

# From Collusion-Resistant Traitor Tracing to Efficient Probabilistic Group Testing

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## Collusion-resistant traitor tracing

Introduction

Score-based construction

Fighting against specific attacks

Results

## Efficient probabilistic group testing

Introduction

A traitor-tracing based solution

Results

## Conclusion

# Collusion-resistant traitor tracing

## Illegal redistribution

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User	Copyrighted content																
Antonino	0	1	1	1	0	0	1	1	1	0	1	1	0	0	1	0	...
Boris	0	1	1	1	0	0	1	1	1	0	1	1	0	0	1	0	...
Caroline	0	1	1	1	0	0	1	1	1	0	1	1	0	0	1	0	...
David	0	1	1	1	0	0	1	1	1	0	1	1	0	0	1	0	...
Eve	0	1	1	1	0	0	1	1	1	0	1	1	0	0	1	0	...
Fred	0	1	1	1	0	0	1	1	1	0	1	1	0	0	1	0	...
Gábor	0	1	1	1	0	0	1	1	1	0	1	1	0	0	1	0	...
Henry	0	1	1	1	0	0	1	1	1	0	1	1	0	0	1	0	...

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# Collusion-resistant traitor tracing

## Illegal redistribution

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Antonino	0	1	1	1	0	0	1	1	1	0	1	1	0	0	1	0	...
Boris	0	1	1	1	0	0	1	1	1	0	1	1	0	0	1	0	...
Caroline	0	1	1	1	0	0	1	1	1	0	1	1	0	0	1	0	...
David	0	1	1	1	0	0	1	1	1	0	1	1	0	0	1	0	...
Eve	0	1	1	1	0	0	1	1	1	0	1	1	0	0	1	0	...
Fred	0	1	1	1	0	0	1	1	1	0	1	1	0	0	1	0	...
Gábor	0	1	1	1	0	0	1	1	1	0	1	1	0	0	1	0	...
Henry	0	1	1	1	0	0	1	1	1	0	1	1	0	0	1	0	...
Copy	0	1	1	1	0	0	1	1	1	0	1	1	0	0	1	0	...

# Collusion-resistant traitor tracing

## Embedding fingerprints

User	Copyrighted content (fingerprinted)																
Antonino	0	1	1	1	0	0	1	1	1	0	1	1	0	1	0	0	...
Boris	0	1	1	1	0	1	0	1	1	0	1	1	1	1	1	0	...
Caroline	0	1	0	1	0	1	0	1	1	0	0	1	1	0	1	0	...
David	0	1	1	1	0	0	0	1	1	0	1	1	0	0	0	0	...
Eve	0	1	0	1	0	1	0	1	1	0	1	1	1	0	0	0	...
Fred	0	1	0	1	0	0	1	1	1	0	0	1	0	1	0	0	...
Gábor	0	1	1	1	0	1	1	1	1	0	1	1	0	0	1	0	...
Henry	0	1	0	1	0	1	1	1	1	0	0	1	0	1	1	0	...

# Collusion-resistant traitor tracing

## Embedding fingerprints

User	Copyrighted content (fingerprinted)																
Antonino	0	1	1	1	0	0	1	1	1	0	1	1	0	1	0	0	...
Boris	0	1	1	1	0	1	0	1	1	0	1	1	1	1	1	0	...
Caroline	0	1	0	1	0	1	0	1	1	0	0	1	1	0	1	0	...
David	0	1	1	1	0	0	0	1	1	0	1	1	0	0	0	0	...
Eve	0	1	0	1	0	1	0	1	1	0	1	1	1	0	0	0	...
Fred	0	1	0	1	0	0	1	1	1	0	0	1	0	1	0	0	...
Gábor	0	1	1	1	0	1	1	1	1	0	1	1	0	0	1	0	...
Henry	0	1	0	1	0	1	1	1	1	0	0	1	0	1	1	0	...
Copy	0	1	0	1	0	1	0	1	1	0	1	1	1	0	0	0	...

# Collusion-resistant traitor tracing

## Embedding fingerprints

User	Copyrighted content (fingerprinted)																
Antonino	0	1	<b>1</b>	1	0	<b>0</b>	<b>1</b>	1	1	0	<b>1</b>	1	<b>0</b>	<b>1</b>	<b>0</b>	0	...
Boris	0	1	<b>1</b>	1	0	<b>1</b>	<b>0</b>	1	1	0	<b>1</b>	1	<b>1</b>	<b>1</b>	<b>1</b>	0	...
Caroline	0	1	<b>0</b>	1	0	<b>1</b>	<b>0</b>	1	1	0	<b>0</b>	1	<b>1</b>	<b>0</b>	<b>1</b>	0	...
David	0	1	<b>1</b>	1	0	<b>0</b>	<b>0</b>	1	1	0	<b>1</b>	1	<b>0</b>	<b>0</b>	<b>0</b>	0	...
Eve	0	1	<b>0</b>	1	0	<b>1</b>	<b>0</b>	1	1	0	<b>1</b>	1	<b>1</b>	<b>0</b>	<b>0</b>	0	...
Fred	0	1	<b>0</b>	1	0	<b>0</b>	<b>1</b>	1	1	0	<b>0</b>	1	<b>0</b>	<b>1</b>	<b>0</b>	0	...
Gábor	0	1	<b>1</b>	1	0	<b>1</b>	<b>1</b>	1	1	0	<b>1</b>	1	<b>0</b>	<b>0</b>	<b>1</b>	0	...
Henry	0	1	<b>0</b>	1	0	<b>1</b>	<b>1</b>	1	1	0	<b>0</b>	1	<b>0</b>	<b>1</b>	<b>1</b>	0	...
Copy	0	1	<b>0</b>	1	0	<b>1</b>	<b>0</b>	1	1	0	<b>1</b>	1	<b>1</b>	<b>0</b>	<b>0</b>	0	...

# Collusion-resistant traitor tracing

## Embedding fingerprints

User	Copyrighted content (fingerprinted)																
Antonino	0	1	<b>1</b>	1	0	<b>0</b>	<b>1</b>	1	1	0	<b>1</b>	1	<b>0</b>	<b>1</b>	<b>0</b>	0	...
Boris	0	1	<b>1</b>	1	0	<b>1</b>	<b>0</b>	1	1	0	<b>1</b>	1	<b>1</b>	<b>1</b>	<b>1</b>	0	...
Caroline	0	1	<b>0</b>	1	0	<b>1</b>	<b>0</b>	1	1	0	<b>0</b>	1	<b>1</b>	<b>0</b>	<b>1</b>	0	...
David	0	1	<b>1</b>	1	0	<b>0</b>	<b>0</b>	1	1	0	<b>1</b>	1	<b>0</b>	<b>0</b>	<b>0</b>	0	...
Eve	0	1	<b>0</b>	1	0	<b>1</b>	<b>0</b>	1	1	0	<b>1</b>	1	<b>1</b>	<b>0</b>	<b>0</b>	0	...
Fred	0	1	<b>0</b>	1	0	<b>0</b>	<b>1</b>	1	1	0	<b>0</b>	1	<b>0</b>	<b>1</b>	<b>0</b>	0	...
Gábor	0	1	<b>1</b>	1	0	<b>1</b>	<b>1</b>	1	1	0	<b>1</b>	1	<b>0</b>	<b>0</b>	<b>1</b>	0	...
Henry	0	1	<b>0</b>	1	0	<b>1</b>	<b>1</b>	1	1	0	<b>0</b>	1	<b>0</b>	<b>1</b>	<b>1</b>	0	...
Copy	0	1	<b>0</b>	1	0	<b>1</b>	<b>0</b>	1	1	0	<b>1</b>	1	<b>1</b>	<b>0</b>	<b>0</b>	0	...



# Collusion-resistant traitor tracing

## Collusion attacks

User	Copyrighted content (fingerprinted)															
Antonino	0	1	1	1	0	0	1	1	1	0	1	1	0	1	0	...
Boris	0	1	1	1	0	1	0	1	1	0	1	1	1	1	0	...
Caroline	0	1	0	1	0	1	0	1	1	0	0	1	1	0	0	...
David	0	1	1	1	0	0	0	1	1	0	1	1	0	0	0	...
Eve	0	1	0	1	0	1	0	1	1	0	1	1	1	0	0	...
Fred	0	1	0	1	0	0	1	1	1	0	0	1	0	1	0	...
Gábor	0	1	1	1	0	1	1	1	1	0	1	1	0	0	1	...
Henry	0	1	0	1	0	1	1	1	1	0	0	1	0	1	1	...

# Collusion-resistant traitor tracing

## Collusion attacks

User	Copyrighted content (fingerprinted)																
Antonino	0	1	1	1	0	0	1	1	1	0	1	1	0	1	0	0	...
Boris	0	1	1	1	0	1	0	1	1	0	1	1	1	1	1	0	...
Caroline	0	1	0	1	0	1	0	1	1	0	0	1	1	0	1	0	...
David	0	1	1	1	0	0	0	1	1	0	1	1	0	0	0	0	...
Eve	0	1	0	1	0	1	0	1	1	0	1	1	1	0	0	0	...
Fred	0	1	0	1	0	0	1	1	1	0	0	1	0	1	0	0	...
Gábor	0	1	1	1	0	1	1	1	1	0	1	1	0	0	1	0	...
Henry	0	1	0	1	0	1	1	1	1	0	0	1	0	1	1	0	...
Copy	0	1	1	1	0	1	0	1	1	0	1	1	0	1	0	0	...

# Collusion-resistant traitor tracing

## Collusion attacks

User	Copyrighted content (fingerprinted)																
Antonino	0	1	<b>1</b>	1	0	<b>0</b>	<b>1</b>	1	1	0	<b>1</b>	1	<b>0</b>	<b>1</b>	<b>0</b>	0	...
Boris	0	1	<b>1</b>	1	0	<b>1</b>	<b>0</b>	1	1	0	<b>1</b>	1	<b>1</b>	<b>1</b>	<b>1</b>	0	...
Caroline	0	1	<b>0</b>	1	0	<b>1</b>	<b>0</b>	1	1	0	<b>0</b>	1	<b>1</b>	<b>0</b>	<b>1</b>	0	...
David	0	1	<b>1</b>	1	0	<b>0</b>	<b>0</b>	1	1	0	<b>1</b>	1	<b>0</b>	<b>0</b>	<b>0</b>	0	...
Eve	0	1	<b>0</b>	1	0	<b>1</b>	<b>0</b>	1	1	0	<b>1</b>	1	<b>1</b>	<b>0</b>	<b>0</b>	0	...
Fred	0	1	<b>0</b>	1	0	<b>0</b>	<b>1</b>	1	1	0	<b>0</b>	1	<b>0</b>	<b>1</b>	<b>0</b>	0	...
Gábor	0	1	<b>1</b>	1	0	<b>1</b>	<b>1</b>	1	1	0	<b>1</b>	1	<b>0</b>	<b>0</b>	<b>1</b>	0	...
Henry	0	1	<b>0</b>	1	0	<b>1</b>	<b>1</b>	1	1	0	<b>0</b>	1	<b>0</b>	<b>1</b>	<b>1</b>	0	...
Copy	0	1	<b>1</b>	1	0	<b>1</b>	<b>0</b>	1	1	0	<b>1</b>	1	<b>0</b>	<b>1</b>	<b>0</b>	0	...

# Collusion-resistant traitor tracing

Schemes resistant against collusion attacks

User	Copyrighted content (fingerprinted)																
Antonino	0	1	?	1	0	?	?	1	1	0	?	1	?	?	?	0	...
Boris	0	1	?	1	0	?	?	1	1	0	?	1	?	?	?	0	...
Caroline	0	1	?	1	0	?	?	1	1	0	?	1	?	?	?	0	...
David	0	1	?	1	0	?	?	1	1	0	?	1	?	?	?	0	...
Eve	0	1	?	1	0	?	?	1	1	0	?	1	?	?	?	0	...
Fred	0	1	?	1	0	?	?	1	1	0	?	1	?	?	?	0	...
Gábor	0	1	?	1	0	?	?	1	1	0	?	1	?	?	?	0	...
Henry	0	1	?	1	0	?	?	1	1	0	?	1	?	?	?	0	...
Copy	0	1	?	1	0	?	?	1	1	0	?	1	?	?	?	0	...

# Collusion-resistant traitor tracing

Schemes resistant against collusion attacks

User	Copyrighted content (fingerprinted)																
Antonino	0	1	?	1	0	?	?	1	1	0	?	1	?	?	?	0	...
Boris	0	1	?	1	0	?	?	1	1	0	?	1	?	?	?	0	...
Caroline	0	1	?	1	0	?	?	1	1	0	?	1	?	?	?	0	...
David	0	1	?	1	0	?	?	1	1	0	?	1	?	?	?	0	...
Eve	0	1	?	1	0	?	?	1	1	0	?	1	?	?	?	0	...
Fred	0	1	?	1	0	?	?	1	1	0	?	1	?	?	?	0	...
Gábor	0	1	?	1	0	?	?	1	1	0	?	1	?	?	?	0	...
Henry	0	1	?	1	0	?	?	1	1	0	?	1	?	?	?	0	...
Copy	0	1	?	1	0	?	?	1	1	0	?	1	?	?	?	0	...

1. An algorithm to construct collusion-resistant codes

# Collusion-resistant traitor tracing

Schemes resistant against collusion attacks

User	Copyrighted content (fingerprinted)																
Antonino	0	1	?	1	0	?	?	1	1	0	?	1	?	?	?	0	...
Boris	0	1	?	1	0	?	?	1	1	0	?	1	?	?	?	0	...
Caroline	0	1	?	1	0	?	?	1	1	0	?	1	?	?	?	0	...
David	0	1	?	1	0	?	?	1	1	0	?	1	?	?	?	0	...
Eve	0	1	?	1	0	?	?	1	1	0	?	1	?	?	?	0	...
Fred	0	1	?	1	0	?	?	1	1	0	?	1	?	?	?	0	...
Gábor	0	1	?	1	0	?	?	1	1	0	?	1	?	?	?	0	...
Henry	0	1	?	1	0	?	?	1	1	0	?	1	?	?	?	0	...
Copy	0	1	?	1	0	?	?	1	1	0	?	1	?	?	?	0	...

1. An algorithm to construct collusion-resistant codes
2. An algorithm to trace pirate copies to colluders

# Collusion-resistant traitor tracing

Schemes resistant against collusion attacks

User	Copyrighted content (fingerprinted)						
Antonino	?	?	?	?	?	?	...
Boris	?	?	?	?	?	?	...
Caroline	?	?	?	?	?	?	...
David	?	?	?	?	?	?	...
Eve	?	?	?	?	?	?	...
Fred	?	?	?	?	?	?	...
Gábor	?	?	?	?	?	?	...
Henry	?	?	?	?	?	?	...
Copy	?	?	?	?	?	?	...

1. An algorithm to construct collusion-resistant codes
2. An algorithm to trace pirate copies to colluders

# Collusion-resistant traitor tracing

Schemes resistant against collusion attacks

User	Copyrighted content (fingerprinted)		
Antonino		...	
Boris		...	
Caroline		...	
David	$X \in \{0, 1\}^{n \times \ell}$	...	
Eve		...	
Fred		...	
Gábor		...	
Henry		...	
Copy		$y \in \{0, 1\}^{\ell}$	...

1. An algorithm to construct collusion-resistant codes
2. An algorithm to trace pirate copies to colluders



# Collusion-resistant traitor tracing

Schemes resistant against collusion attacks

1. An algorithm to construct collusion-resistant codes
2. An algorithm to trace pirate copies to colluders

# Collusion-resistant traitor tracing

Schemes resistant against collusion attacks

1. An algorithm to construct collusion-resistant codes
2. An algorithm to trace pirate copies to colluders

# Score-based construction

## Overview

1. An algorithm to construct collusion-resistant codes
2. An algorithm to trace pirate copies to colluders

# Score-based construction

## Overview

1. An algorithm to construct collusion-resistant codes
  - 1a. For each segment  $i$ , generate  $p_i \sim F$ .
  - 1b. For each segment  $i$ , user  $j$ , choose  $X_{j,i} = 1$  with prob.  $p_i$ .
2. An algorithm to trace pirate copies to colluders

# Score-based construction

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  - 1a. For each segment  $i$ , generate  $p_i \sim F$ .
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2. An algorithm to trace pirate copies to colluders
  - 2a. For each segment  $i$ , user  $j$ , calculate  $S_{j,i} = g(X_{j,i}, y_i, p_i)$ .
  - 2b. For each user  $j$ , accuse user  $j$  iff  $\sum_i S_{j,i}$  is “large”.

## Score-based construction

### Arbitrary attacks

1. An algorithm to construct collusion-resistant codes
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$$F(p_i) = \frac{2}{\pi} \arcsin \sqrt{p_i}$$

$$g(X_{j,i}, y_i, p_i) = \begin{cases} +0 & (X_{j,i}, y_i) = (0, 0) \\ -\sqrt{p_i/(1-p_i)} & (X_{j,i}, y_i) = (0, 1) \\ -0 & (X_{j,i}, y_i) = (1, 0) \\ +\sqrt{(1-p_i)/p_i} & (X_{j,i}, y_i) = (1, 1) \end{cases}$$

## Score-based construction

### Arbitrary attacks

1. An algorithm to construct collusion-resistant codes
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$$\ell = 100c^2 \ln n \quad [\text{Tar'03}]$$

## Score-based construction

### Arbitrary attacks

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$$\ell \sim 4\pi^2 c^2 \ln n \quad [\text{S+'06}]$$



## Score-based construction

### Arbitrary attacks

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$$F(p_i) = \frac{2}{\pi} \arcsin \sqrt{p_i}$$

$$g(X_{j,i}, y_i, p_i) = \begin{cases} +0 & (X_{j,i}, y_i) = (0, 0) \\ -\sqrt{p_i/(1-p_i)} & (X_{j,i}, y_i) = (0, 1) \\ -0 & (X_{j,i}, y_i) = (1, 0) \\ +\sqrt{(1-p_i)/p_i} & (X_{j,i}, y_i) = (1, 1) \end{cases}$$

$$\ell \sim 2\pi^2 c^2 \ln n \quad [\text{BT}'08]$$

## Score-based construction

### Arbitrary attacks

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  - 2b. For each user  $j$ , accuse user  $j$  iff  $\sum_i S_{j,i}$  is "large".

$$F(p_i) = \frac{2}{\pi} \arcsin \sqrt{p_i}$$

$$g(X_{j,i}, y_i, p_i) = \begin{cases} +\sqrt{p_i/(1-p_i)} & (X_{j,i}, y_i) = (0, 0) \\ -\sqrt{p_i/(1-p_i)} & (X_{j,i}, y_i) = (0, 1) \\ -\sqrt{(1-p_i)/p_i} & (X_{j,i}, y_i) = (1, 0) \\ +\sqrt{(1-p_i)/p_i} & (X_{j,i}, y_i) = (1, 1) \end{cases}$$

## Score-based construction

### Arbitrary attacks

1. An algorithm to construct collusion-resistant codes
  - 1a. For each segment  $i$ , generate  $p_i \sim F$ .
  - 1b. For each segment  $i$ , user  $j$ , choose  $X_{j,i} = 1$  with prob.  $p_i$ .
2. An algorithm to trace pirate copies to colluders
  - 2a. For each segment  $i$ , user  $j$ , calculate  $S_{j,i} = g(X_{j,i}, y_i, p_i)$ .
  - 2b. For each user  $j$ , accuse user  $j$  iff  $\sum_i S_{j,i}$  is “large”.

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$$\ell \sim \pi^2 c^2 \ln n \quad [\text{S+'08}]$$

## Score-based construction

### Arbitrary attacks

1. An algorithm to construct collusion-resistant codes
  - 1a. For each segment  $i$ , generate  $p_i \sim F$ .
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$$F(p_i) = \frac{2}{\pi} \arcsin \sqrt{p_i}$$

$$g(X_{j,i}, y_i, p_i) = \begin{cases} +\sqrt{p_i/(1-p_i)} & (X_{j,i}, y_i) = (0, 0) \\ -\sqrt{p_i/(1-p_i)} & (X_{j,i}, y_i) = (0, 1) \\ -\sqrt{(1-p_i)/p_i} & (X_{j,i}, y_i) = (1, 0) \\ +\sqrt{(1-p_i)/p_i} & (X_{j,i}, y_i) = (1, 1) \end{cases}$$

$$\ell \sim \frac{1}{2} \pi^2 c^2 \ln n \quad [\text{LdW}'13]$$

## Score-based construction

### Arbitrary attacks

1. An algorithm to construct collusion-resistant codes
  - 1a. For each segment  $i$ , generate  $p_i \sim F$ .
  - 1b. For each segment  $i$ , user  $j$ , choose  $X_{j,i} = 1$  with prob.  $p_i$ .
2. An algorithm to trace pirate copies to colluders
  - 2a. For each segment  $i$ , user  $j$ , calculate  $S_{j,i} = g(X_{j,i}, y_i, p_i)$ .
  - 2b. For each user  $j$ , accuse user  $j$  iff  $\sum_i S_{j,i}$  is "large".

$$F(p_i) = \frac{2}{\pi} \arcsin \sqrt{p_i}$$

$$g(X_{j,i}, y_i, p_i) = \begin{cases} +p_i/(1-p_i) & (X_{j,i}, y_i) = (0, 0) \\ -1 & (X_{j,i}, y_i) = (0, 1) \\ -1 & (X_{j,i}, y_i) = (1, 0) \\ +(1-p_i)/p_i & (X_{j,i}, y_i) = (1, 1) \end{cases}$$

# Score-based construction

## Arbitrary attacks

1. An algorithm to construct collusion-resistant codes
  - 1a. For each segment  $i$ , generate  $p_i \sim F$ .
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$$\ell \sim 2c^2 \ln n \quad [\text{OSD}'13]$$

## Fighting against specific attacks

What can pirates do?

---

Antonino	$X_{1,1}$	$X_{1,2}$	$X_{1,3}$	$X_{1,4}$	$X_{1,5}$	...	$X_{1,l}$
Boris	$X_{2,1}$	$X_{2,2}$	$X_{2,3}$	$X_{2,4}$	$X_{2,5}$	...	$X_{2,l}$
Caroline	$X_{3,1}$	$X_{3,2}$	$X_{3,3}$	$X_{3,4}$	$X_{3,5}$	...	$X_{3,l}$
David	$X_{4,1}$	$X_{4,2}$	$X_{4,3}$	$X_{4,4}$	$X_{4,5}$	...	$X_{4,l}$
Eve	$X_{5,1}$	$X_{5,2}$	$X_{5,3}$	$X_{5,4}$	$X_{5,5}$	...	$X_{5,l}$
Fred	$X_{6,1}$	$X_{6,2}$	$X_{6,3}$	$X_{6,4}$	$X_{6,5}$	...	$X_{6,l}$
Gábor	$X_{7,1}$	$X_{7,2}$	$X_{7,3}$	$X_{7,4}$	$X_{7,5}$	...	$X_{7,l}$
Henry	$X_{8,1}$	$X_{8,2}$	$X_{8,3}$	$X_{8,4}$	$X_{8,5}$	...	$X_{8,l}$

---

Copy	$y_1$	$y_2$	$y_3$	$y_4$	$y_5$	...	$y_l$
------	-------	-------	-------	-------	-------	-----	-------

## Fighting against specific attacks

What can pirates do?

---

Antonino	0	0	1	1	1	...	0
Boris	1	0	1	1	1	...	1
Caroline	1	0	0	1	0	...	0
David	0	0	1	1	1	...	0
Eve	0	0	1	0	1	...	0
Fred	1	0	1	0	1	...	0
Gábor	0	0	1	0	1	...	0
Henry	1	0	0	0	1	...	0

---

Copy	$y_1$	$y_2$	$y_3$	$y_4$	$y_5$	...	$y_\ell$
------	-------	-------	-------	-------	-------	-----	----------



## Fighting against specific attacks

What can pirates do?

---

Antonino	.	.	.	.	.	...	.
Boris	.	.	.	.	.	...	.
Caroline	1	0	0	1	0	...	0
David	.	.	.	.	.	...	.
Eve	0	0	1	0	1	...	0
Fred	.	.	.	.	.	...	.
Gábor	.	.	.	.	.	...	.
Henry	1	0	0	0	1	...	0

---

Copy	$y_1$	$y_2$	$y_3$	$y_4$	$y_5$	...	$y_\ell$
------	-------	-------	-------	-------	-------	-----	----------

## Fighting against specific attacks

What can pirates do?

---

Antonino	.	.	.	.	.	...	.
Boris	.	.	.	.	.	...	.
Caroline	<b>1</b>	0	<b>0</b>	<b>1</b>	<b>0</b>	...	0
David	.	.	.	.	.	...	.
Eve	<b>0</b>	0	<b>1</b>	<b>0</b>	<b>1</b>	...	0
Fred	.	.	.	.	.	...	.
Gábor	.	.	.	.	.	...	.
Henry	<b>1</b>	0	<b>0</b>	<b>0</b>	<b>1</b>	...	0

---

Copy	$y_1$	$y_2$	$y_3$	$y_4$	$y_5$	...	$y_\ell$
------	-------	-------	-------	-------	-------	-----	----------

## Fighting against specific attacks

What can pirates do?

---

Antonino	.	.	.	.	.	...	.
Boris	.	.	.	.	.	...	.
Caroline	<b>1</b>	0	<b>0</b>	<b>1</b>	<b>0</b>	...	0
David	.	.	.	.	.	...	.
Eve	<b>0</b>	0	<b>1</b>	<b>0</b>	<b>1</b>	...	0
Fred	.	.	.	.	.	...	.
Gábor	.	.	.	.	.	...	.
Henry	<b>1</b>	0	<b>0</b>	<b>0</b>	<b>1</b>	...	0

---

Copy	0/1	0	0/1	0/1	0/1	...	0
------	-----	---	-----	-----	-----	-----	---

## Fighting against specific attacks

What can pirates do?

---

Antonino	.	.	.	.	.	...	.
Boris	.	.	.	.	.	...	.
Caroline	<b>1</b>	0	<b>0</b>	<b>1</b>	<b>0</b>	...	0
David	.	.	.	.	.	...	.
Eve	<b>0</b>	0	<b>1</b>	<b>0</b>	<b>1</b>	...	0
Fred	.	.	.	.	.	...	.
Gábor	.	.	.	.	.	...	.
Henry	<b>1</b>	0	<b>0</b>	<b>0</b>	<b>1</b>	...	0

---

Copy	0/ <b>1</b>	0	0/ <b>1</b>	0/ <b>1</b>	0/ <b>1</b>	...	0
All-1	1	0	1	1	1	...	0

## Fighting against specific attacks

What can pirates do?

Antonino	.	.	.	.	.	...	.
Boris	.	.	.	.	.	...	.
Caroline	<b>1</b>	0	<b>0</b>	<b>1</b>	<b>0</b>	...	0
David	.	.	.	.	.	...	.
Eve	<b>0</b>	0	<b>1</b>	<b>0</b>	<b>1</b>	...	0
Fred	.	.	.	.	.	...	.
Gábor	.	.	.	.	.	...	.
Henry	<b>1</b>	0	<b>0</b>	<b>0</b>	<b>1</b>	...	0
Copy	0/1	0	0/1	0/1	0/1	...	0
All-1	1	0	1	1	1	...	0
Minority	0	0	1	1	0	...	0

# Fighting against specific attacks

What can pirates do?

Antonino	.	.	.	.	.	...	.
Boris	.	.	.	.	.	...	.
Caroline	<b>1</b>	0	<b>0</b>	<b>1</b>	<b>0</b>	...	0
David	.	.	.	.	.	...	.
Eve	<b>0</b>	0	<b>1</b>	<b>0</b>	<b>1</b>	...	0
Fred	.	.	.	.	.	...	.
Gábor	.	.	.	.	.	...	.
Henry	<b>1</b>	0	<b>0</b>	<b>0</b>	<b>1</b>	...	0
Copy	0/1	0	0/1	0/1	0/1	...	0
All-1	1	0	1	1	1	...	0
Minority	0	0	1	1	0	...	0
Majority	1	0	0	0	1	...	0

# Fighting against specific attacks

What can pirates do?

Antonino	.	.	.	.	.	...	.
Boris	.	.	.	.	.	...	.
Caroline	<b>1</b>	0	<b>0</b>	<b>1</b>	<b>0</b>	...	0
David	.	.	.	.	.	...	.
Eve	<b>0</b>	0	<b>1</b>	<b>0</b>	<b>1</b>	...	0
Fred	.	.	.	.	.	...	.
Gábor	.	.	.	.	.	...	.
Henry	<b>1</b>	0	<b>0</b>	<b>0</b>	<b>1</b>	...	0
Copy	<b>0/1</b>	0	<b>0/1</b>	<b>0/1</b>	<b>0/1</b>	...	0
All-1	1	0	1	1	1	...	0
Minority	0	0	1	1	0	...	0
Majority	1	0	0	0	1	...	0
Coin-flip	1	0	1	0	0	...	0

# Fighting against specific attacks

What can pirates do?

Antonino	.	.	.	.	.	...	.
Boris	.	.	.	.	.	...	.
Caroline	<b>1</b>	0	<b>0</b>	<b>1</b>	<b>0</b>	...	0
David	.	.	.	.	.	...	.
Eve	<b>0</b>	0	<b>1</b>	<b>0</b>	<b>1</b>	...	0
Fred	.	.	.	.	.	...	.
Gábor	.	.	.	.	.	...	.
Henry	<b>1</b>	0	<b>0</b>	<b>0</b>	<b>1</b>	...	0
Copy	0/1	0	0/1	0/1	0/1	...	0
All-1	1	0	1	1	1	...	0
Minority	0	0	1	1	0	...	0
Majority	1	0	0	0	1	...	0
Coin-flip	1	0	1	0	0	...	0
Interleaving	0	0	0	0	1	...	0



# Fighting against specific attacks

## Arbitrary attacks

1. An algorithm to construct collusion-resistant codes
  - 1a. For each segment  $i$ , generate  $p_i \sim F$ .
  - 1b. For each segment  $i$ , user  $j$ , choose  $X_{j,i} = 1$  with prob.  $p_i$ .
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  - 2a. For each segment  $i$ , user  $j$ , calculate  $S_{j,i} = g(X_{j,i}, y_i, p_i)$ .
  - 2b. For each user  $j$ , accuse user  $j$  iff  $\sum_i S_{j,i}$  is "large".

$$F(p_i) = \frac{2}{\pi} \arcsin \sqrt{p_i}$$

$$g(X_{j,i}, y_i, p_i) = \begin{cases} +p_i/(1-p_i) & (X_{j,i}, y_i) = (0, 0) \\ -1 & (X_{j,i}, y_i) = (0, 1) \\ -1 & (X_{j,i}, y_i) = (1, 0) \\ +(1-p_i)/p_i & (X_{j,i}, y_i) = (1, 1) \end{cases}$$

$$\ell \sim 2c^2 \ln n \quad [\text{OSD}'13]$$

# Fighting against specific attacks

## The interleaving attack

1. An algorithm to construct collusion-resistant codes
  - 1a. For each segment  $i$ , generate  $p_i \sim F$ .
  - 1b. For each segment  $i$ , user  $j$ , choose  $X_{j,i} = 1$  with prob.  $p_i$ .
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  - 2a. For each segment  $i$ , user  $j$ , calculate  $S_{j,i} = g(X_{j,i}, y_i, p_i)$ .
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$$\ell \sim 2c^2 \ln n \quad [\text{OSD}'13]$$

## Fighting against specific attacks

### The interleaving attack

1. An algorithm to construct collusion-resistant codes
  - 1a. For each segment  $i$ , generate  $p_i \sim F$ .
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  - 2a. For each segment  $i$ , user  $j$ , calculate  $S_{j,i} = g(X_{j,i}, y_i, p_i)$ .
  - 2b. For each user  $j$ , accuse user  $j$  iff  $\sum_i S_{j,i}$  is "large".

$$F(p_i) = 1 \left\{ p_i \geq p = \frac{1}{2} \right\}$$

$$g(X_{j,i}, y_i, p_i) = \begin{cases} +1 & (X_{j,i}, y_i) = (0, 0) \\ -1 & (X_{j,i}, y_i) = (0, 1) \\ -1 & (X_{j,i}, y_i) = (1, 0) \\ +1 & (X_{j,i}, y_i) = (1, 1) \end{cases}$$

$$\ell \sim 2c^2 \ln n$$

# Fighting against specific attacks

## The all-1 attack

1. An algorithm to construct collusion-resistant codes
  - 1a. For each segment  $i$ , generate  $p_i \sim F$ .
  - 1b. For each segment  $i$ , user  $j$ , choose  $X_{j,i} = 1$  with prob.  $p_i$ .
2. An algorithm to trace pirate copies to colluders
  - 2a. For each segment  $i$ , user  $j$ , calculate  $S_{j,i} = g(X_{j,i}, y_i, p_i)$ .
  - 2b. For each user  $j$ , accuse user  $j$  iff  $\sum_i S_{j,i}$  is "large".

$$F(p_i) = \frac{2}{\pi} \arcsin \sqrt{p_i}$$

$$g(X_{j,i}, y_i, p_i) = \begin{cases} +p_i/(1-p_i) & (X_{j,i}, y_i) = (0, 0) \\ -p_i(1-p_i)^{c-1}/(1-(1-p_i)^c) & (X_{j,i}, y_i) = (0, 1) \\ -1 & (X_{j,i}, y_i) = (1, 0) \\ +(1-p_i)^c/(1-(1-p_i)^c) & (X_{j,i}, y_i) = (1, 1) \end{cases}$$

$$\ell \sim O(c^{1.5} \ln n) \quad [\text{OSD}'13]$$

## Fighting against specific attacks

### The all-1 attack

1. An algorithm to construct collusion-resistant codes
  - 1a. For each segment  $i$ , generate  $p_i \sim F$ .
  - 1b. For each segment  $i$ , user  $j$ , choose  $X_{j,i} = 1$  with prob.  $p_i$ .
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  - 2a. For each segment  $i$ , user  $j$ , calculate  $S_{j,i} = g(X_{j,i}, y_i, p_i)$ .
  - 2b. For each user  $j$ , accuse user  $j$  iff  $\sum_i S_{j,i}$  is "large".

$$F(p_i) = 1 \left\{ p_i \geq p \approx \frac{1}{c} \right\}$$

$$g(X_{j,i}, y_i, p_i) = \begin{cases} +p_i/(1-p_i) & (X_{j,i}, y_i) = (0, 0) \\ -p_i(1-p_i)^{c-1}/(1-(1-p_i)^c) & (X_{j,i}, y_i) = (0, 1) \\ -1 & (X_{j,i}, y_i) = (1, 0) \\ +(1-p_i)^c/(1-(1-p_i)^c) & (X_{j,i}, y_i) = (1, 1) \end{cases}$$

$$\ell \sim 2c \ln n$$

# Fighting against specific attacks

## The minority voting attack

1. An algorithm to construct collusion-resistant codes
  - 1a. For each segment  $i$ , generate  $p_i \sim F$ .
  - 1b. For each segment  $i$ , user  $j$ , choose  $X_{j,i} = 1$  with prob.  $p_i$ .
2. An algorithm to trace pirate copies to colluders
  - 2a. For each segment  $i$ , user  $j$ , calculate  $S_{j,i} = g(X_{j,i}, y_i, p_i)$ .
  - 2b. For each user  $j$ , accuse user  $j$  iff  $\sum_i S_{j,i}$  is "large".

$$F(p_i) = \frac{2}{\pi} \arcsin \sqrt{p_i}$$

$$g(X_{j,i}, y_i, p_i) = \begin{cases} + \dots & (X_{j,i}, y_i) = (0, 0) \\ - \dots & (X_{j,i}, y_i) = (0, 1) \\ - \dots & (X_{j,i}, y_i) = (1, 0) \\ + \dots & (X_{j,i}, y_i) = (1, 1) \end{cases}$$

$$\ell \sim O(c^{1.5} \ln n) \quad [\text{OSD}'13]$$

## Fighting against specific attacks

### The minority voting attack

1. An algorithm to construct collusion-resistant codes
  - 1a. For each segment  $i$ , generate  $p_i \sim F$ .
  - 1b. For each segment  $i$ , user  $j$ , choose  $X_{j,i} = 1$  with prob.  $p_i$ .
2. An algorithm to trace pirate copies to colluders
  - 2a. For each segment  $i$ , user  $j$ , calculate  $S_{j,i} = g(X_{j,i}, y_i, p_i)$ .
  - 2b. For each user  $j$ , accuse user  $j$  iff  $\sum_i S_{j,i}$  is "large".

$$F(p_i) = 1 \left\{ p_i \geq p \approx \frac{1}{c} \right\}$$

$$g(X_{j,i}, y_i, p_i) \approx \begin{cases} +p_i/(1-p_i) & (X_{j,i}, y_i) = (0, 0) \\ -p_i(1-p_i)^{c-1}/(1-(1-p_i)^c) & (X_{j,i}, y_i) = (0, 1) \\ -1 & (X_{j,i}, y_i) = (1, 0) \\ +(1-p_i)^c/(1-(1-p_i)^c) & (X_{j,i}, y_i) = (1, 1) \end{cases}$$

$$\ell \sim 2c \ln n$$

# Fighting against specific attacks

## The majority voting attack

1. An algorithm to construct collusion-resistant codes
  - 1a. For each segment  $i$ , generate  $p_i \sim F$ .
  - 1b. For each segment  $i$ , user  $j$ , choose  $X_{j,i} = 1$  with prob.  $p_i$ .
2. An algorithm to trace pirate copies to colluders
  - 2a. For each segment  $i$ , user  $j$ , calculate  $S_{j,i} = g(X_{j,i}, y_i, p_i)$ .
  - 2b. For each user  $j$ , accuse user  $j$  iff  $\sum_i S_{j,i}$  is "large".

$$F(p_i) = \frac{2}{\pi} \arcsin \sqrt{p_i}$$

$$g(X_{j,i}, y_i, p_i) = \begin{cases} + \dots & (X_{j,i}, y_i) = (0, 0) \\ - \dots & (X_{j,i}, y_i) = (0, 1) \\ - \dots & (X_{j,i}, y_i) = (1, 0) \\ + \dots & (X_{j,i}, y_i) = (1, 1) \end{cases}$$

$$\ell \sim O(c^{1.5} \ln n) \quad [\text{OSD}'13]$$



## Fighting against specific attacks

### The majority voting attack

1. An algorithm to construct collusion-resistant codes
  - 1a. For each segment  $i$ , generate  $p_i \sim F$ .
  - 1b. For each segment  $i$ , user  $j$ , choose  $X_{j,i} = 1$  with prob.  $p_i$ .
2. An algorithm to trace pirate copies to colluders
  - 2a. For each segment  $i$ , user  $j$ , calculate  $S_{j,i} = g(X_{j,i}, y_i, p_i)$ .
  - 2b. For each user  $j$ , accuse user  $j$  iff  $\sum_i S_{j,i}$  is "large".

$$F(p_i) = \mathbf{1} \left\{ p_i \geq p = \frac{1}{2} \right\}$$

$$g(X_{j,i}, y_i, p_i) \approx \begin{cases} +1 & (X_{j,i}, y_i) = (0, 0) \\ -1 & (X_{j,i}, y_i) = (0, 1) \\ -1 & (X_{j,i}, y_i) = (1, 0) \\ +1 & (X_{j,i}, y_i) = (1, 1) \end{cases}$$

$$\ell \sim \pi c \ln n$$

# Fighting against specific attacks

## The coin-flip attack

1. An algorithm to construct collusion-resistant codes
  - 1a. For each segment  $i$ , generate  $p_i \sim F$ .
  - 1b. For each segment  $i$ , user  $j$ , choose  $X_{j,i} = 1$  with prob.  $p_i$ .
2. An algorithm to trace pirate copies to colluders
  - 2a. For each segment  $i$ , user  $j$ , calculate  $S_{j,i} = g(X_{j,i}, y_i, p_i)$ .
  - 2b. For each user  $j$ , accuse user  $j$  iff  $\sum_i S_{j,i}$  is "large".

$$F(p_i) = \frac{2}{\pi} \arcsin \sqrt{p_i}$$

$$g(X_{j,i}, y_i, p_i) = \begin{cases} +p_i/(1 - p_i^c + (1 - p_i)^c) & (X_{j,i}, y_i) = (0, 0) \\ -p_i/(1 + p_i^c - (1 - p_i)^c) & (X_{j,i}, y_i) = (0, 1) \\ -(1 - p_i)/(1 - p_i^c + (1 - p_i)^c) & (X_{j,i}, y_i) = (1, 0) \\ +(1 - p_i)/(1 + p_i^c - (1 - p_i)^c) & (X_{j,i}, y_i) = (1, 1) \end{cases}$$

$$\ell \sim O(c^{1.5} \ln n) \quad [\text{OSD}'13]$$

# Fighting against specific attacks

## The coin-flip attack

1. An algorithm to construct collusion-resistant codes
  - 1a. For each segment  $i$ , generate  $p_i \sim F$ .
  - 1b. For each segment  $i$ , user  $j$ , choose  $X_{j,i} = 1$  with prob.  $p_i$ .
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$$F(p_i) = 1 \left\{ p_i \geq p \approx \frac{1}{c} \right\}$$

$$g(X_{j,i}, y_i, p_i) \approx \begin{cases} +p_i/(1 - p_i^c + (1 - p_i)^c) & (X_{j,i}, y_i) = (0, 0) \\ -p_i/(1 + p_i^c - (1 - p_i)^c) & (X_{j,i}, y_i) = (0, 1) \\ -(1 - p_i)/(1 - p_i^c + (1 - p_i)^c) & (X_{j,i}, y_i) = (1, 0) \\ +(1 - p_i)/(1 + p_i^c - (1 - p_i)^c) & (X_{j,i}, y_i) = (1, 1) \end{cases}$$

$$\ell \sim 4c \ln n$$

# Results

## The Tardos scheme

Table: Asymptotics of  $L = \frac{\ell}{\ln n}$  for large  $n$ .

	Efficient constr.		Lower bounds	
Arbitrary attacks	$100c^2$	[Tar'03]	$\Omega(c^2)$	[Tar'03]
Interleaving attack	$100c^2$	[Tar'03]	$\Omega(c)$	
All-1 attack	$100c^2$	[Tar'03]	$\Omega(c)$	
Minority voting	$100c^2$	[Tar'03]	$\Omega(c)$	
Majority voting	$100c^2$	[Tar'03]	$\Omega(c)$	
Coin-flip attack	$100c^2$	[Tar'03]	$\Omega(c)$	

# Results

## Improvements of the Tardos scheme

Table: Asymptotics of  $L = \frac{\ell}{\ln n}$  for large  $n$ .

	Efficient constr.		Lower bounds	
Arbitrary attacks	$2c^2$	[OSD'13]	$2c^2$	[HM'12]
Interleaving attack	$2c^2$	[OSD'13]	$2c^2$	[HM'12]
All-1 attack	$O(c^{1.5})$	[OSD'13]	$\Omega(c)$	
Minority voting	$O(c^{1.5})$	[OSD'13]	$\Omega(c)$	
Majority voting	$O(c^{1.5})$	[OSD'13]	$\Omega(c)$	
Coin-flip attack	$O(c^{1.5})$	[OSD'13]	$\Omega(c)$	

# Results

## Results from group testing

Table: Asymptotics of  $L = \frac{\ell}{\ln n}$  for large  $n$ .

	Efficient constr.		Lower bounds	
Arbitrary attacks	$2c^2$	[OSD'13]	$2c^2$	[HM'12]
Interleaving attack	$2c^2$	[OSD'13]	$2c^2$	[HM'12]
All-1 attack	$ec$	[C+'11]	$\log_2(e)c$	[Seb'85]
Minority voting	$O(c^{1.5})$	[OSD'13]	$\Omega(c)$	
Majority voting	$O(c^{1.5})$	[OSD'13]	$\Omega(c)$	
Coin-flip attack	$O(c^{1.5})$	[OSD'13]	$\Omega(c)$	

# Results

## Contributions

Table: Asymptotics of  $L = \frac{\ell}{\ln n}$  for large  $n$ .

	Efficient constr.		Lower bounds	
Arbitrary attacks	$2c^2$	[OSD'13]	$2c^2$	[HM'12]
Interleaving attack	$2c^2$	[OSD'13]	$2c^2$	[HM'12]
All-1 attack	$2c$	[Laa'13]	$\log_2(e)c$	[Seb'85]
Minority voting	$2c$	[Laa'13]	$\Omega(c)$	
Majority voting	$\pi c$	[Laa'13]	$\Omega(c)$	
Coin-flip attack	$4c$	[Laa'13]	$\Omega(c)$	

# Efficient probabilistic group testing

Problem: Blood testing



Antonino

Boris

Caroline

David

Eve

Fred

Gábor

Henry

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# Efficient probabilistic group testing

Problem: Blood testing

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Antonino	1	0	0	0	0	0	0	0	0	0	0
Boris	0	1	0	0	0	0	0	0	0	0	0
Caroline	0	0	1	0	0	0	0	0	0	0	0
David	0	0	0	1	0	0	0	0	0	0	0
Eve	0	0	0	0	1	0	0	0	0	0	0
Fred	0	0	0	0	0	1	0	0	0	0	0
Gábor	0	0	0	0	0	0	1	0	0	0	0
Henry	0	0	0	0	0	0	0	0	1	0	0

---












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# Efficient probabilistic group testing

Problem: Blood testing

---

											
Antonino	1	0	0	0	0	0	0	0			
Boris	0	1	0	0	0	0	0	0			
Caroline	0	0	1	0	0	0	0	0			
David	0	0	0	1	0	0	0	0			
Eve	0	0	0	0	1	0	0	0			
Fred	0	0	0	0	0	1	0	0			
Gábor	0	0	0	0	0	0	1	0			
Henry	0	0	0	0	0	0	0	1			
Results	0	0	1	0	0	0	0	0			

---

# Efficient probabilistic group testing

Solution: Using pools

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	1	2	3
Antonino	0	0	0
Boris	0	0	1
Caroline	0	1	0
David	0	1	1
Eve	1	0	0
Fred	1	0	1
Gábor	1	1	0
Henry	1	1	1

---

Results

---















# Efficient probabilistic group testing

Solution: Group testing

	1	2	3	4	5	6	7	8	9	10	11
Antonino	?	?	?	?	?	?	?	?	?	?	?
Boris	?	?	?	?	?	?	?	?	?	?	?
Caroline	?	?	?	?	?	?	?	?	?	?	?
David	?	?	?	?	?	?	?	?	?	?	?
Eve	?	?	?	?	?	?	?	?	?	?	?
Fred	?	?	?	?	?	?	?	?	?	?	?
Gábor	?	?	?	?	?	?	?	?	?	?	?
Henry	?	?	?	?	?	?	?	?	?	?	?
Results	?	?	?	?	?	?	?	?	?	?	?

1. An algorithm to construct group testing matrices

# Efficient probabilistic group testing

Solution: Group testing

	1	2	3	4	5	6	7	8	9	10	11
Antonino	?	?	?	?	?	?	?	?	?	?	?
Boris	?	?	?	?	?	?	?	?	?	?	?
Caroline	?	?	?	?	?	?	?	?	?	?	?
David	?	?	?	?	?	?	?	?	?	?	?
Eve	?	?	?	?	?	?	?	?	?	?	?
Fred	?	?	?	?	?	?	?	?	?	?	?
Gábor	?	?	?	?	?	?	?	?	?	?	?
Henry	?	?	?	?	?	?	?	?	?	?	?
Results	?	?	?	?	?	?	?	?	?	?	?

1. An algorithm to construct group testing matrices
2. An algorithm to link test results to infected people

# Efficient probabilistic group testing

Solution: Group testing

1. An algorithm to construct group testing matrices
2. An algorithm to link test results to infected people

## Efficient probabilistic group testing

Solution: Group testing

1. An algorithm to construct group testing matrices
2. An algorithm to link test results to infected people

# Score-based construction

## Overview

1. An algorithm to construct group testing matrices
2. An algorithm to link test results to infected people

# Score-based construction

## Overview

1. An algorithm to construct group testing matrices
  - 1a. For each test  $i$ , generate  $p_i \sim F$ .
  - 1b. For each test  $i$ , person  $j$ , choose  $X_{j,i} = 1$  with prob.  $p_i$ .
2. An algorithm to link test results to infected people
  - 2a. For each test  $i$ , person  $j$ , calculate  $S_{j,i} = g(X_{j,i}, y_i, p_i)$ .
  - 2b. Mark person  $j$  infected iff  $\sum_i S_{j,i}$  is “large”.



## Score-based construction

The classical model  $\cong$  The all-1 attack

1. An algorithm to construct group testing matrices
  - 1a. For each test  $i$ , generate  $p_i \sim F$ .
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## Score-based construction

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$$F(p_i) = \mathbf{1} \left\{ p_i \geq p \approx \frac{1}{c} \right\}$$

$$g(X_{j,i}, y_i, p_i) = \begin{cases} +p_i/(1-p_i) & (X_{j,i}, y_i) = (0, 0) \\ -p_i(1-p_i)^{c-1}/(1-(1-p_i)^c) & (X_{j,i}, y_i) = (0, 1) \\ -1 & (X_{j,i}, y_i) = (1, 0) \\ +(1-p_i)^c/(1-(1-p_i)^c) & (X_{j,i}, y_i) = (1, 1) \end{cases}$$

$$\ell \sim 2c \ln n$$

# Results

## Previous results from group testing

Table: Asymptotics of  $L = \frac{\ell}{\ln n}$  for large  $n$ .

	Efficient constr.		Lower bounds
Classical model	$ec$	[C+'11]	$\log_2(e)c$ [Seb'85]
Noisy model	$\frac{4.36c}{(1-r)^2}$	[C+'11]	$\Omega\left(\frac{c}{(1-r)^2}\right)$ [AS'09]
Majority model	$O(c)$		$\Omega(c)$
Bernoulli gap	$\frac{4e^8 \ln(2)c}{\pi^2}$	[C+'13]	$\Omega(c)$
Linear gap	$2 \log_2(e)c^2$	[D+'05]	$\Omega(c)$

# Results

## Contributions

Table: Asymptotics of  $L = \frac{\ell}{\ln n}$  for large  $n$ .

	Efficient constr.		Lower bounds
Classical model	$2c$	[Laa'13]	$\log_2(e)c$ [Seb'85]
Noisy model	$\frac{2c}{(1-r)^2}$	[Laa'13]	$\Omega\left(\frac{c}{(1-r)^2}\right)$ [AS'09]
Majority model	$\pi c$	[Laa'13]	$\Omega(c)$
Bernoulli gap	$4c$	[Laa'13]	$\Omega(c)$
Linear gap	$2c^2$	[OSD'13]	$2c^2$ [HM'12]

## Conclusion

### **Fighting against specific attacks in traitor tracing**

- If you know the attack, you can often find pirates much faster!
- Trick: Not only optimize  $g$ , but also  $p$
- Code length often linear in  $c$  with small constants

### **Applications to probabilistic group testing**

- Group testing models  $\cong$  Specific attacks in traitor tracing
- Classical model: Asymptotic improvement over best result
- Improvements for various other models as well

Questions?