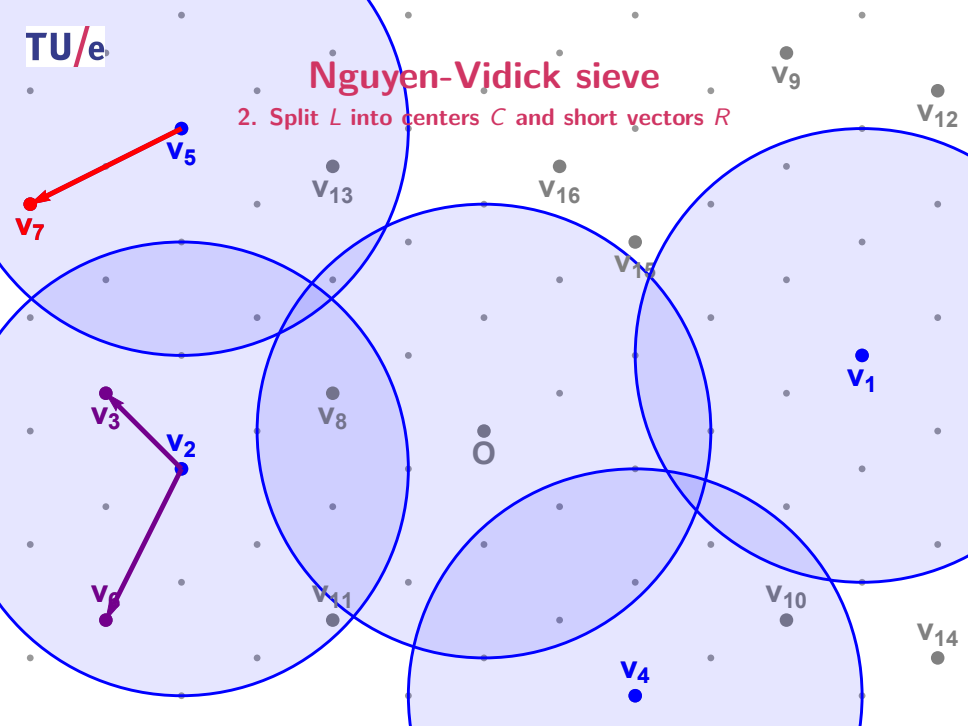
Nguyen-Vidick sieve

2. Split L into centers C and short vectors R



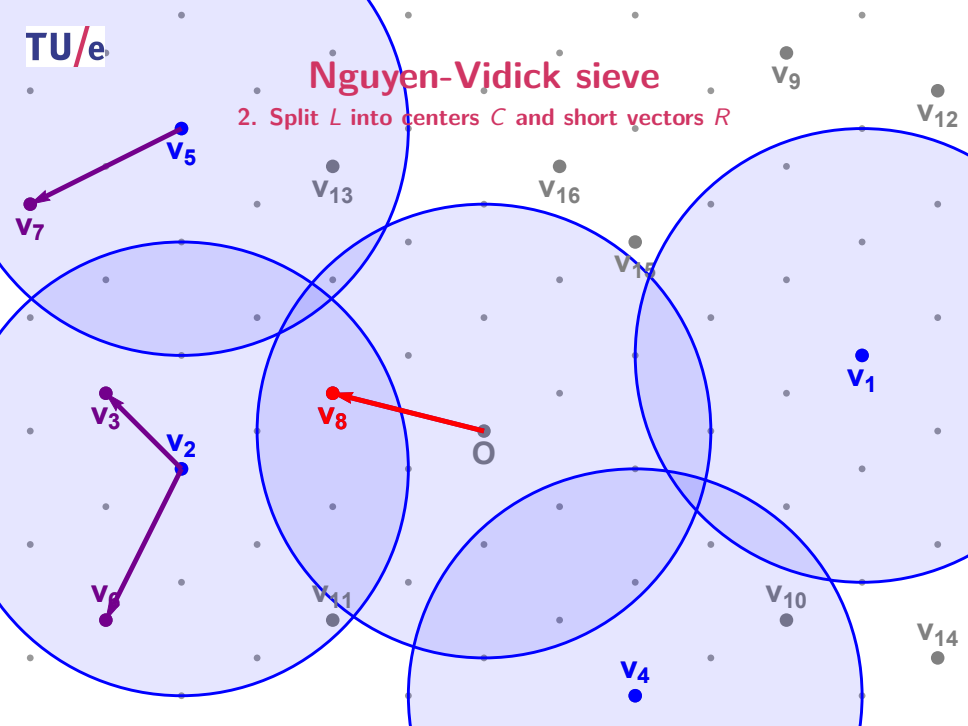
Nguyen-Vidick sieve

2. Split L into centers C and short vectors R



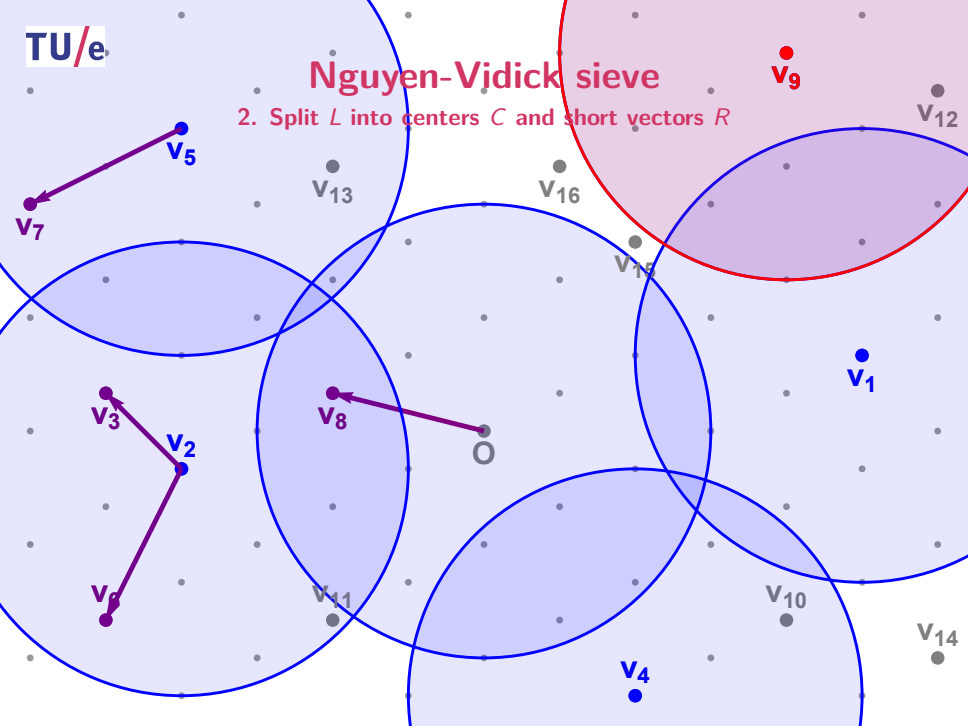
Nguyen-Vidick sieve

2. Split L into centers C and short vectors R



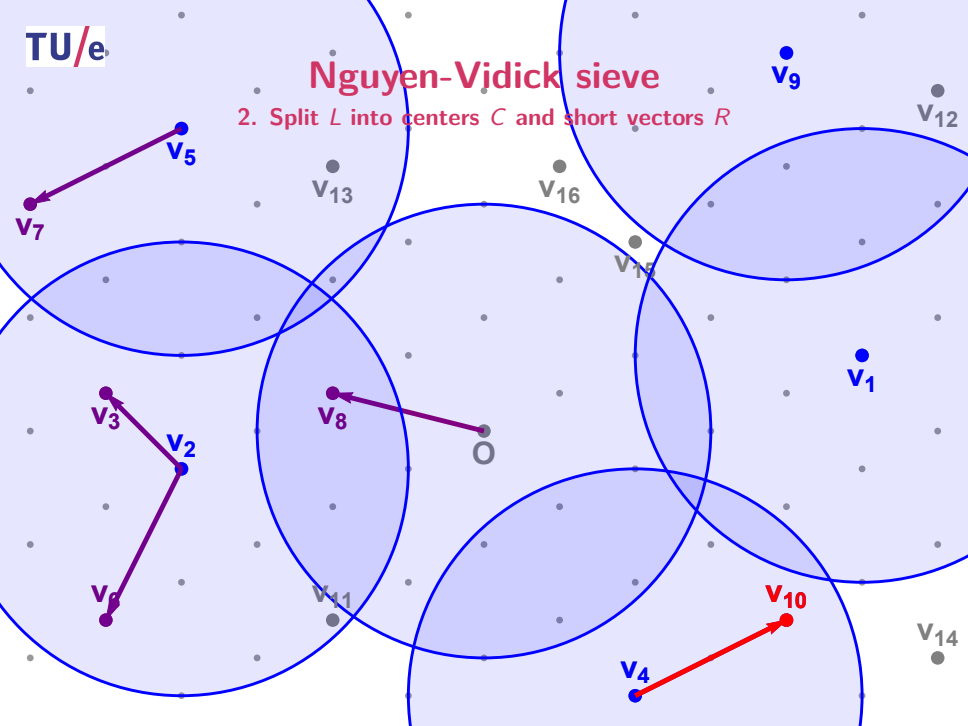
Nguyen-Vidick sieve

2. Split L into centers C and short vectors R



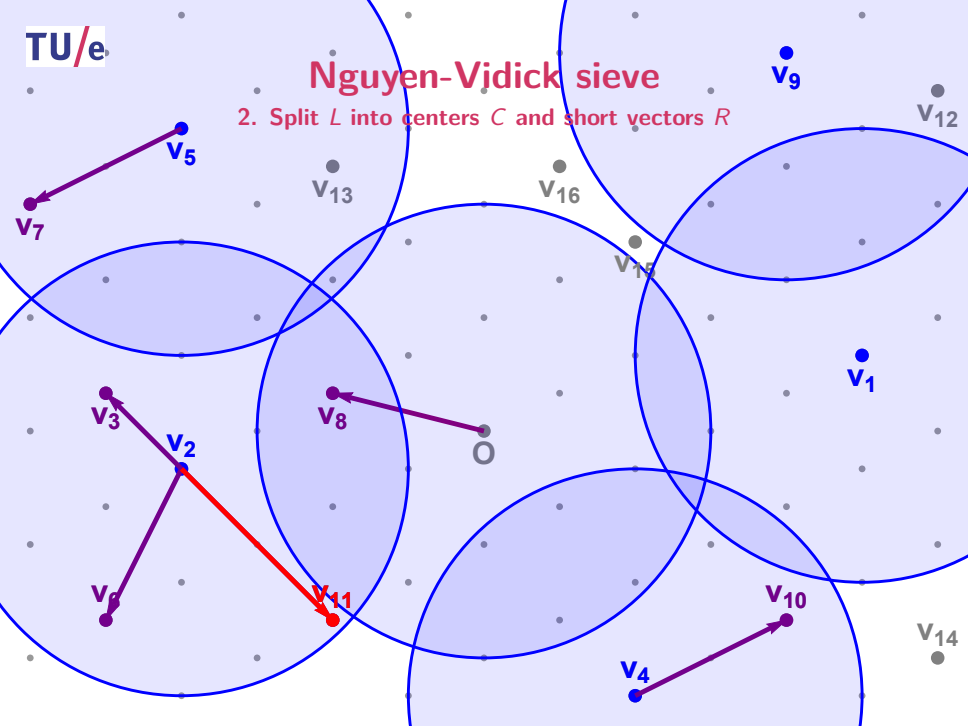
Nguyen-Vidick sieve

2. Split L into centers C and short vectors R



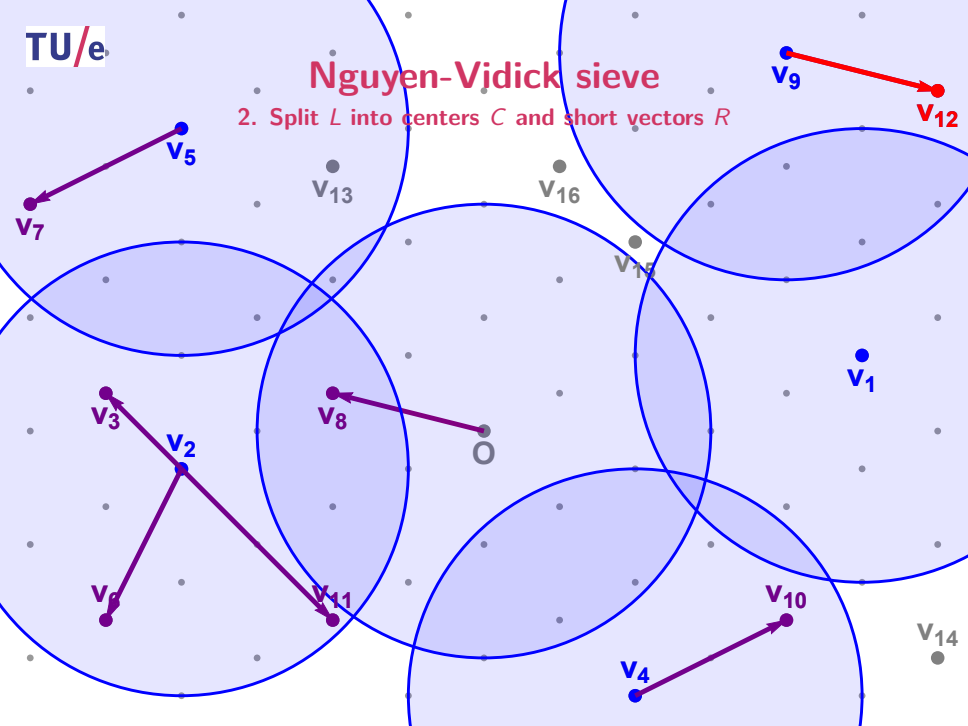
Nguyen-Vidick sieve

2. Split L into centers C and short vectors R



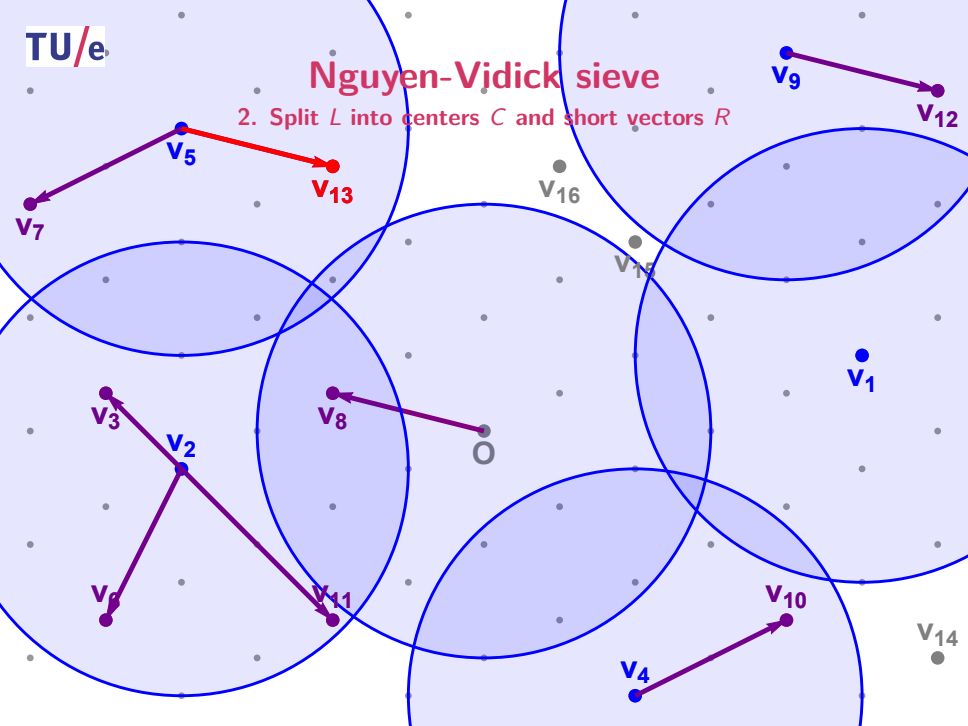
Nguyen-Vidick sieve

2. Split L into centers C and short vectors R



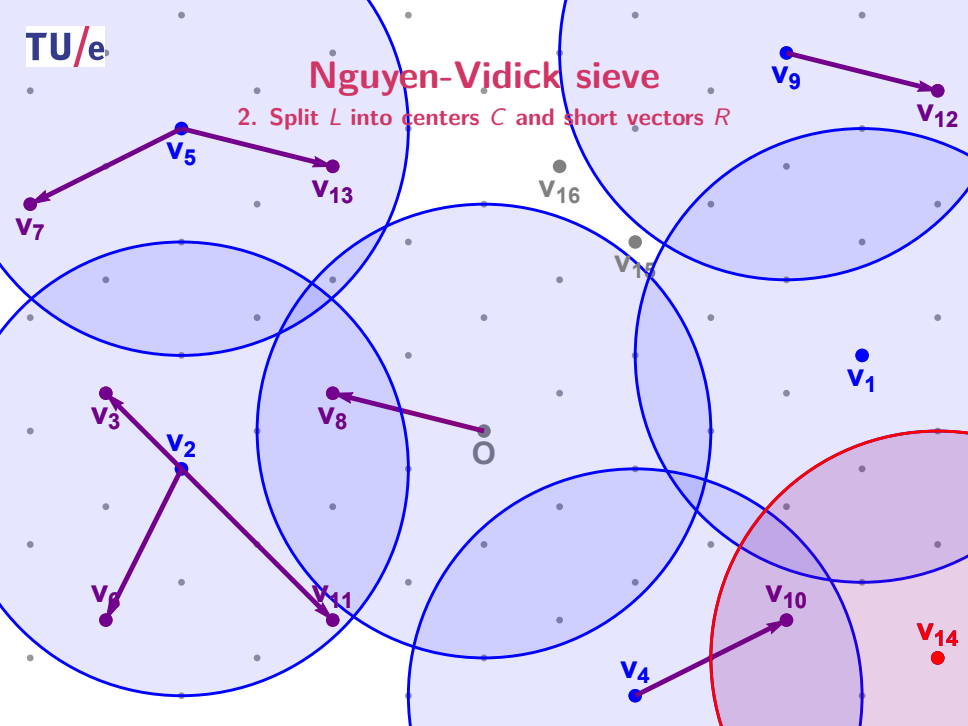
Nguyen-Vidick sieve

2. Split L into centers C and short vectors R



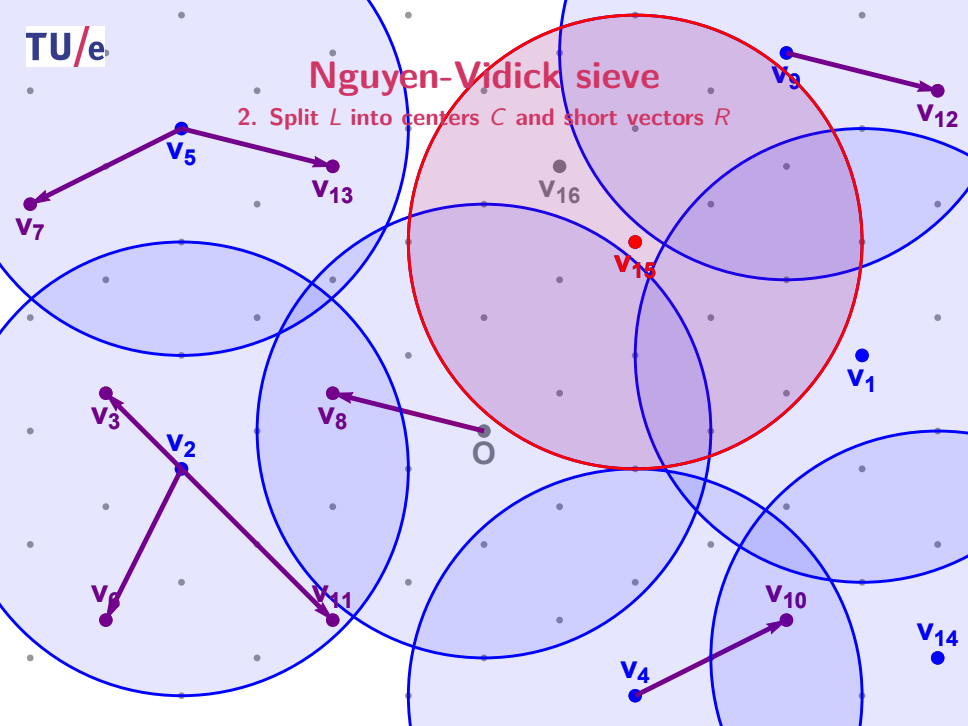
Nguyen-Vidick sieve

2. Split L into centers C and short vectors R



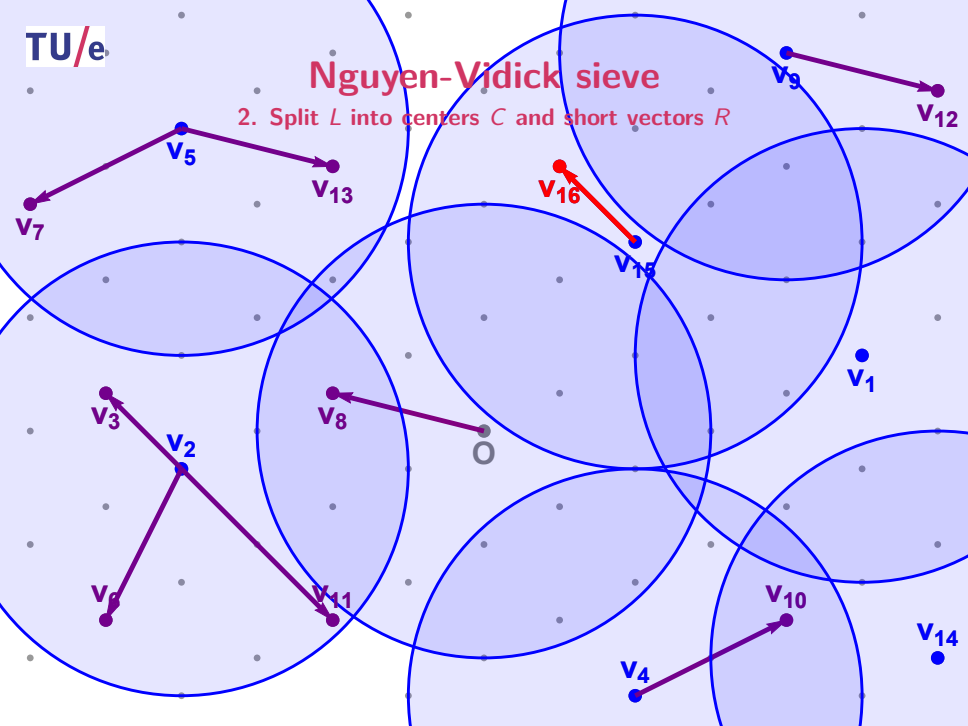
Nguyen-Vidick sieve

2. Split L into centers C and short vectors R



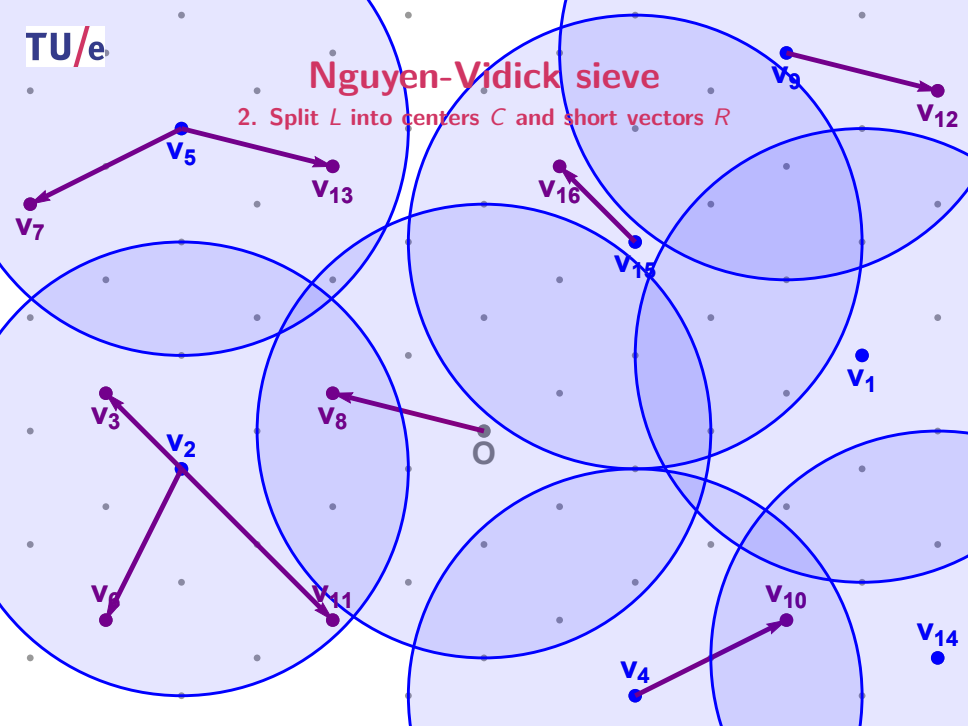
Nguyen-Vidick sieve

2. Split L into centers C and short vectors R



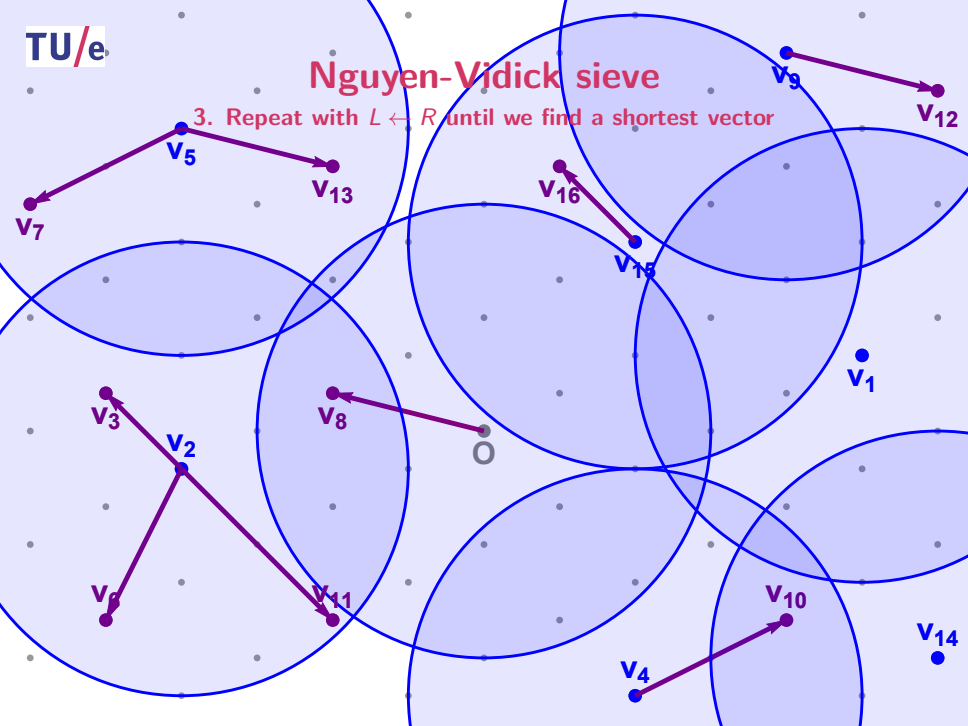
Nguyen-Vidick sieve

2. Split L into centers C and short vectors R



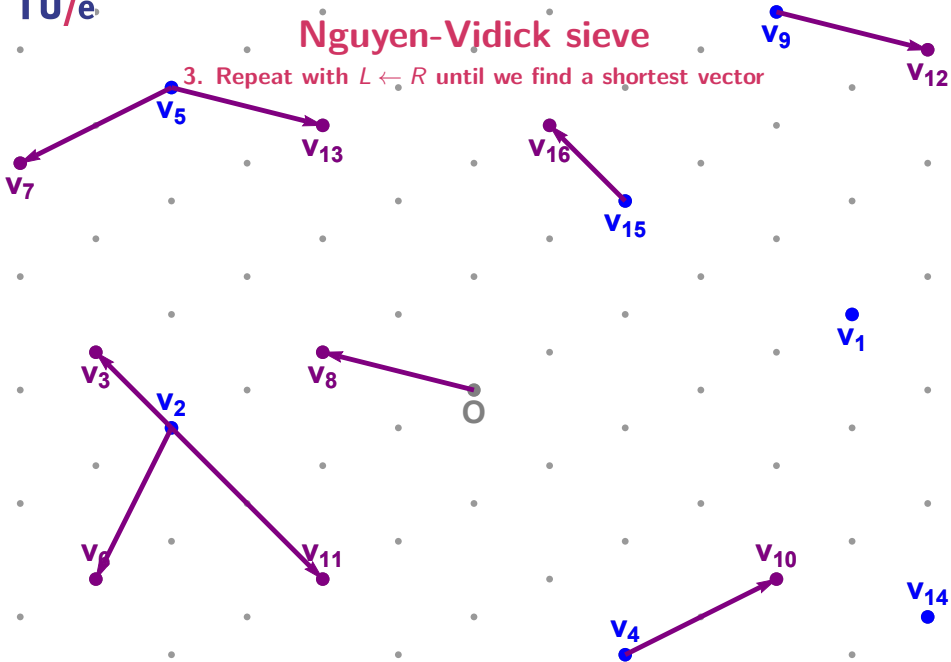
Nguyen-Vidick sieve

3. Repeat with $L \leftarrow R$ until we find a shortest vector



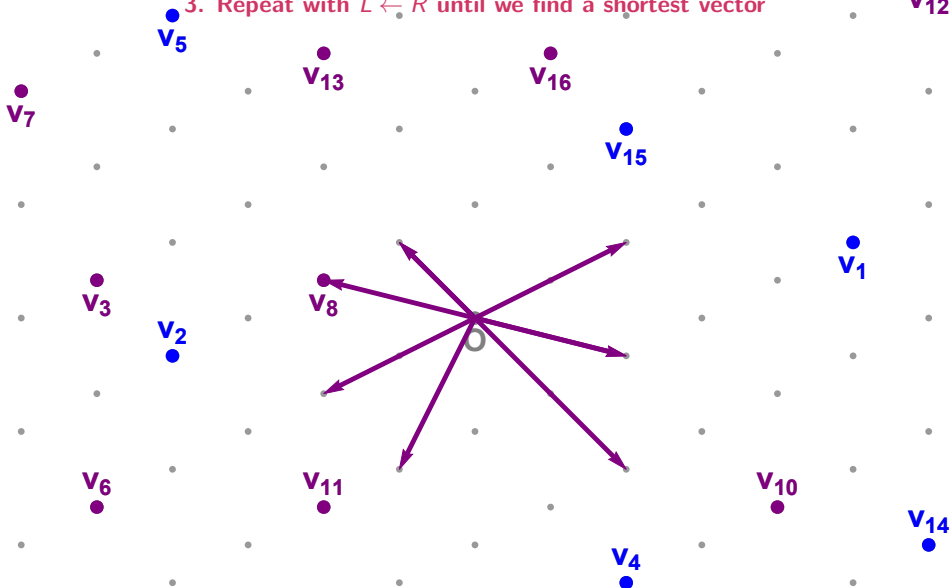
Nguyen-Vidick sieve

3. Repeat with $L \leftarrow R$ until we find a shortest vector



Nguyen-Vidick sieve

3. Repeat with $L \leftarrow R$ until we find a shortest vector



Nguyen-Vidick sieve

3. Repeat with $L \leftarrow R$ until we find a shortest vector



Nguyen-Vidick sieve

Overview



Nguyen-Vidick sieve

Overview

Heuristic (Nguyen and Vidick, J. Math. Crypt. '08)

The Nguyen-Vidick sieve runs in time $(4/3)^n$ and space $\sqrt{4/3}^n$.

Nguyen-Vidick sieve

Overview

Heuristic (Nguyen and Vidick, J. Math. Crypt. '08)

The Nguyen-Vidick sieve runs in time $2^{0.415n}$ and space $2^{0.208n}$.

Two-level sieve

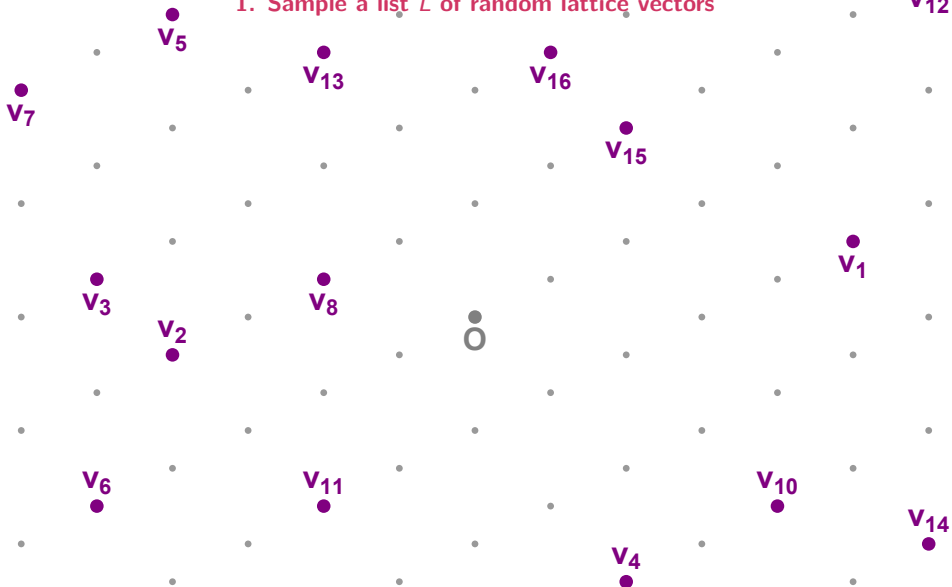
1. Sample a list L of random lattice vectors



O

Two-level sieve

1. Sample a list L of random lattice vectors



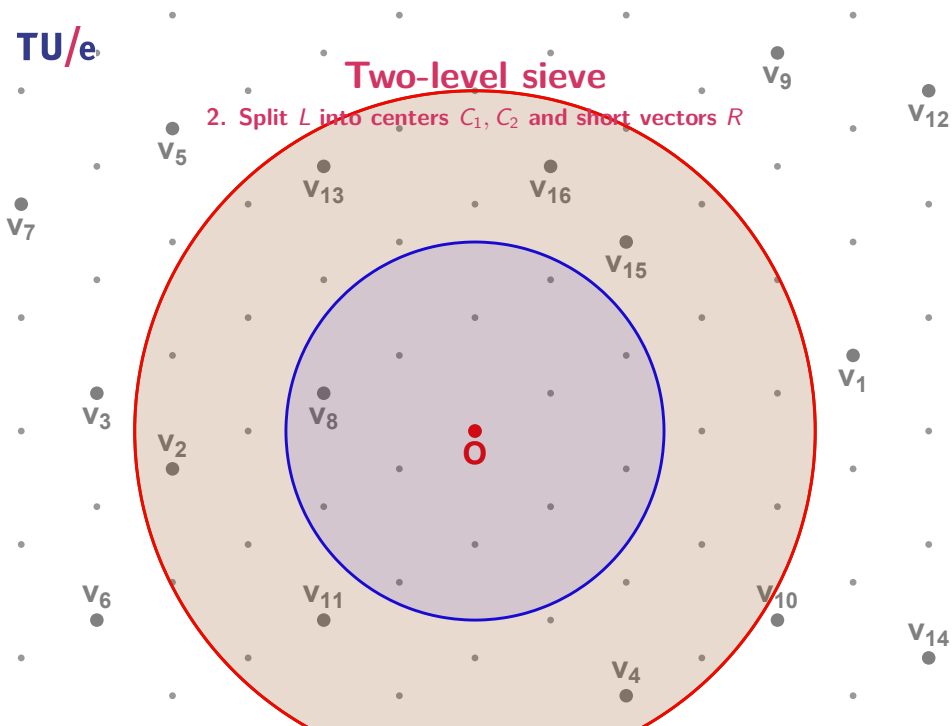
Two-level sieve

2. Split L into centers C_1, C_2 and short vectors R



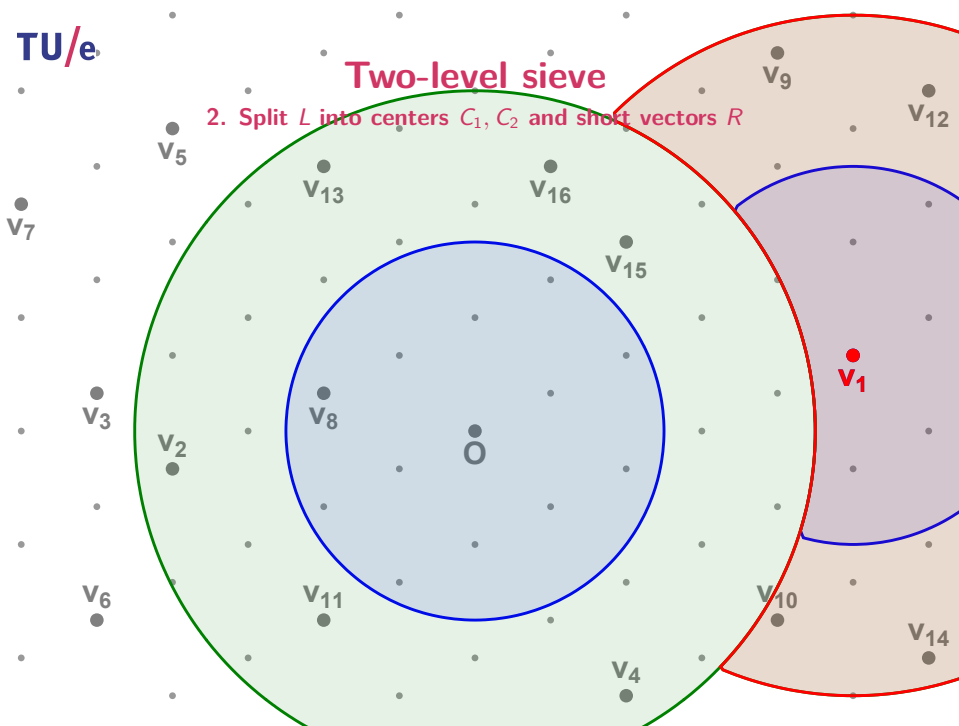
Two-level sieve

2. Split L into centers C_1, C_2 and short vectors R



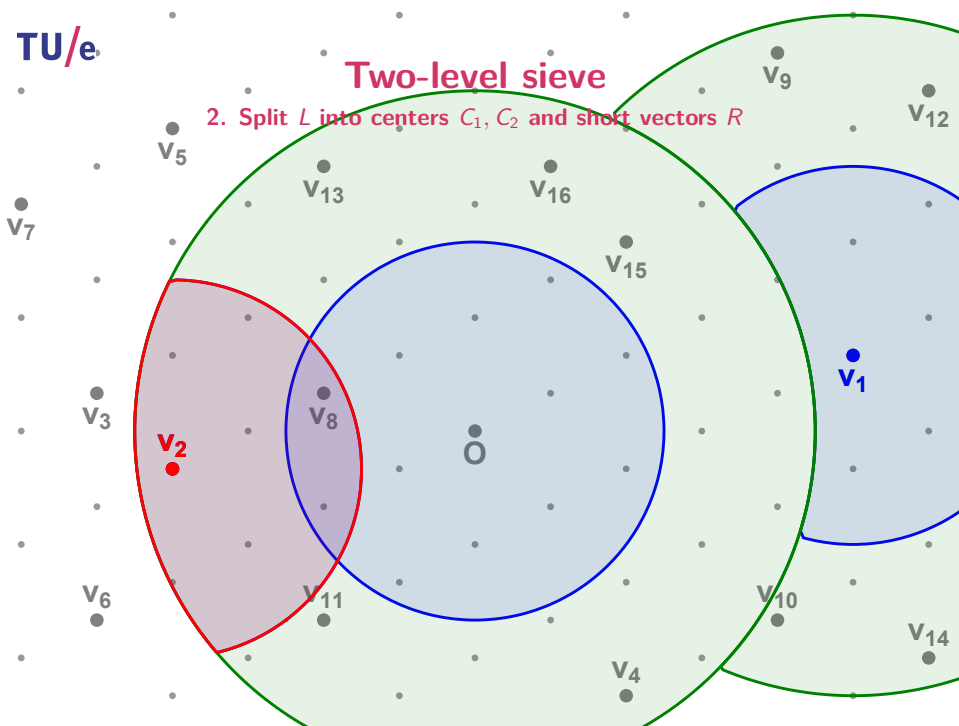
Two-level sieve

2. Split L into centers C_1, C_2 and short vectors R



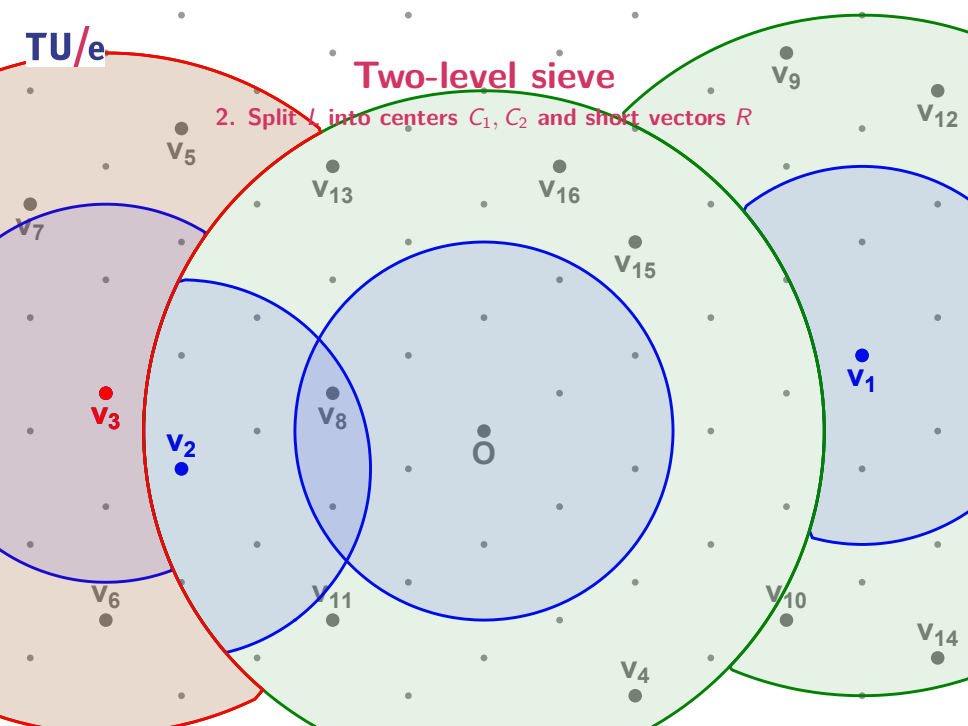
Two-level sieve

2. Split L into centers C_1, C_2 and short vectors R



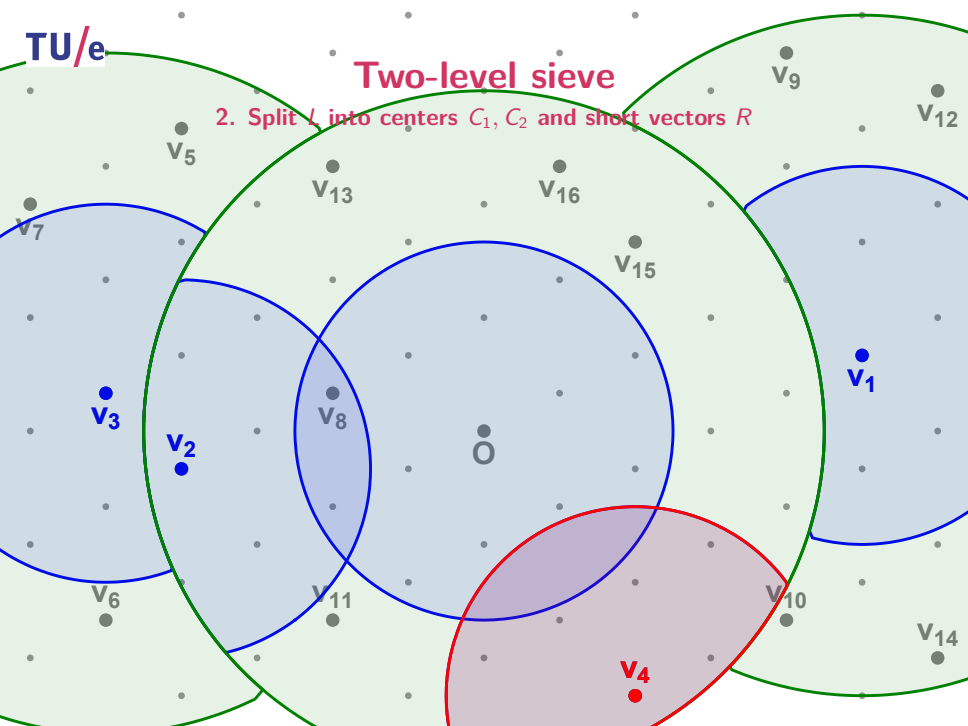
Two-level sieve

2. Split L into centers C_1, C_2 and short vectors R



Two-level sieve

2. Split L into centers C_1, C_2 and short vectors R



TU/e

Two-level sieve

2. Split L into centers C_1, C_2 and short vectors R

TU/e

Two-level sieve

2. Split L into centers C_1, C_2 and short vectors R

The diagram shows a central point O surrounded by several overlapping circles. A red circle labeled v_5 is highlighted, and a green circle labeled v_9 is also highlighted. Other circles are labeled $v_1, v_2, v_3, v_4, v_6, v_7, v_8, v_{10}, v_{11}, v_{12}, v_{13}, v_{14}, v_{15}, v_{16}$. The circles are colored in shades of blue and green.

TU/e

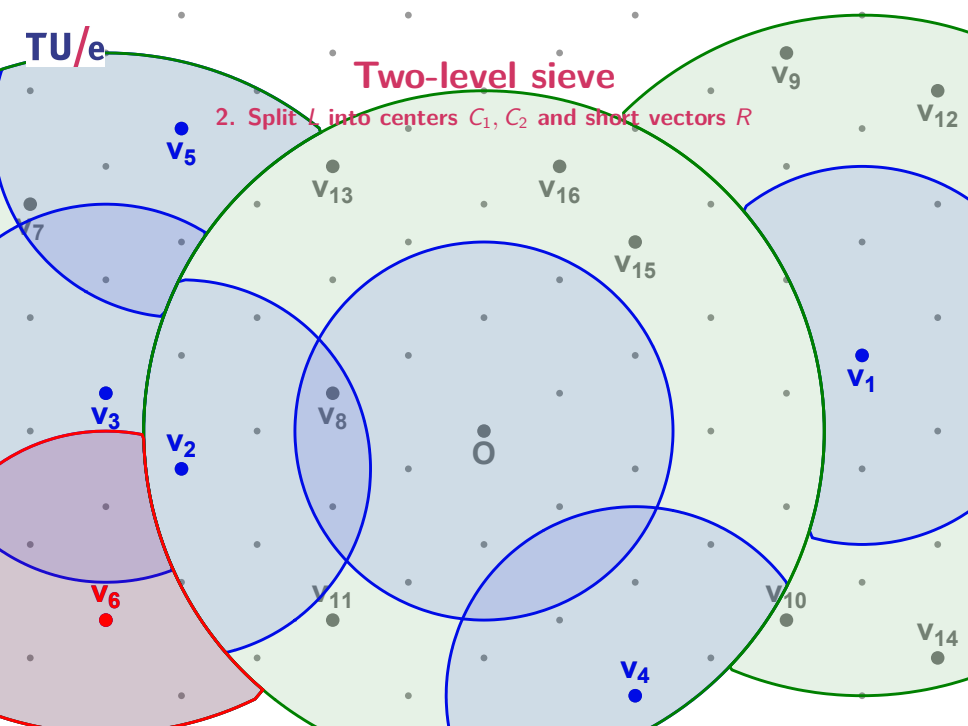
Two-level sieve

2. Split L into centers C_1, C_2 and short vectors R

Diagram illustrating the Two-level sieve method. A central point O is surrounded by several overlapping circles. A large green circle is labeled **TU/e** and contains a red dot labeled v_5 . A blue circle is labeled **2. Split L into centers C_1, C_2 and short vectors R** . Other circles are labeled with vectors $v_1, v_2, v_3, v_4, v_6, v_7, v_8, v_9, v_{10}, v_{11}, v_{12}, v_{13}, v_{14}, v_{15}, v_{16}$. The diagram shows the decomposition of a lattice L into smaller regions defined by centers C_1, C_2 and short vectors R .

Two-level sieve

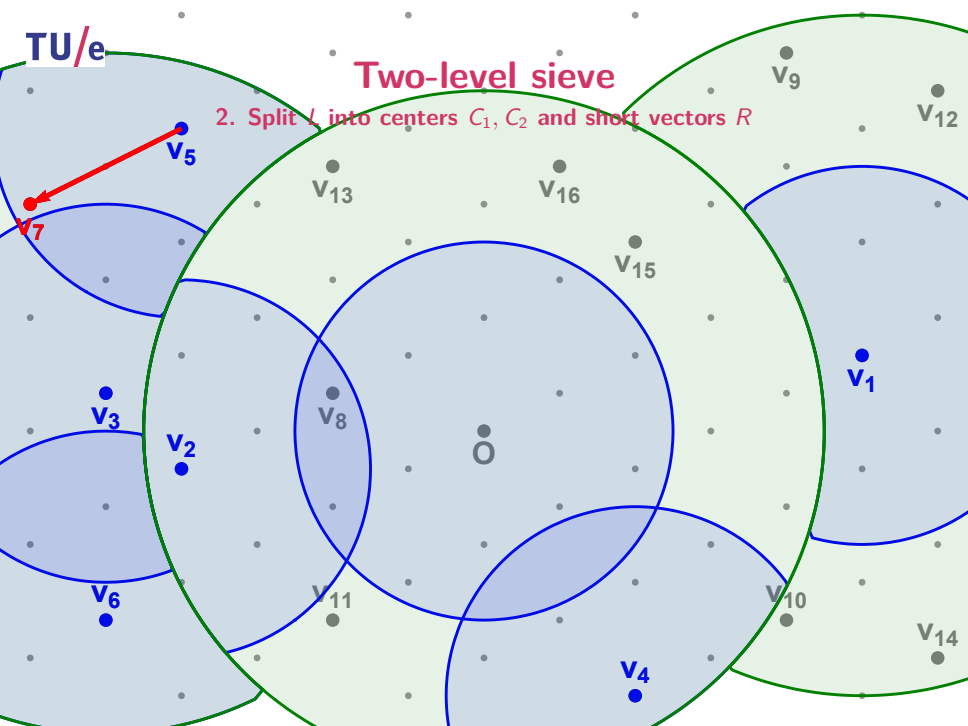
2. Split L into centers C_1, C_2 and short vectors R



TU/e

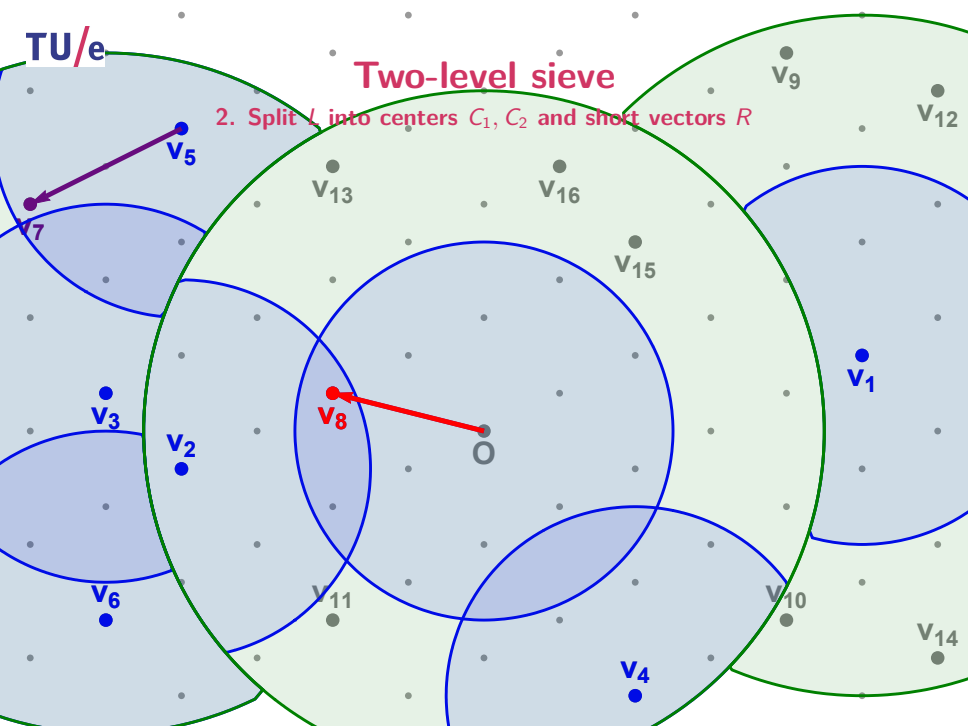
Two-level sieve

2. Split L into centers C_1, C_2 and short vectors R



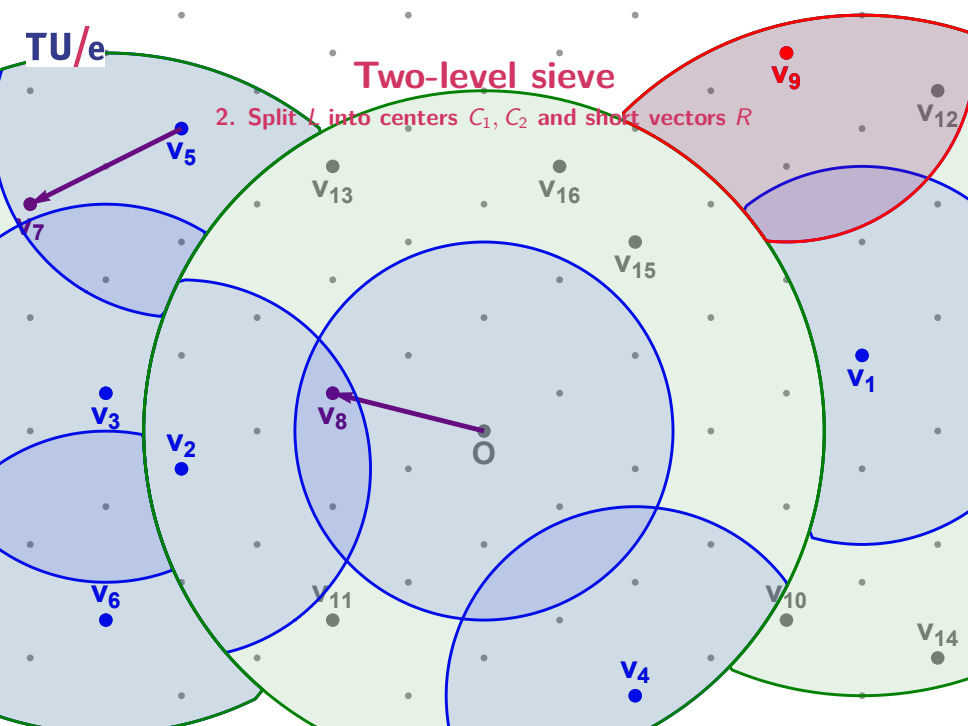
Two-level sieve

2. Split L into centers C_1, C_2 and short vectors R



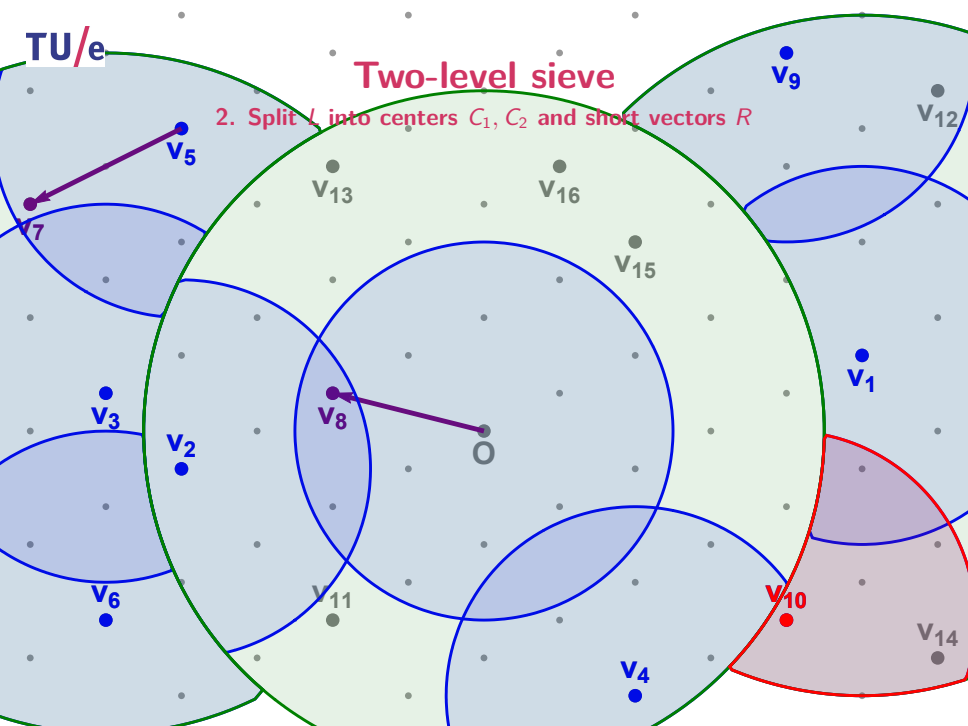
Two-level sieve

2. Split L into centers C_1, C_2 and short vectors R



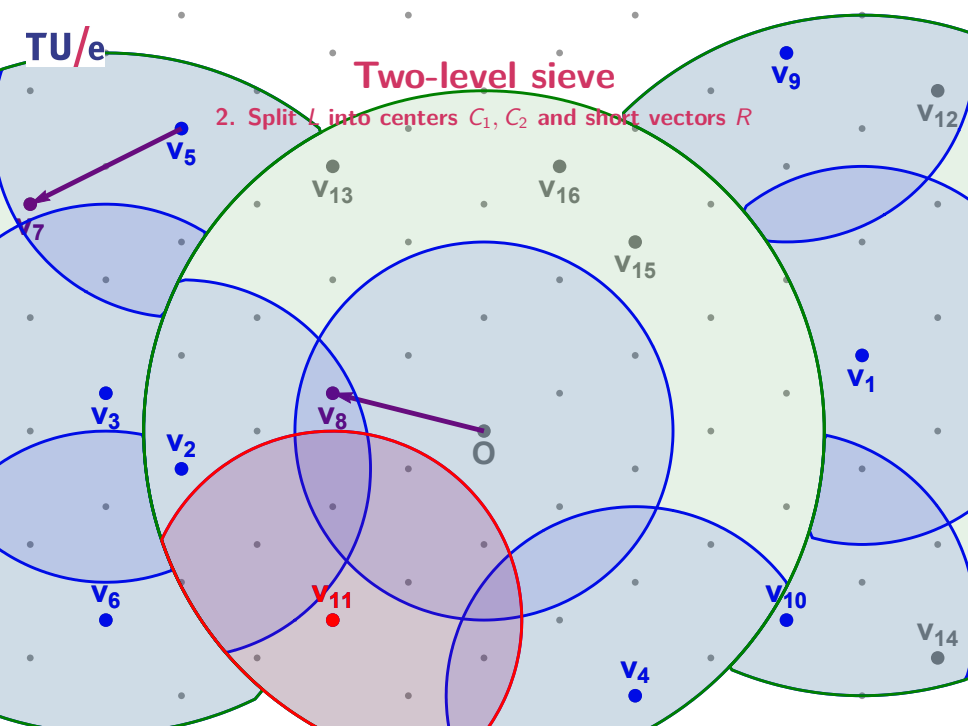
Two-level sieve

2. Split L into centers C_1, C_2 and short vectors R



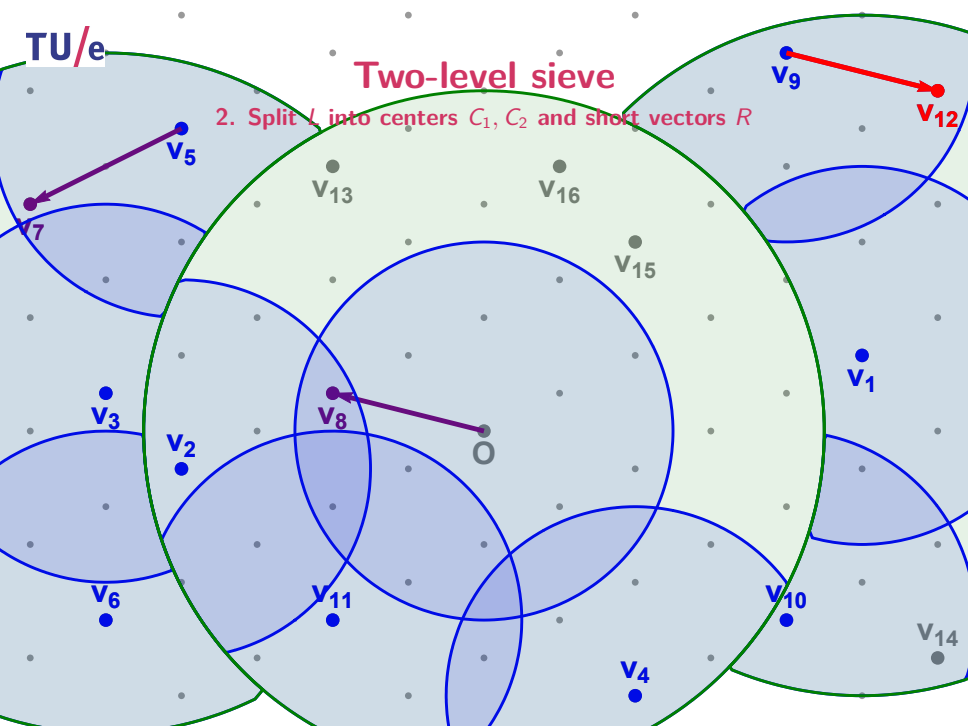
Two-level sieve

2. Split L into centers C_1, C_2 and short vectors R



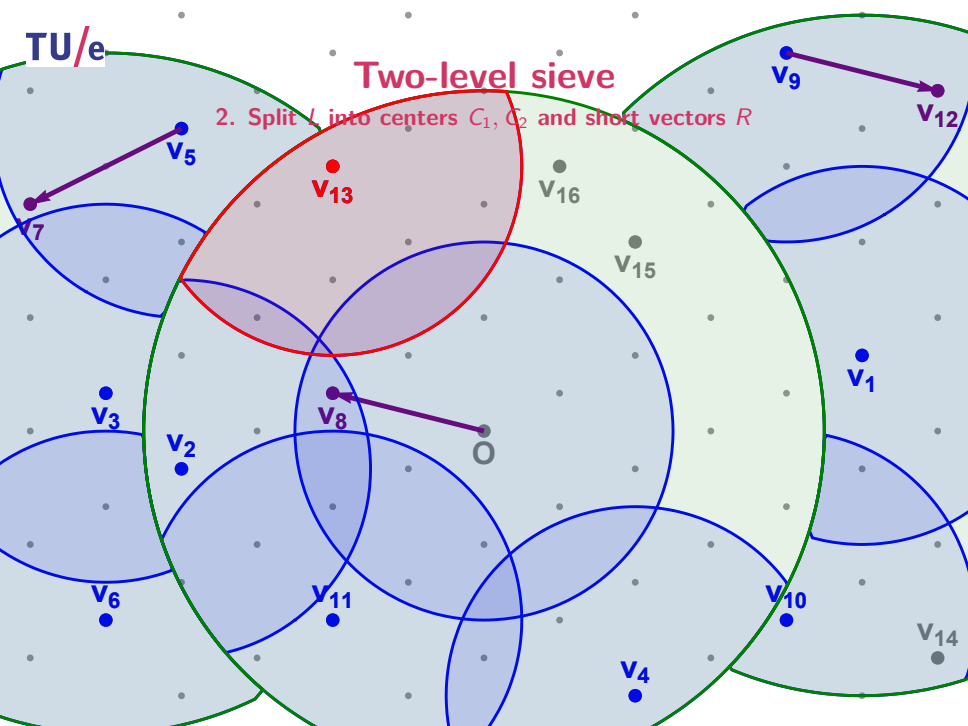
Two-level sieve

2. Split L into centers C_1, C_2 and short vectors R



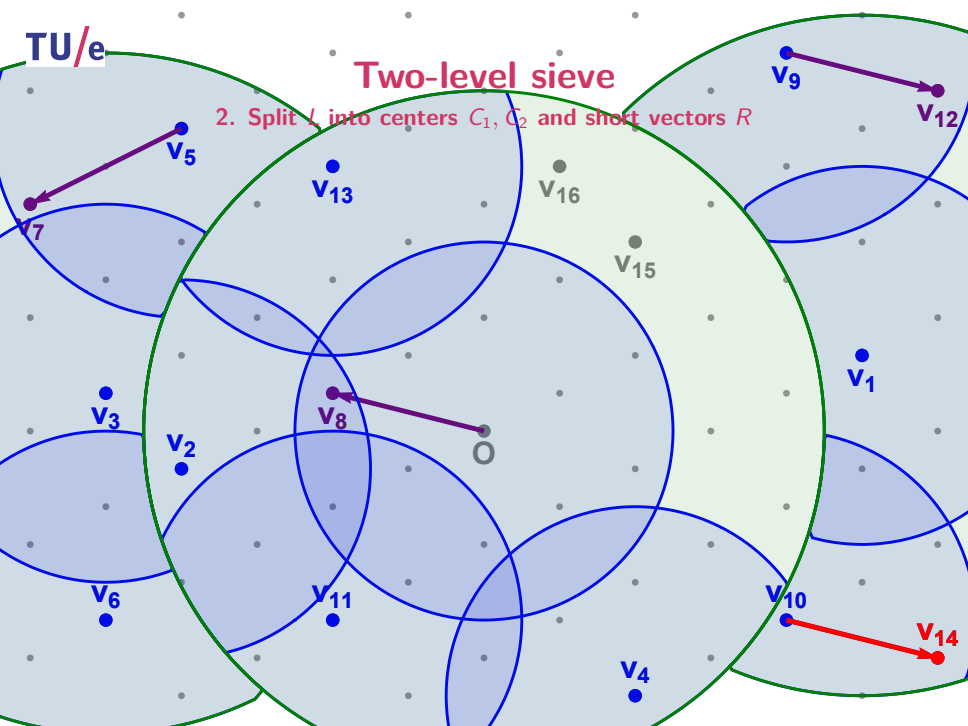
Two-level sieve

2. Split L into centers C_1, C_2 and short vectors R



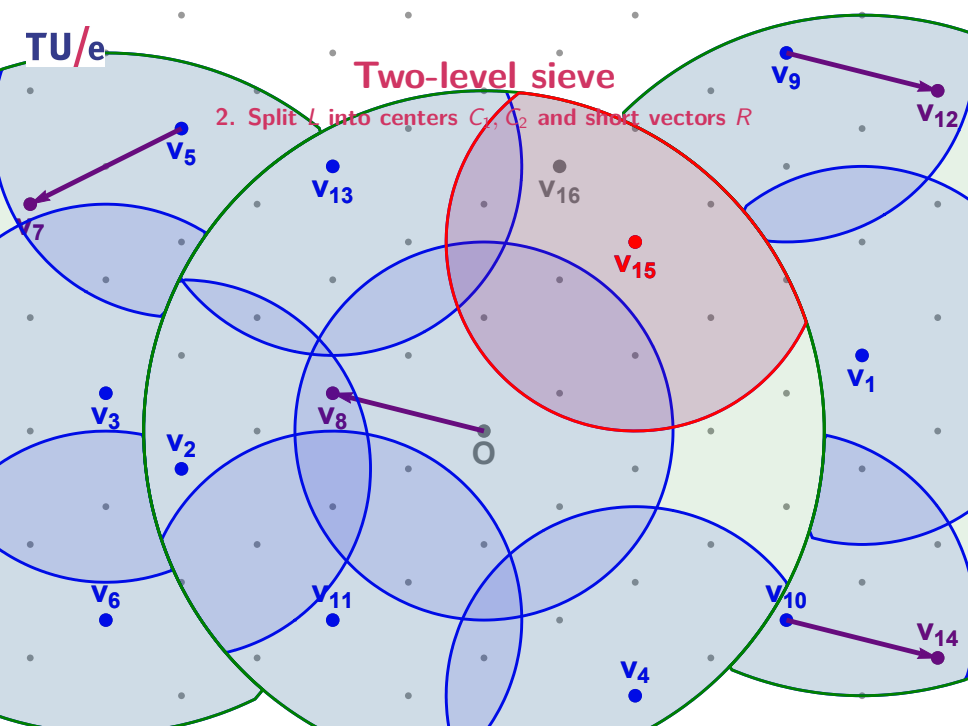
Two-level sieve

2. Split L into centers C_1, C_2 and short vectors R



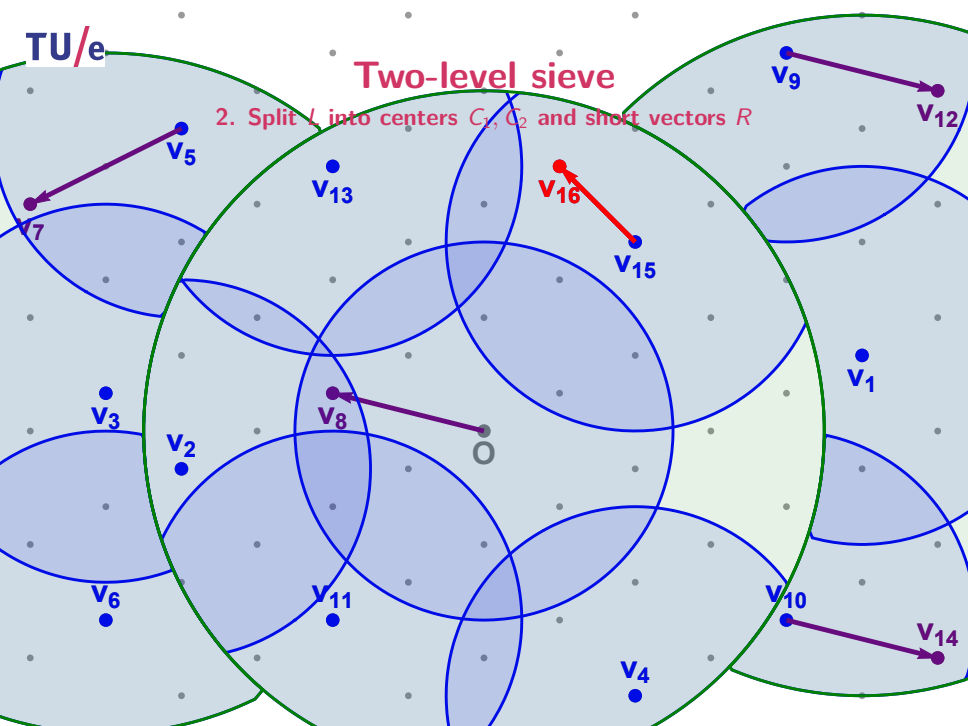
Two-level sieve

2. Split L into centers C_1, C_2 and short vectors R



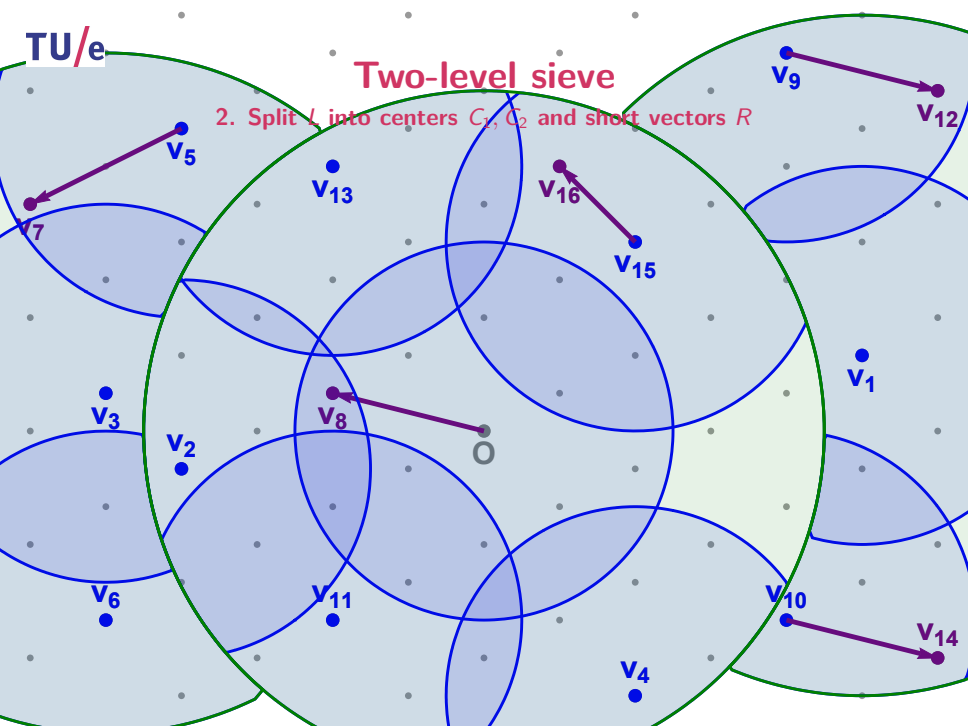
Two-level sieve

2. Split L into centers C_1, C_2 and short vectors R



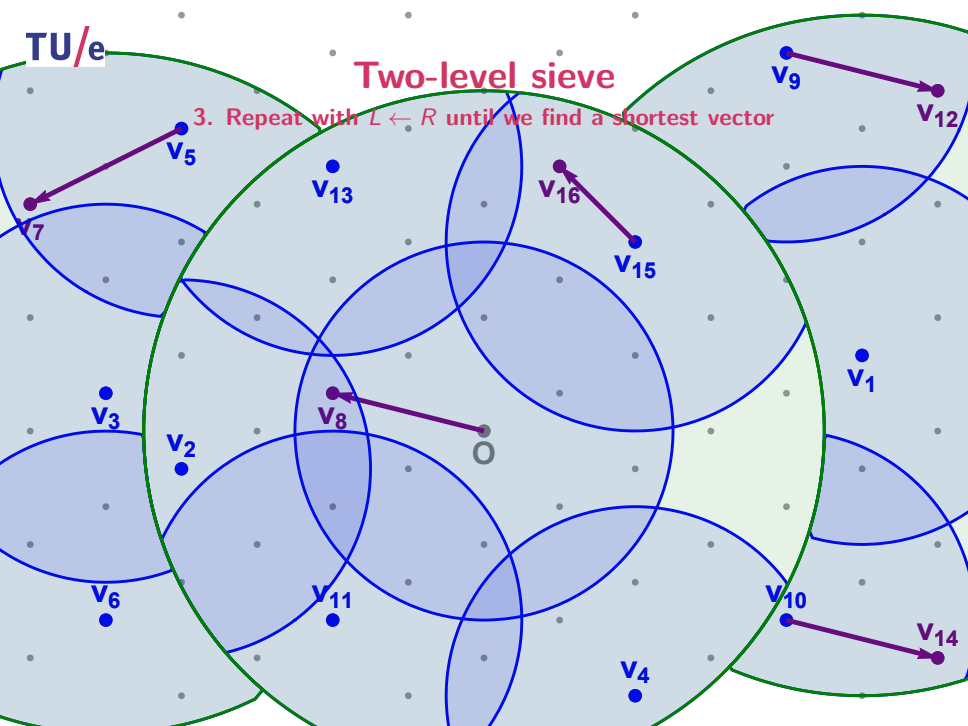
Two-level sieve

2. Split L into centers C_1, C_2 and short vectors R



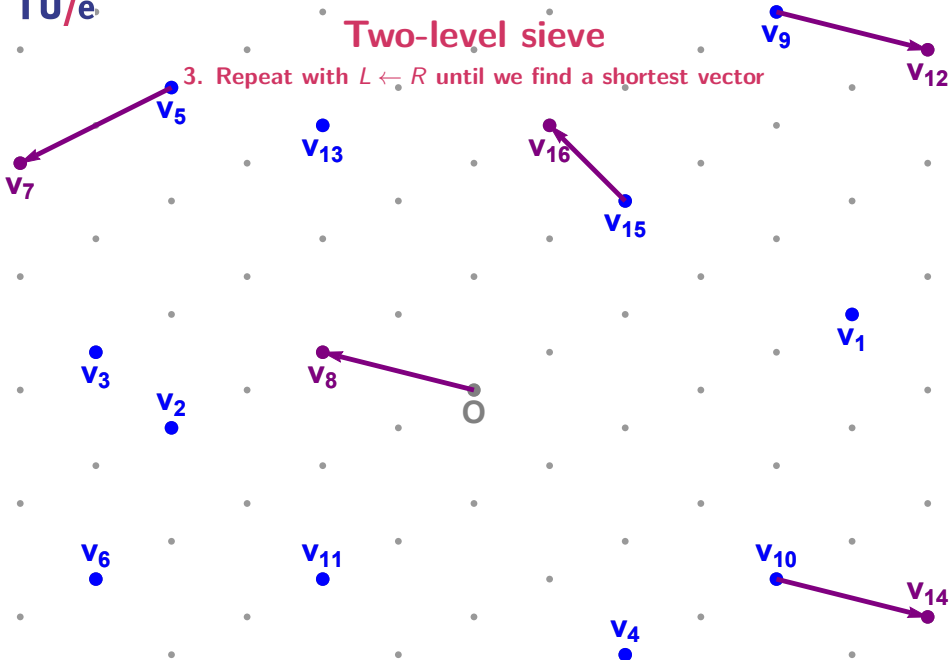
Two-level sieve

3. Repeat with $L \leftarrow R$ until we find a shortest vector



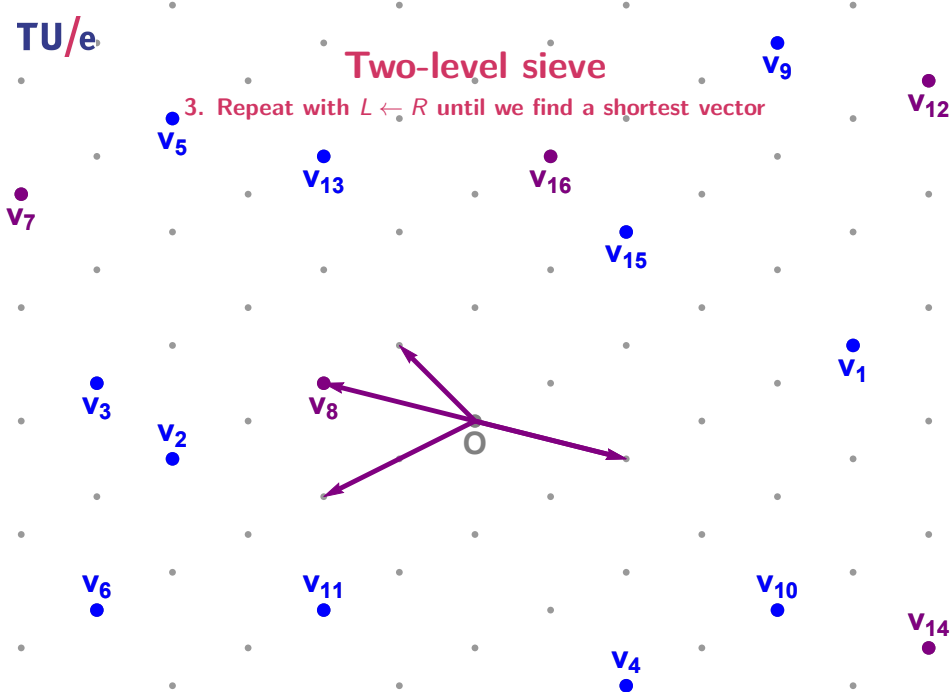
Two-level sieve

3. Repeat with $L \leftarrow R$ until we find a shortest vector



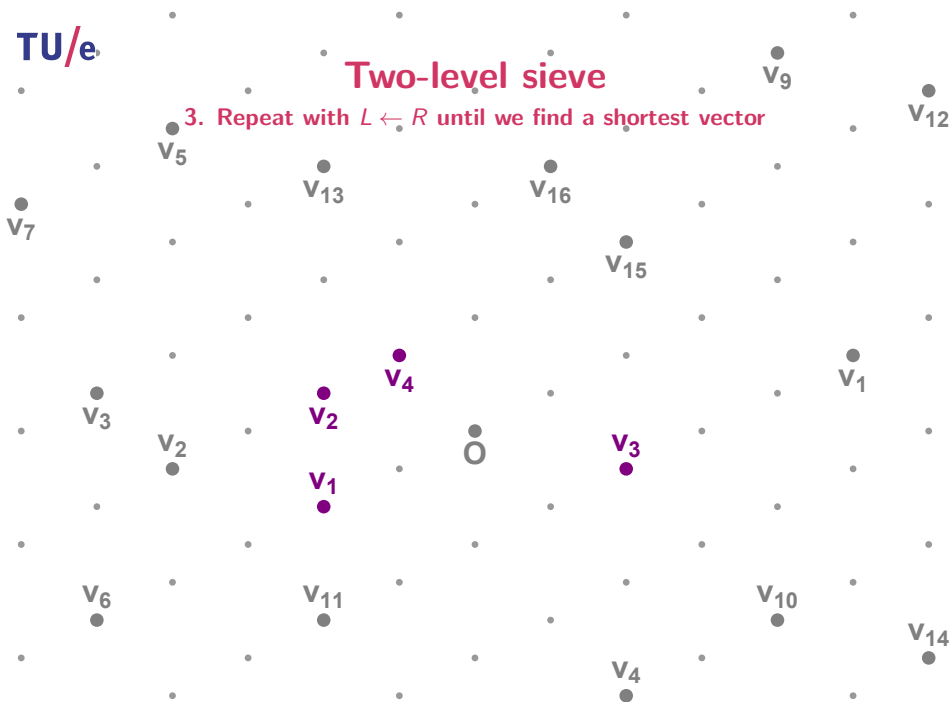
Two-level sieve

3. Repeat with $L \leftarrow R$ until we find a shortest vector



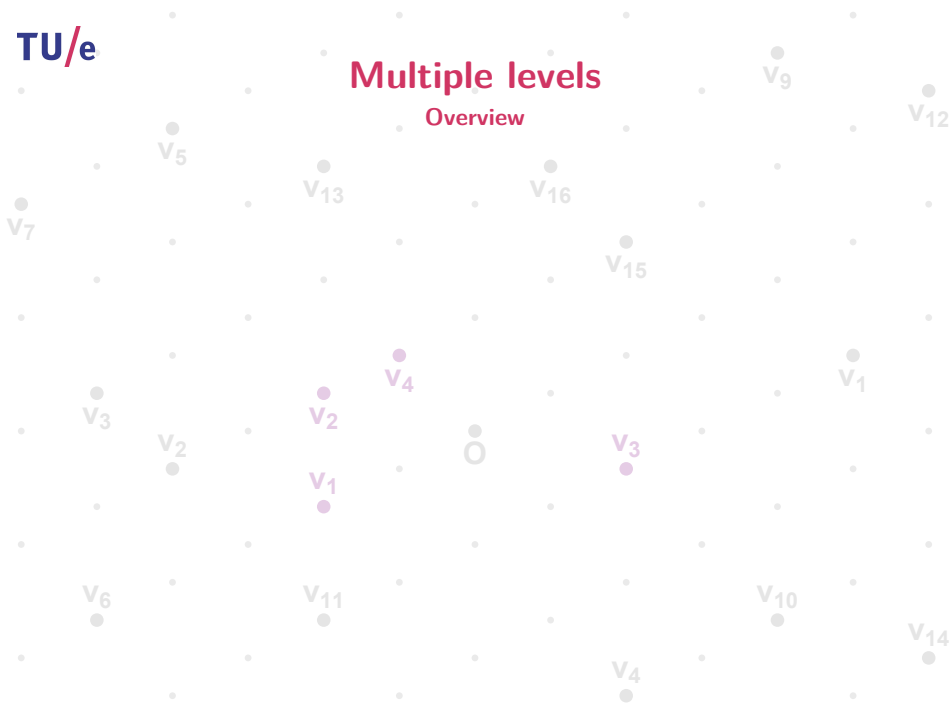
Two-level sieve

3. Repeat with $L \leftarrow R$ until we find a shortest vector



Multiple levels

Overview



Multiple levels

Overview

Heuristic (Nguyen and Vidick, J. Math. Crypt. '08)

The one-level sieve runs in time $2^{0.4150n}$ and space $2^{0.2075n}$.

Multiple levels

Overview

Heuristic (Nguyen and Vidick, J. Math. Crypt. '08)

The one-level sieve runs in time $2^{0.4150n}$ and space $2^{0.2075n}$.

Heuristic (Wang et al., ASIACCS'11)

The two-level sieve runs in time $2^{0.3836n}$ and space $2^{0.2557n}$.

Multiple levels

Overview

Heuristic (Nguyen and Vidick, J. Math. Crypt. '08)

The one-level sieve runs in time $2^{0.4150n}$ and space $2^{0.2075n}$.

Heuristic (Wang et al., ASIACCS'11)

The two-level sieve runs in time $2^{0.3836n}$ and space $2^{0.2557n}$.

Heuristic (Zhang et al., SAC'13)

The three-level sieve runs in time $2^{0.3778n}$ and space $2^{0.2833n}$.

Multiple levels

Overview

Heuristic (Nguyen and Vidick, J. Math. Crypt. '08)

The one-level sieve runs in time $2^{0.4150n}$ and space $2^{0.2075n}$.

Heuristic (Wang et al., ASIACCS'11)

The two-level sieve runs in time $2^{0.3836n}$ and space $2^{0.2557n}$.

Heuristic (Zhang et al., SAC'13)

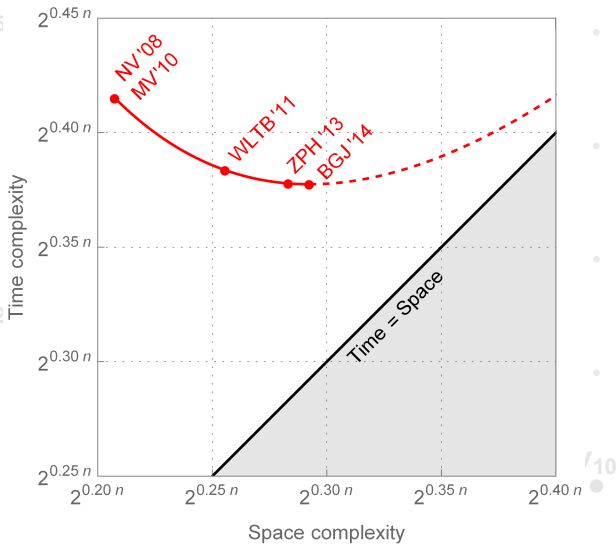
The three-level sieve runs in time $2^{0.3778n}$ and space $2^{0.2833n}$.

Conjecture

The four-level sieve runs in time $2^{0.3774n}$ and space $2^{0.2925n}$, and higher-level sieves are not faster than this.

Sieving

Space/time trade-off



Quantum Search

Classical form

Problem: Given a list L of size N , and a function $f : L \rightarrow \{0, 1\}$ such that there is exactly one element $e \in L$ with $f(e) = 1$. Find this element e .

- Classical search: $\Theta(N)$ time
- Quantum search: $\Theta(\sqrt{N})$ time [Gro96]

Quantum Search

General form

Problem: Given a list L of size N , and a function $f : L \rightarrow \{0, 1\}$ such that there are $c = O(1)$ elements $e \in L$ with $f(e) = 1$. Find one such element e .

- Classical search: $\Theta(N/c)$ time
- Quantum search: $\Theta(\sqrt{N/c})$ time [Gro96]

Quantum Search

General form

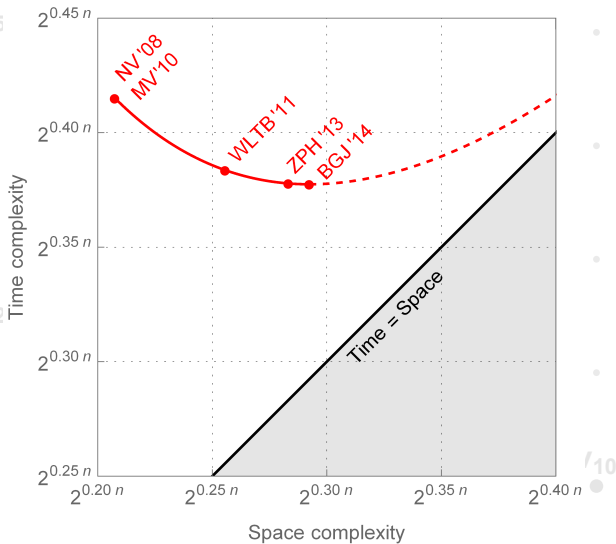
Problem: Given a list L of size N , and a function $f : L \rightarrow \{0, 1\}$ such that there are $c = O(1)$ elements $e \in L$ with $f(e) = 1$. Find one such element e .

- Classical search: $\Theta(N/c)$ time
- Quantum search: $\Theta(\sqrt{N/c})$ time [Gro96]

Potentially speed up search subroutines in sieving

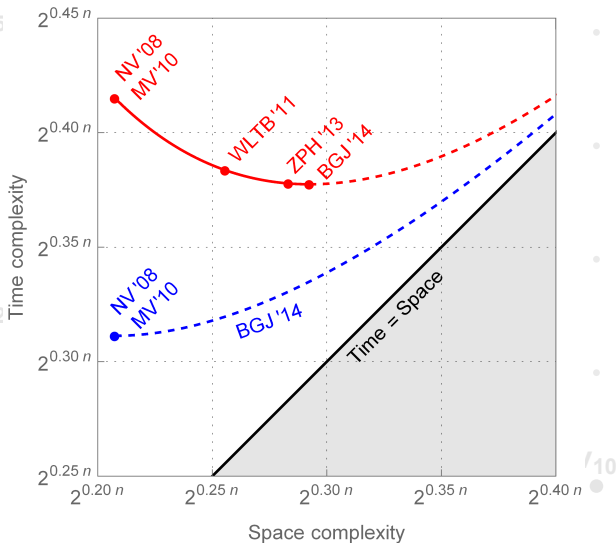
Sieving

Space/time trade-off



Sieving

Space/time trade-off



Nguyen-Vidick sieve with LSH

1. Sample a list L of random lattice vectors



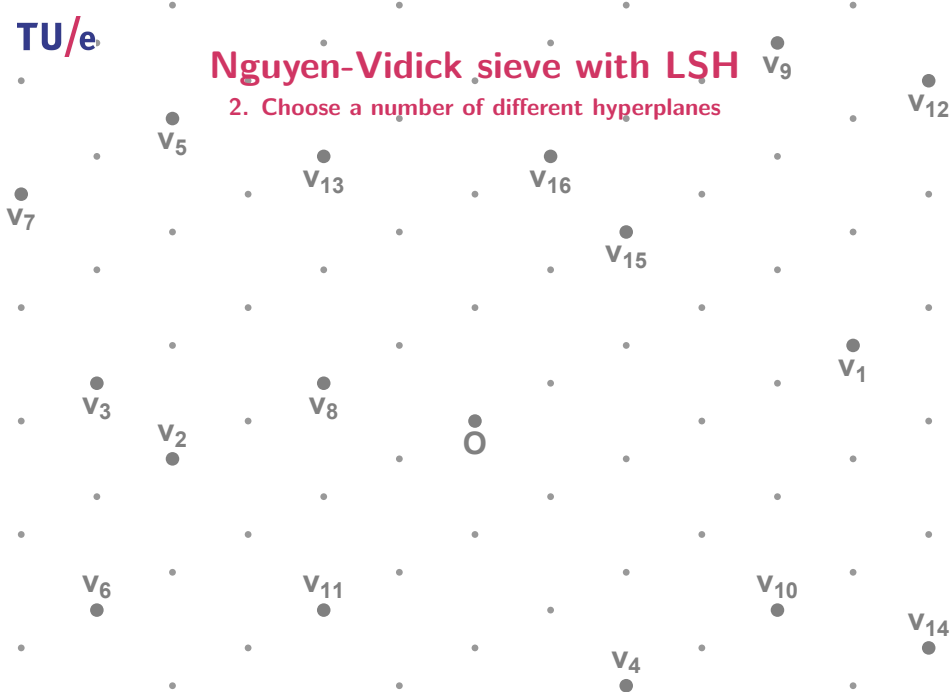
Nguyen-Vidick sieve with LSH

1. Sample a list L of random lattice vectors



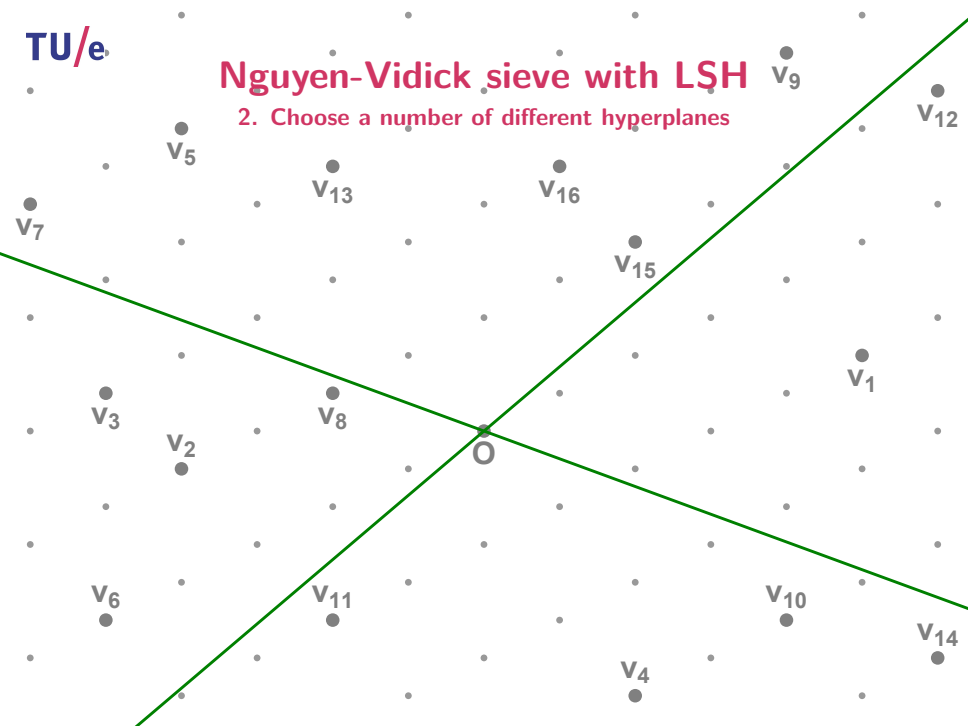
Nguyen-Vidick sieve with LSH

2. Choose a number of different hyperplanes



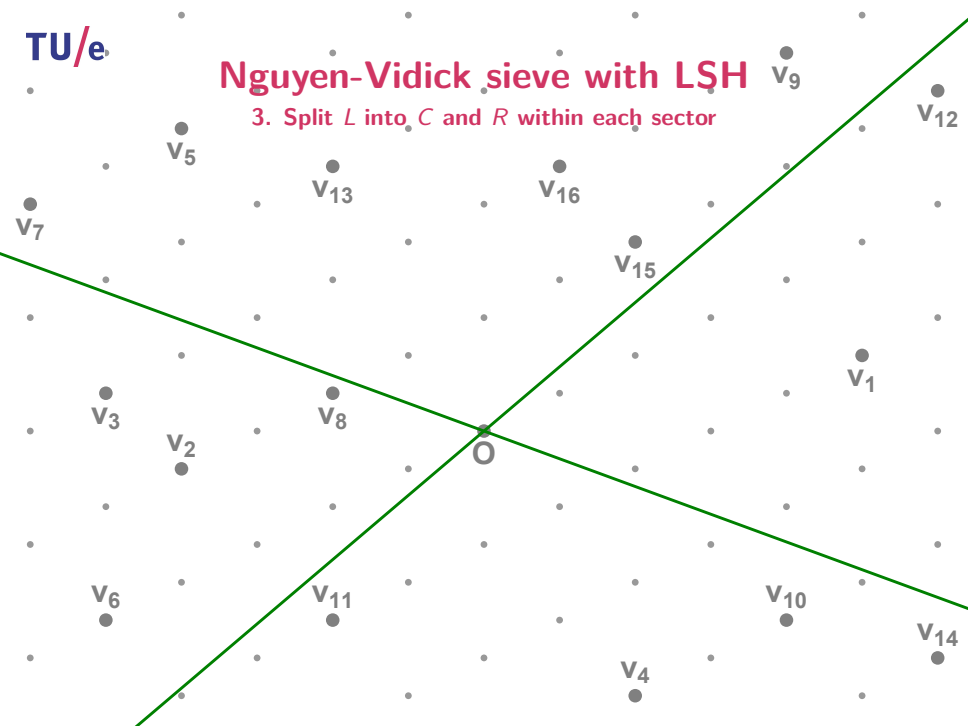
Nguyen-Vidick sieve with LSH

2. Choose a number of different hyperplanes



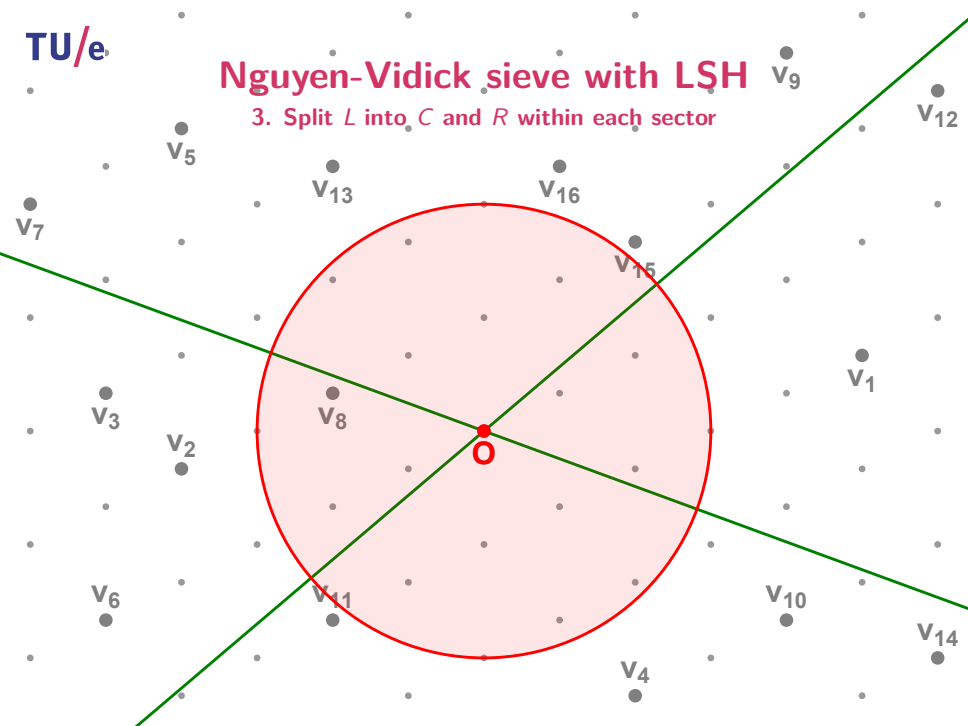
Nguyen-Vidick sieve with LSH

3. Split L into C and R within each sector



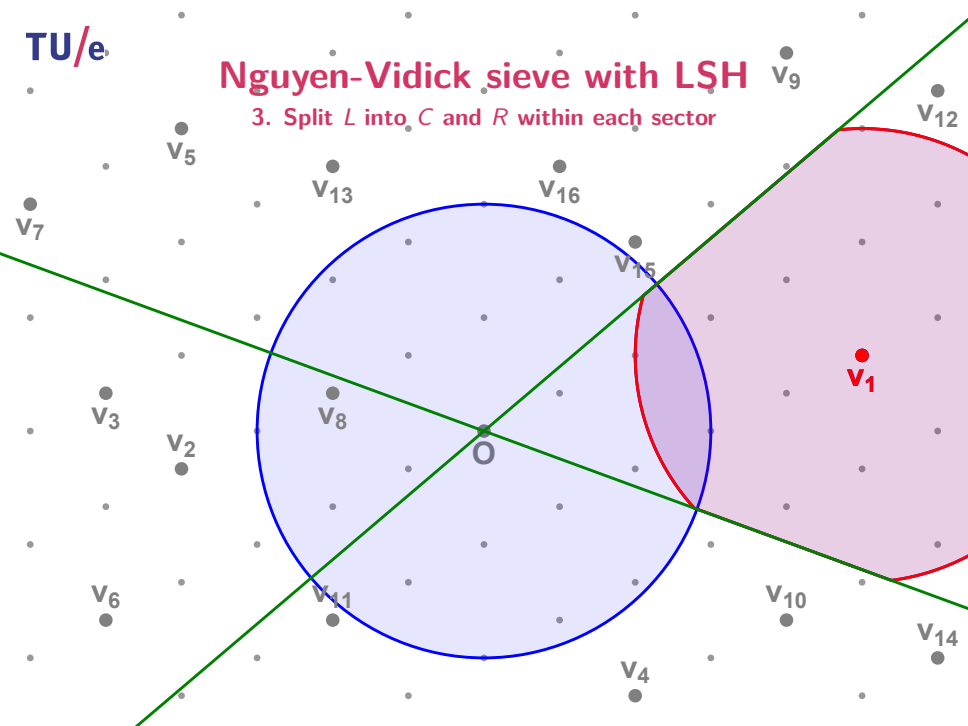
Nguyen-Vidick sieve with LSH

3. Split L into C and R within each sector



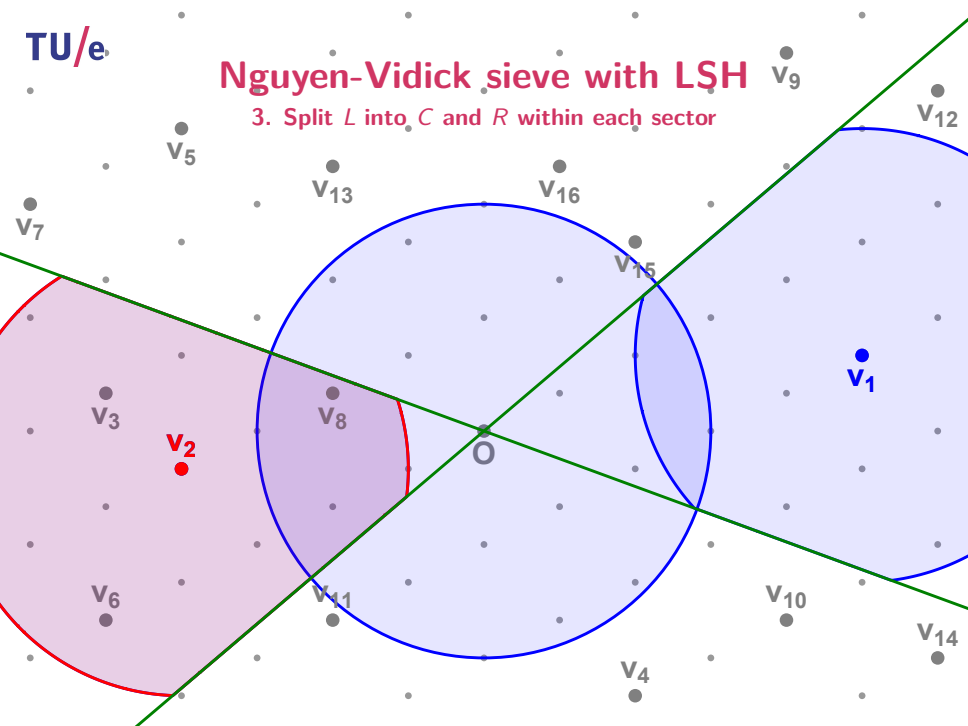
Nguyen-Vidick sieve with LSH

3. Split L into C and R within each sector



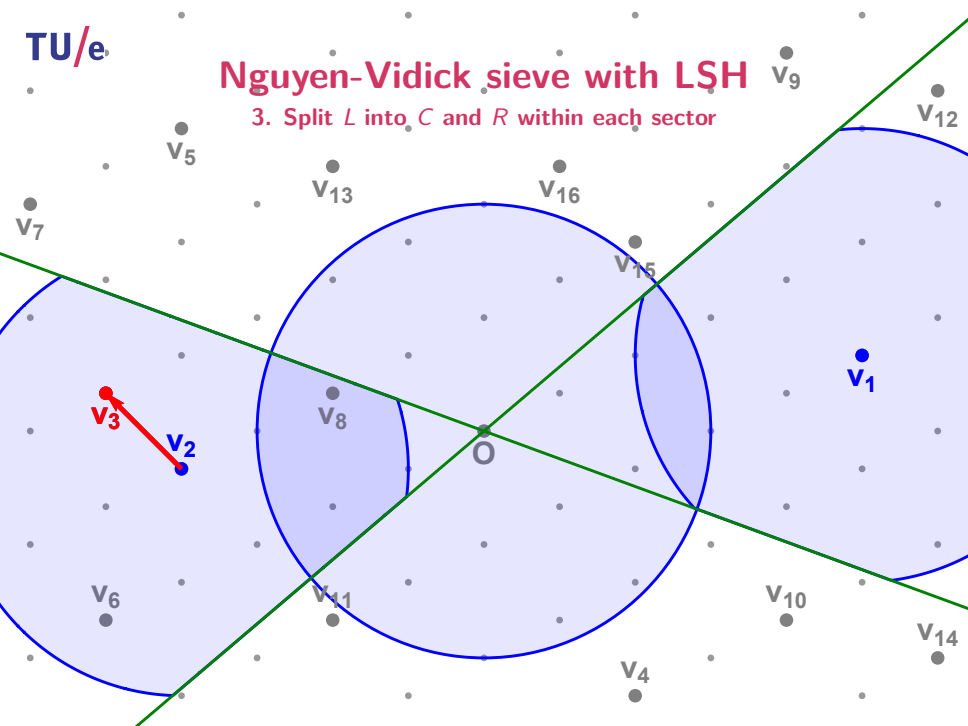
Nguyen-Vidick sieve with LSH

3. Split L into C and R within each sector



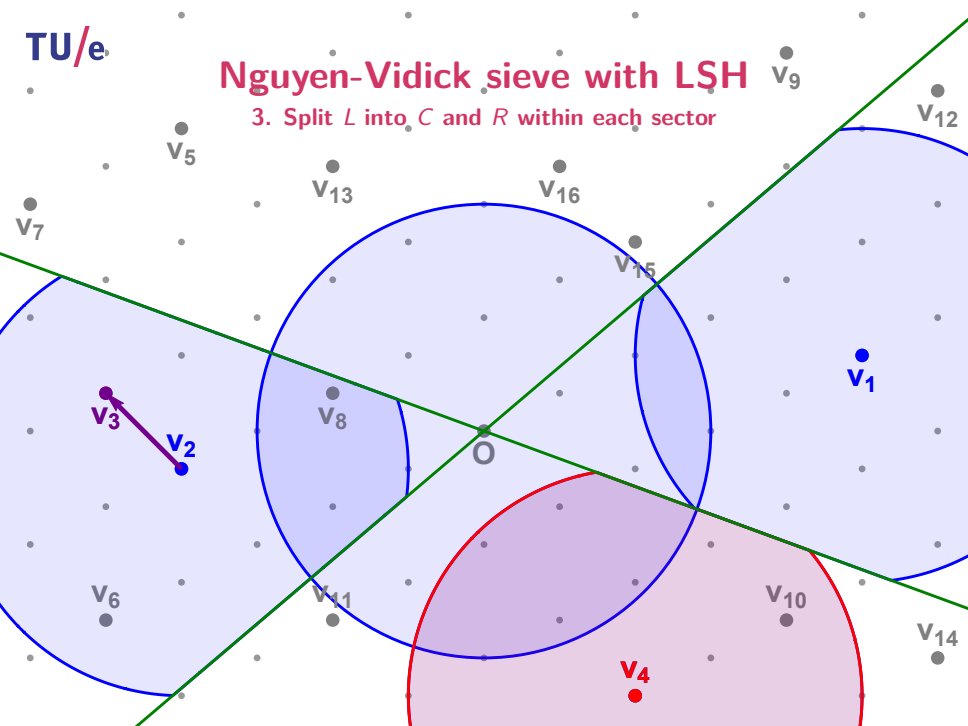
Nguyen-Vidick sieve with LSH

3. Split L into C and R within each sector



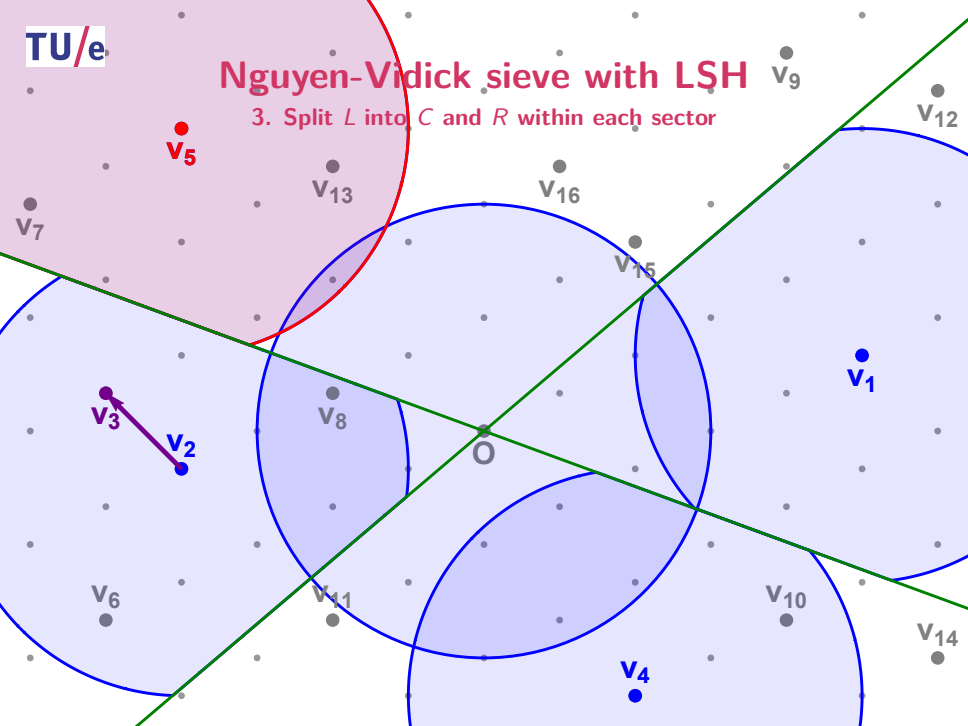
Nguyen-Vidick sieve with LSH

3. Split L into C and R within each sector



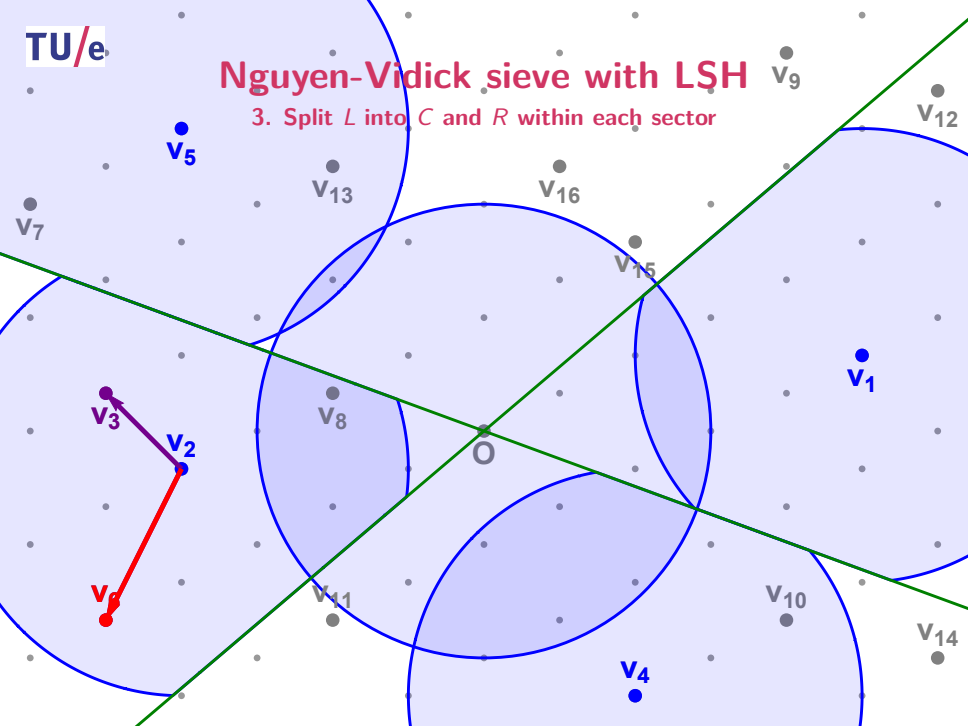
Nguyen-Vidick sieve with LSH

3. Split L into C and R within each sector



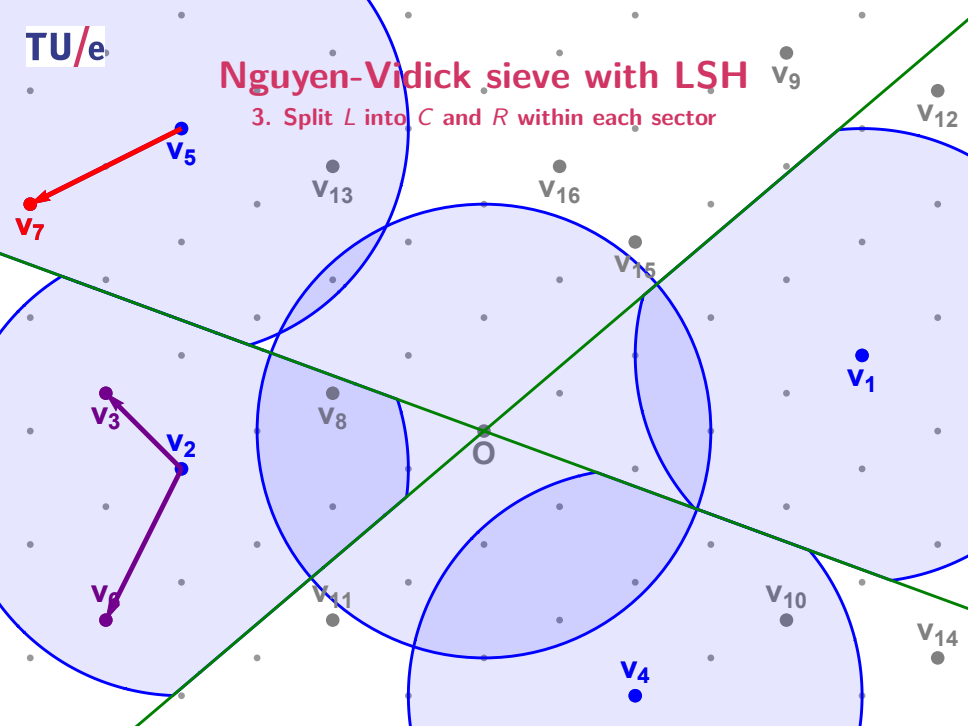
Nguyen-Vidick sieve with LSH

3. Split L into C and R within each sector



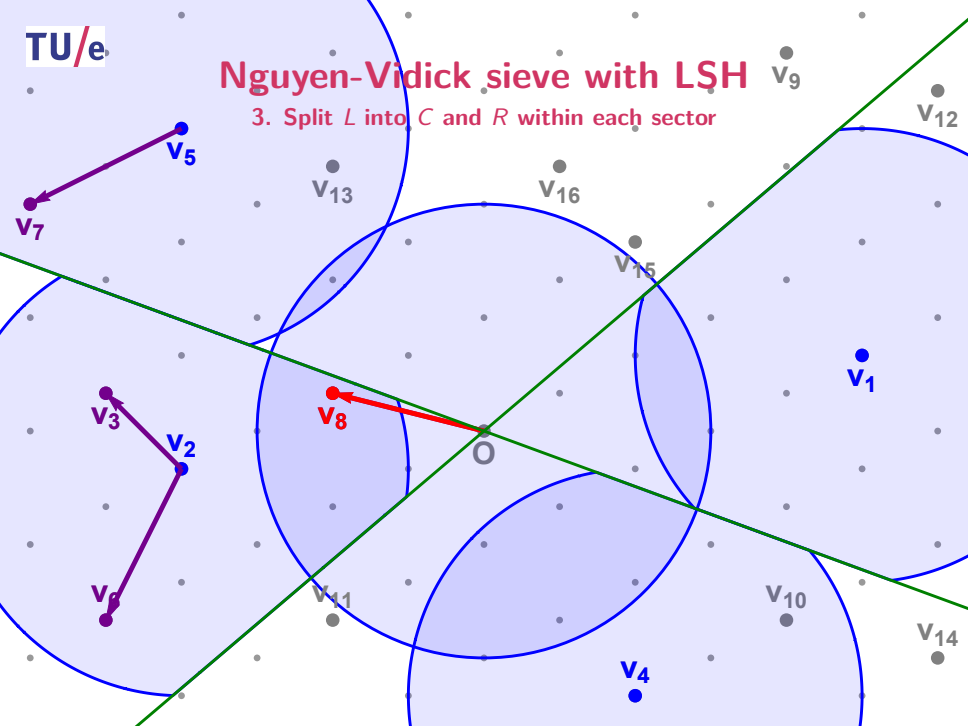
Nguyen-Vidick sieve with LSH

3. Split L into C and R within each sector



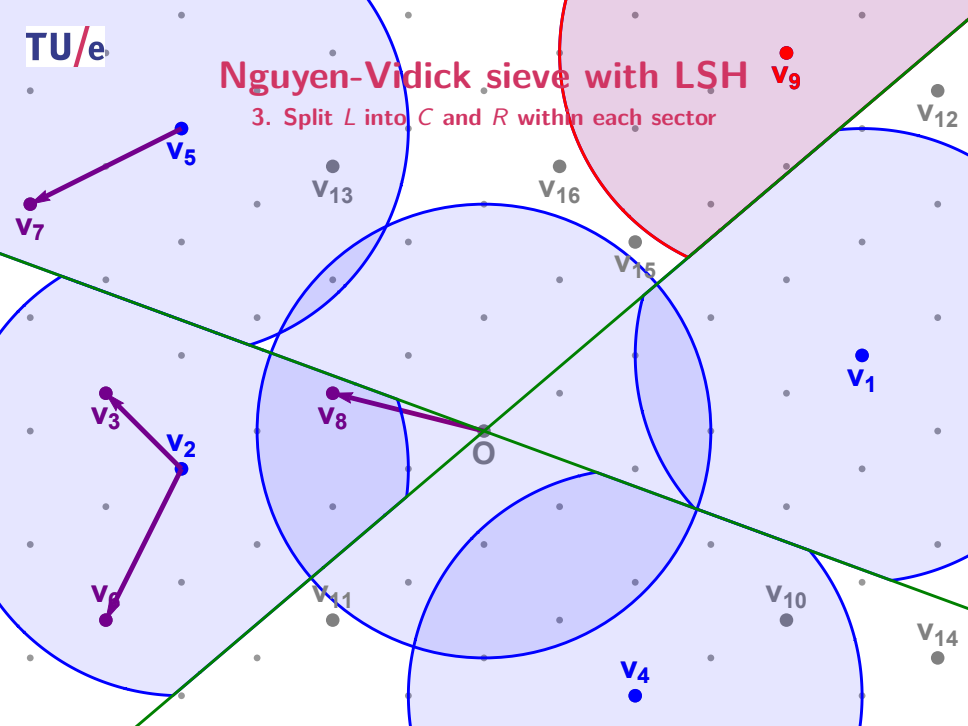
Nguyen-Vidick sieve with LSH

3. Split L into C and R within each sector



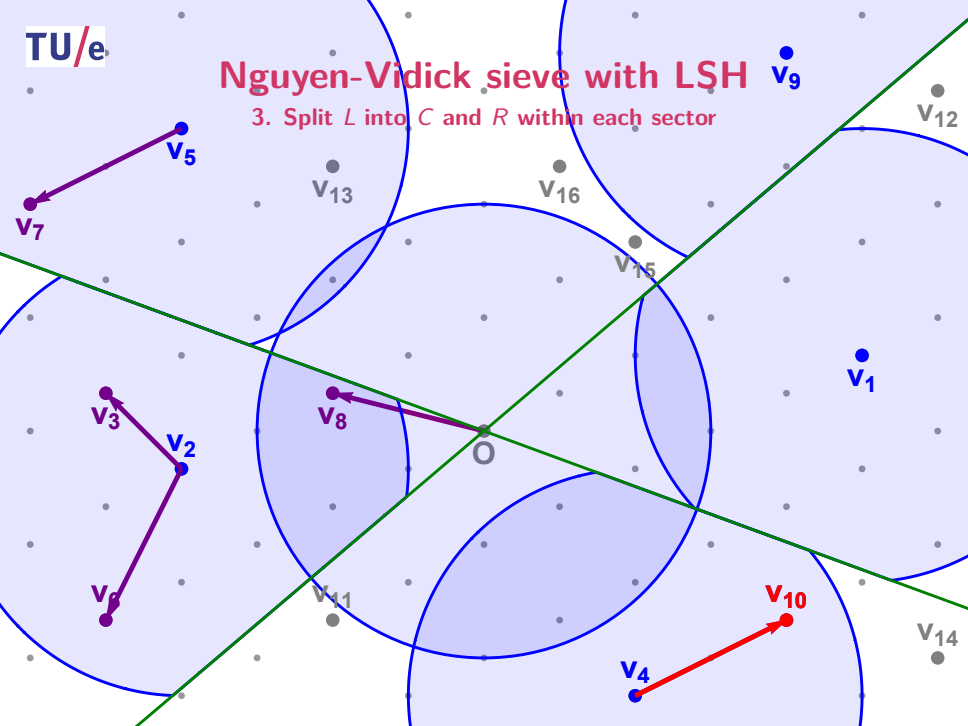
Nguyen-Vidick sieve with LSH

3. Split L into C and R within each sector



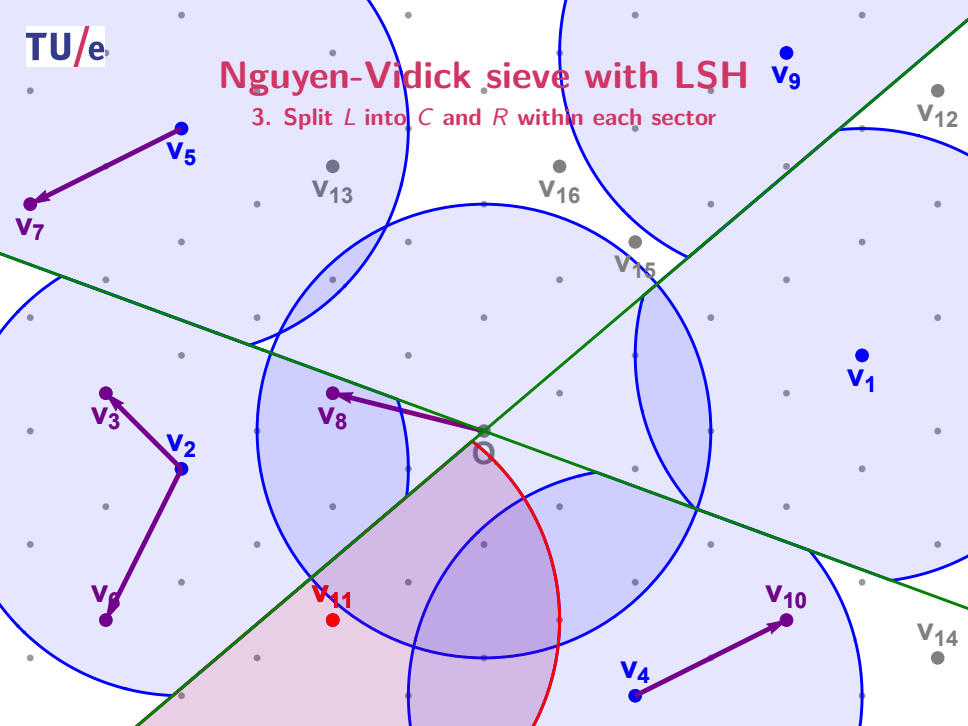
Nguyen-Vidick sieve with LSH

3. Split L into C and R within each sector



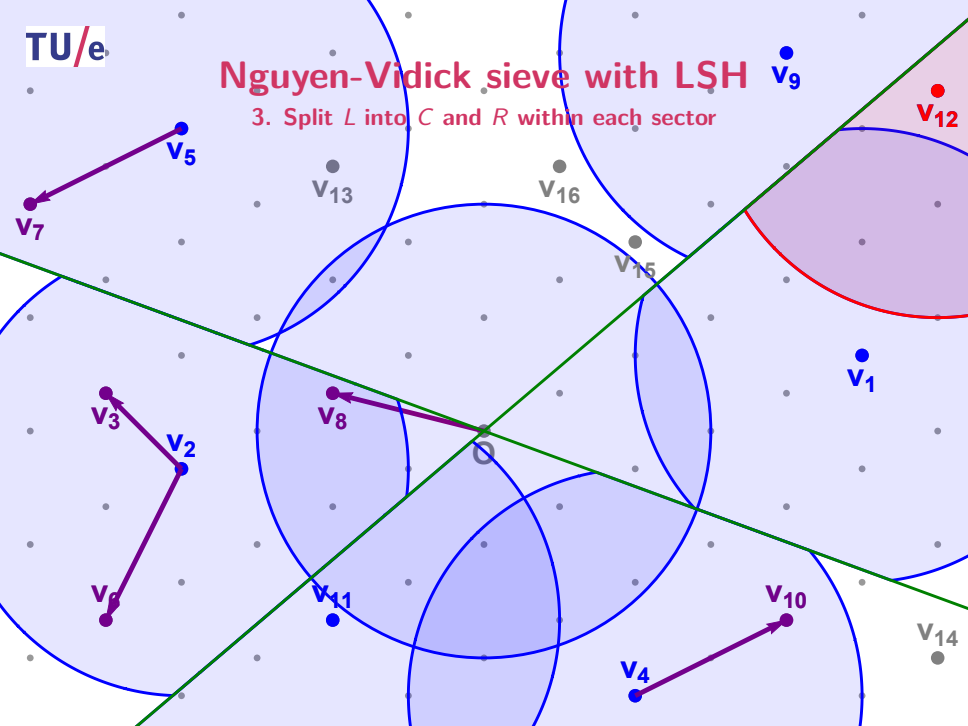
Nguyen-Vidick sieve with LSH

3. Split L into C and R within each sector



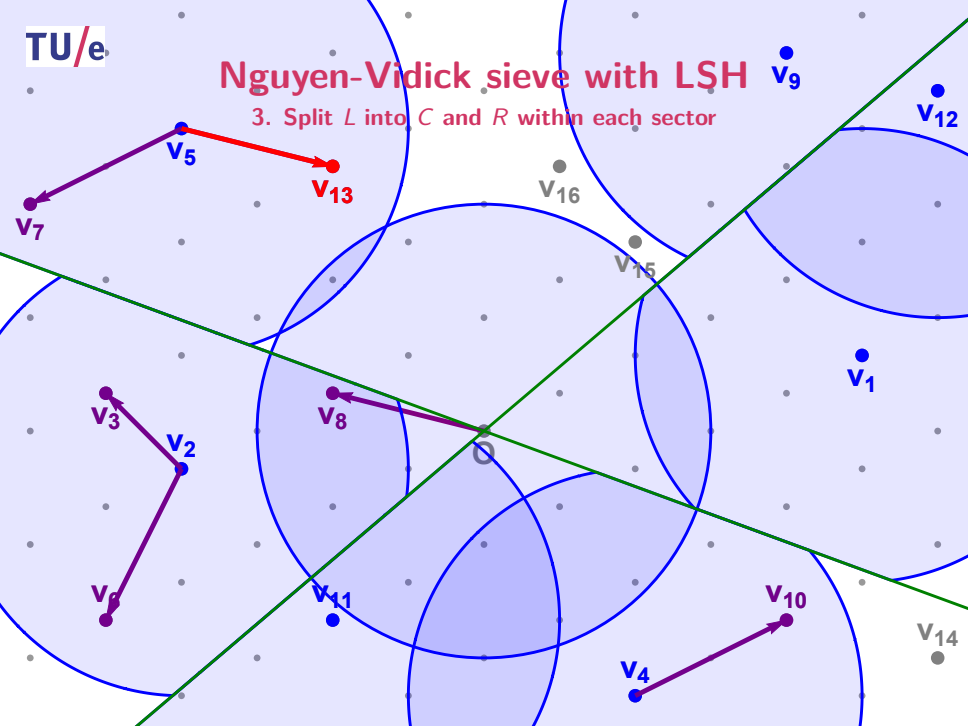
Nguyen-Vidick sieve with LSH

3. Split L into C and R within each sector



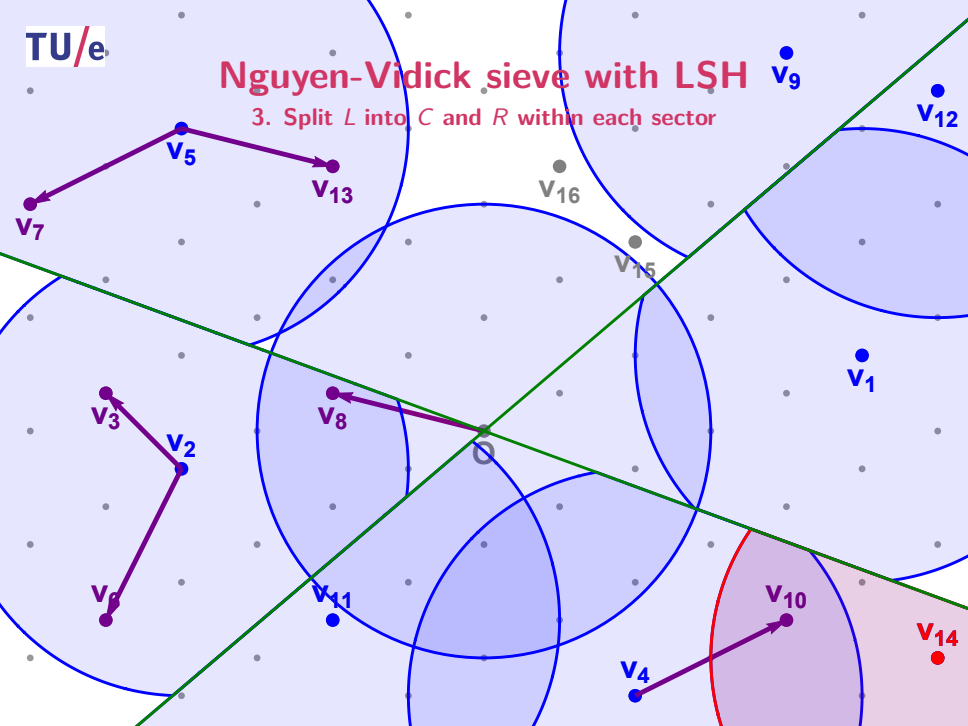
Nguyen-Vidick sieve with LSH

3. Split L into C and R within each sector



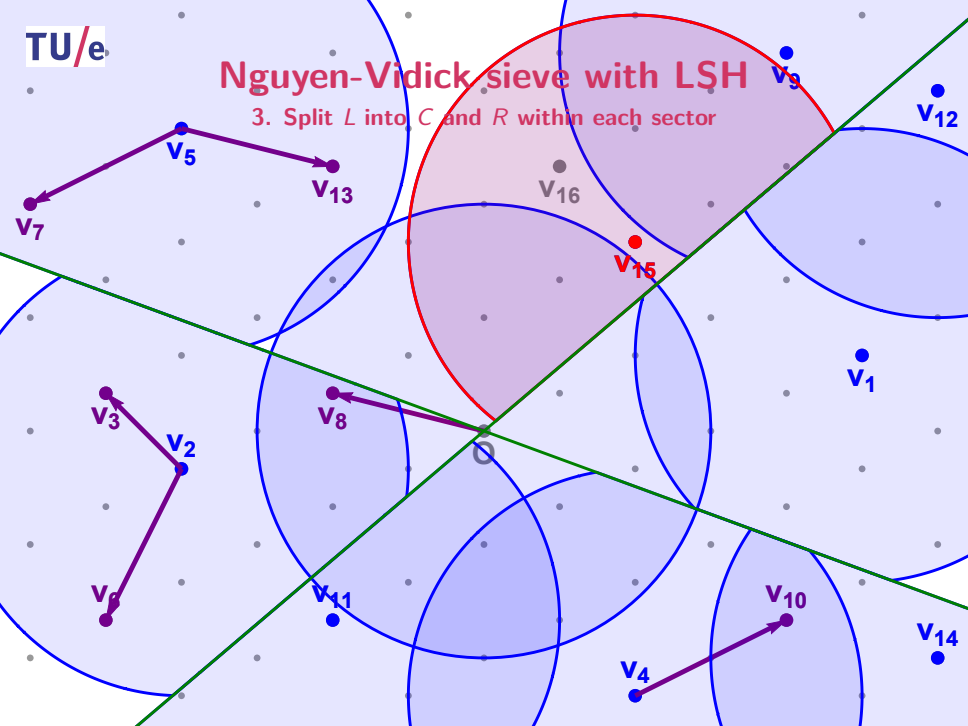
Nguyen-Vidick sieve with LSH

3. Split L into C and R within each sector



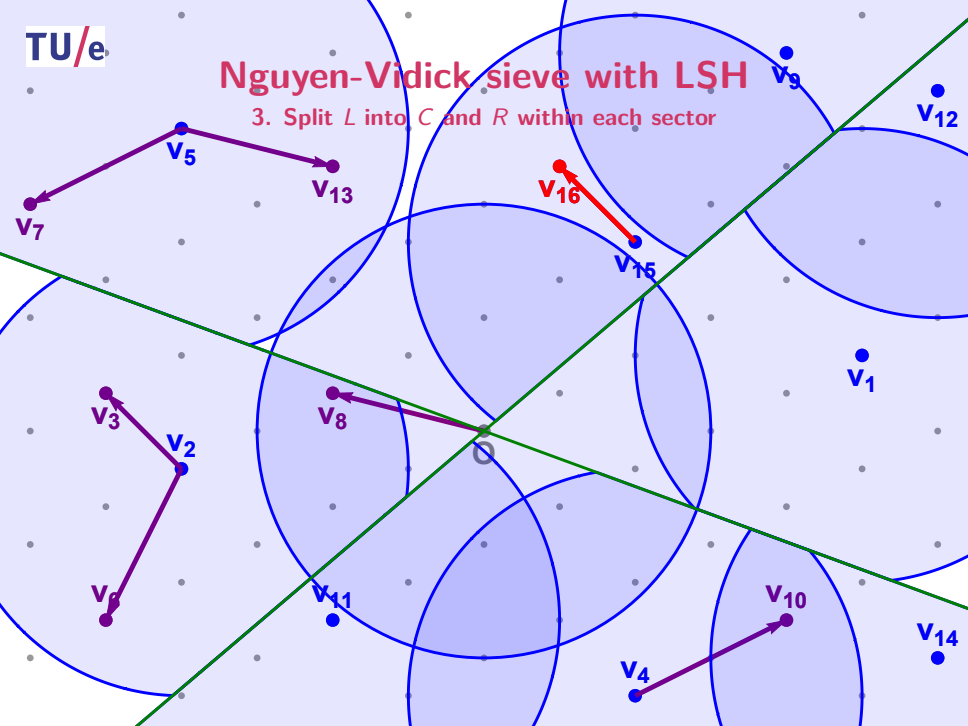
Nguyen-Vidick sieve with LSH

3. Split L into C and R within each sector



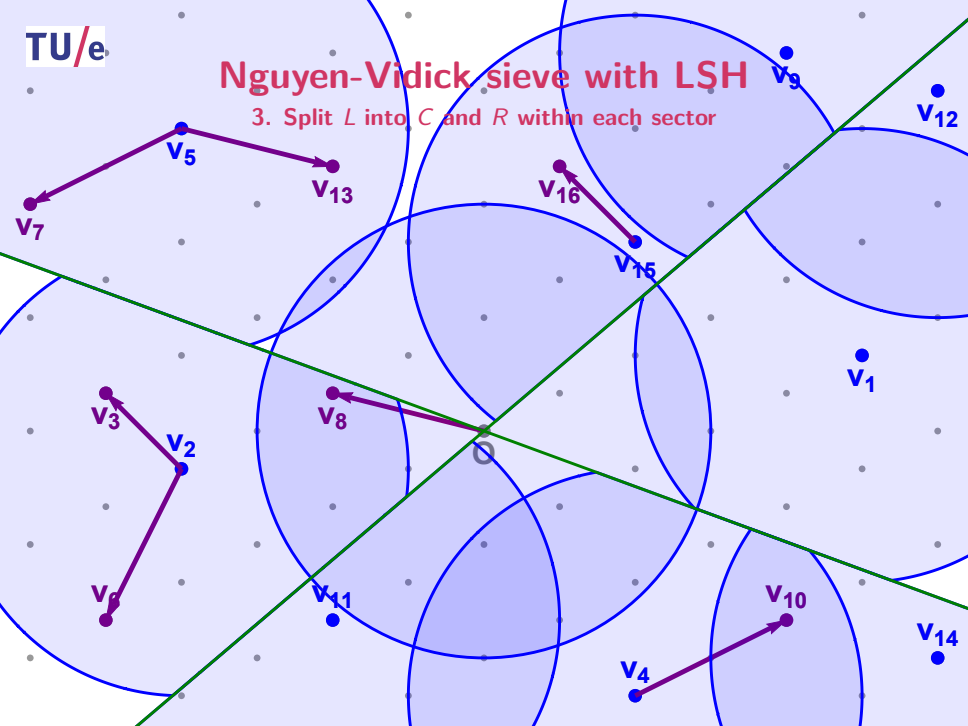
Nguyen-Vidick sieve with LSH

3. Split L into C and R within each sector



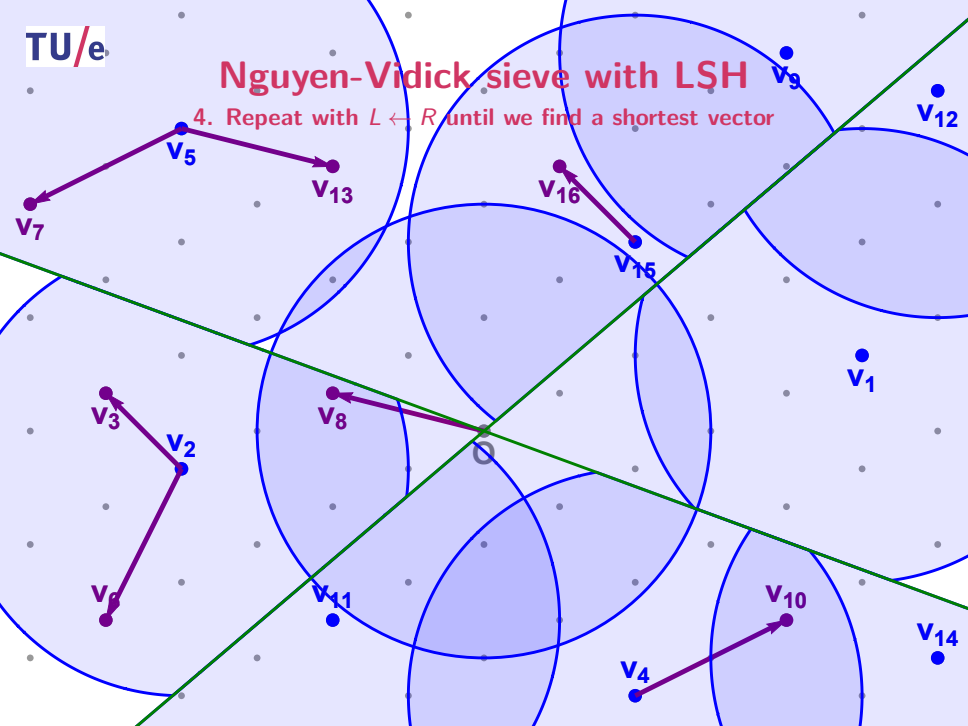
Nguyen-Vidick sieve with LSH

3. Split L into C and R within each sector



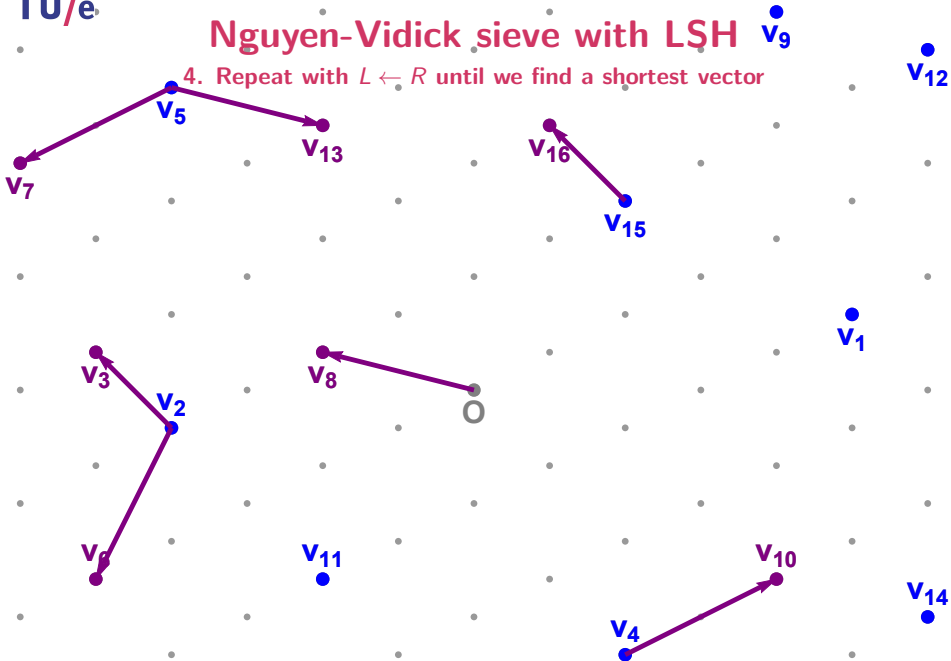
Nguyen-Vidick sieve with LSH

4. Repeat with $L \leftarrow R$ until we find a shortest vector



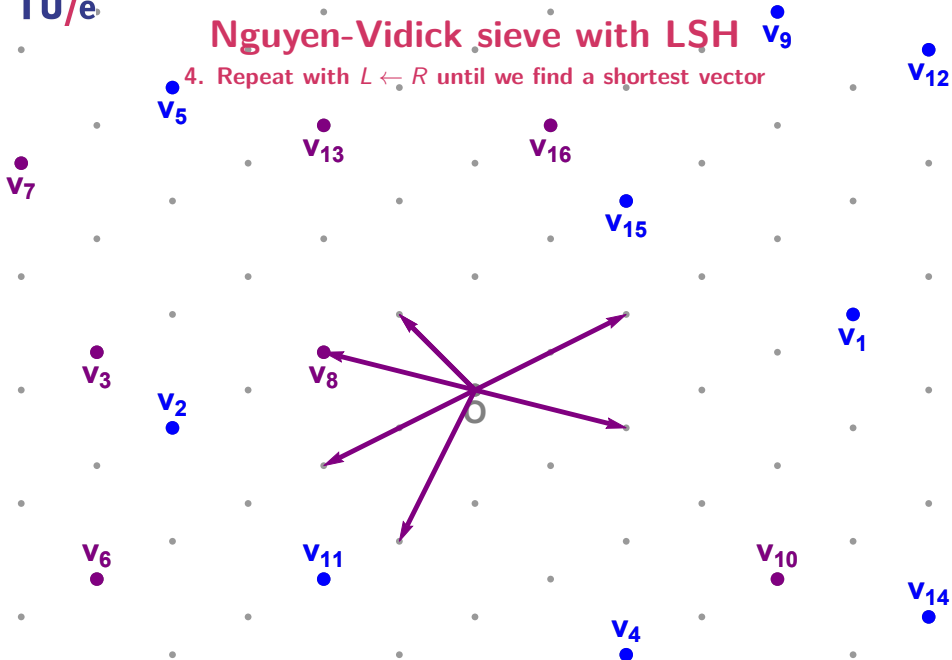
Nguyen-Vidick sieve with LSH

4. Repeat with $L \leftarrow R$ until we find a shortest vector



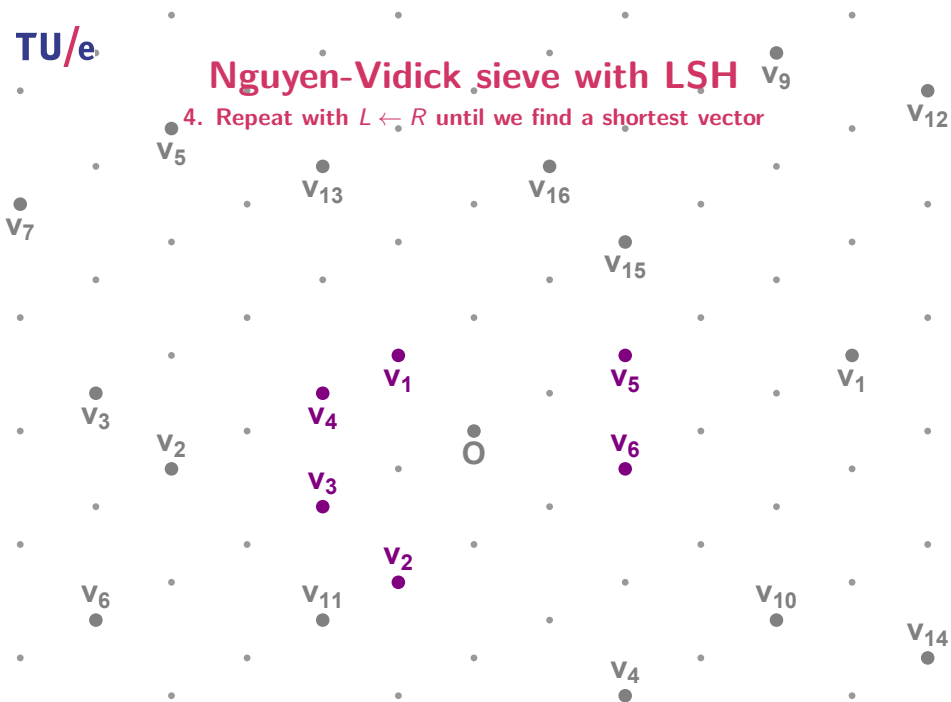
Nguyen-Vidick sieve with LSH

4. Repeat with $L \leftarrow R$ until we find a shortest vector



Nguyen-Vidick sieve with LSH

4. Repeat with $L \leftarrow R$ until we find a shortest vector



Nguyen-Vidick sieve with LSH

Overview



Nguyen-Vidick sieve with LSH

Overview

- Two parameters to tune
 - ▶ $k = O(n)$: Number of hyperplanes, leading to 2^k regions
 - ▶ $t = 2^{O(n)}$: Number of different, independent “hash tables”

Nguyen-Vidick sieve with LSH

Overview

- Two parameters to tune
 - ▶ $k = O(n)$: Number of hyperplanes, leading to 2^k regions
 - ▶ $t = 2^{O(n)}$: Number of different, independent “hash tables”
- Space complexity: $2^{0.34n+o(n)}$
 - ▶ Store $2^{0.13n}$ hash tables, each containing all $2^{0.21n}$ vectors

Nguyen-Vidick sieve with LSH

Overview

- Two parameters to tune
 - ▶ $k = O(n)$: Number of hyperplanes, leading to 2^k regions
 - ▶ $t = 2^{O(n)}$: Number of different, independent “hash tables”
- Space complexity: $2^{0.34n+o(n)}$
 - ▶ Store $2^{0.13n}$ hash tables, each containing all $2^{0.21n}$ vectors
- Time complexity: $2^{0.34n+o(n)}$
 - ▶ Compute $2^{0.13n}$ hashes, and go through $2^{0.13n}$ vectors
 - ▶ Repeat this for each of $2^{0.21n}$ vectors

Nguyen-Vidick sieve with LSH

Overview

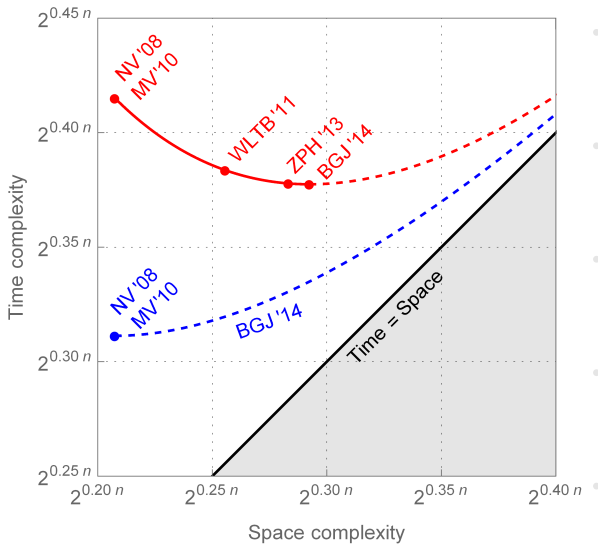
- Two parameters to tune
 - ▶ $k = O(n)$: Number of hyperplanes, leading to 2^k regions
 - ▶ $t = 2^{O(n)}$: Number of different, independent “hash tables”
- Space complexity: $2^{0.34n+o(n)}$
 - ▶ Store $2^{0.13n}$ hash tables, each containing all $2^{0.21n}$ vectors
- Time complexity: $2^{0.34n+o(n)}$
 - ▶ Compute $2^{0.13n}$ hashes, and go through $2^{0.13n}$ vectors
 - ▶ Repeat this for each of $2^{0.21n}$ vectors

Heuristic

Sieving with LSH runs in time $2^{0.34n+o(n)}$ and space $2^{0.34n+o(n)}$.

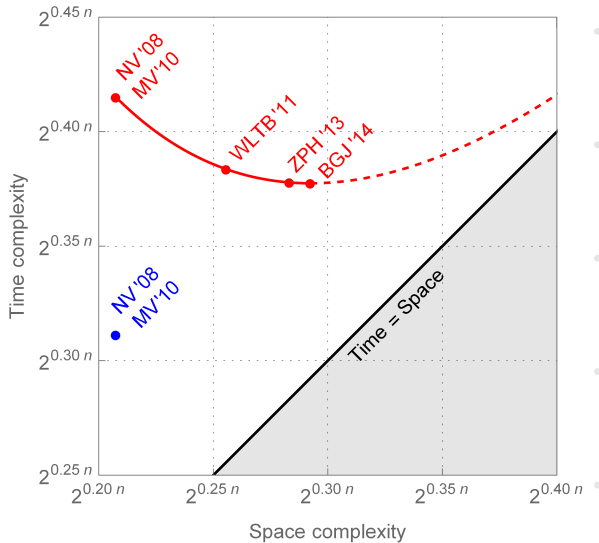
Sieving

Space/time trade-off



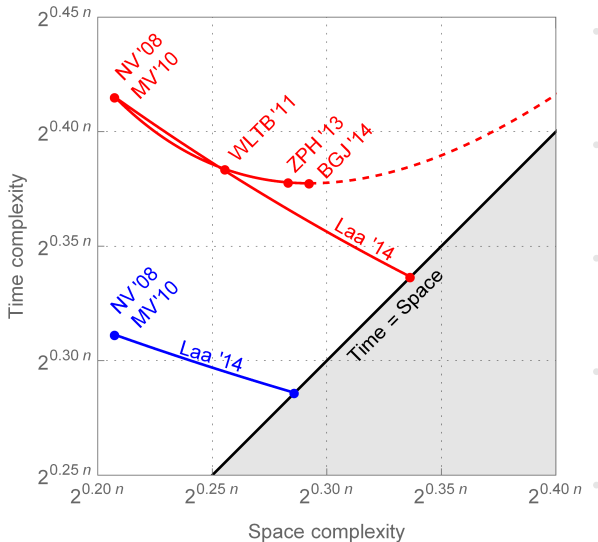
Sieving

Space/time trade-off



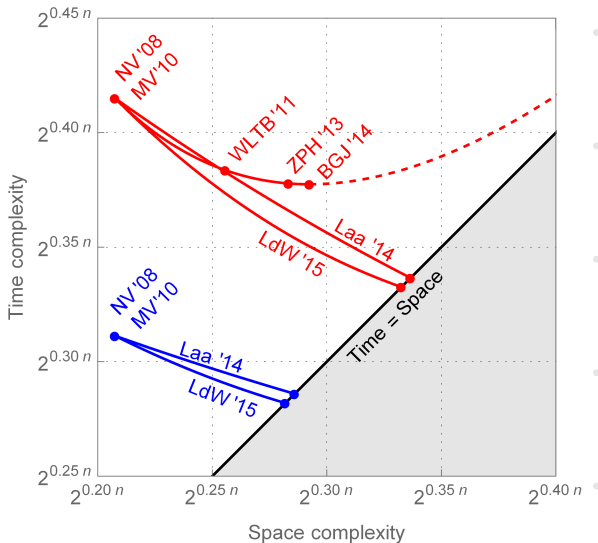
Sieving with LSH

Space/time trade-off



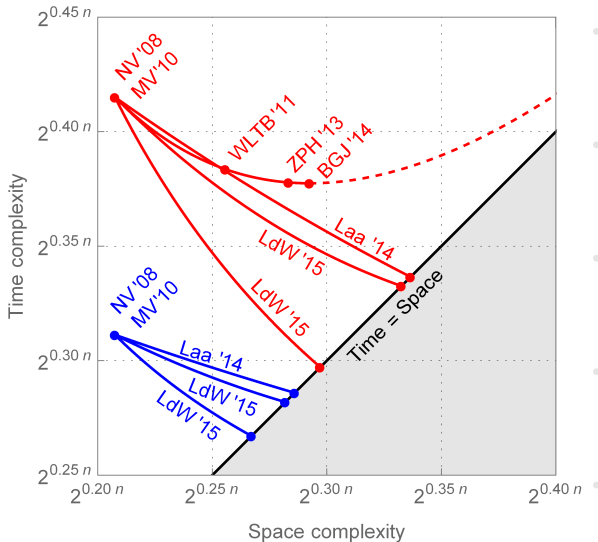
Sieving with LSH

Space/time trade-off



Sieving with LSH

Space/time trade-off



Questions

